

Primality testing: Given an integer  $N$  is  $N$  prime?

See for  $i = 1 \dots \sqrt{N}$  if  $i$  divides  $N$

$O(\sqrt{N})$

$\uparrow$   
 $2^{\frac{n}{2}}$

Input size:  $\log N$  bits  
 $= n$

density of primes: Among  $n$  integers  
 $\sim \frac{n}{\ln n}$  are prime

Pick a large random integer  $k$   
test if  $k$  is prime

Efficient algorithms for primality tests:  
Rabin-Miller      Solovay Strassen

(Monte Carlo) Randomized:  $\sim n^2$

AKS algorithm deterministic:  $n^{12}$

# Selection Problem

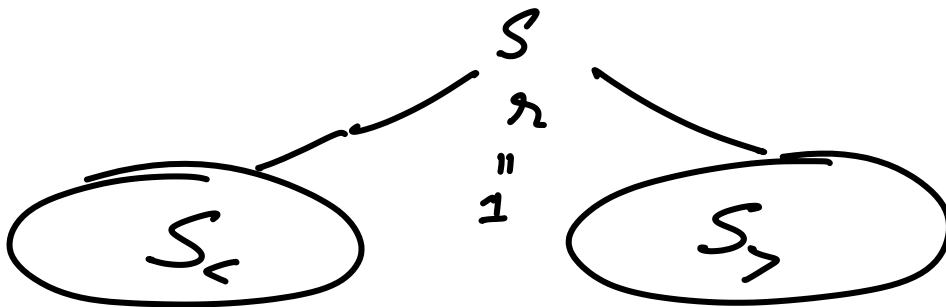
Given  $n$  elements from a set  $S$  (totally ordered) and an integer  $1 \leq k \leq n$ , find an element  $x \in S$  such that  $\text{rank}(x, S) = k$

# elements  $y \in S$   
s.t.  $y < x$

Objective: Design a linear time algorithm for selection

Easy soln: to sort

Pick some element  $r \in S$  and partition  $S$  using  $r$



Prune and search

Analysis: If the subset containing  
the rank  $k$  element decreases  
by some constant factor, say  $\alpha < 1$   
running time:  $\sum_{i \geq 0} \alpha^i n = O(n)$

Approximate median: If the  
rank of  $r$  is between

$$\left[ \frac{1}{4}n, \frac{3}{4}n \right]$$

...  $r$ : random element of  $S$   
 $r$  is an approx median with prob  $\geq \frac{1}{2}$

Find rank of  $r$

If rank is in  $\left[ \frac{1}{4}n, \frac{3}{4}n \right]$

then proceed with the correct  
partition

else

Expected # of iterations = 2

In each iteration we spend  $cn$  time

First	recursive	call	$2cn$	$X_1$
Sec.	"	"	$2c \cdot \frac{3}{4}n$	$X_2$
$\vdots$				
$i$	"	"	$2c \left(\frac{3}{4}\right)^i n$	$X_m$

$$m = \log_{4/3} n$$

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$$O(cn)$$

$X$ : random variable representing the overall running time

$$\begin{aligned} E[X] &= E\left[cn \cdot X_1 + c \frac{3}{4}n \cdot X_2 + \dots + \left(\frac{3}{4}\right)^{m-1} cn \cdot X_m\right] \\ &= cn E[X_1] + \frac{3}{4}cn E[X_2] + \dots + \left(\frac{3}{4}\right)^{m-1} cn E[X_m] \end{aligned}$$

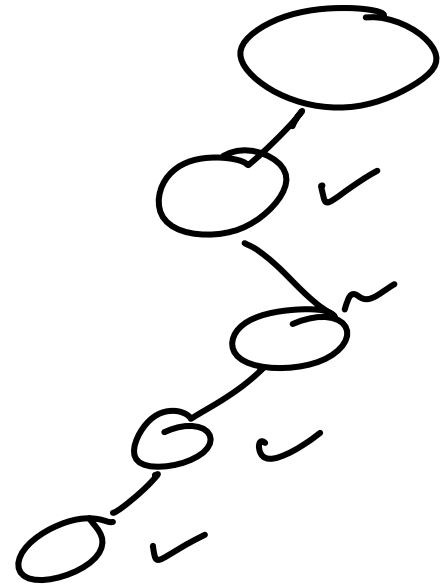
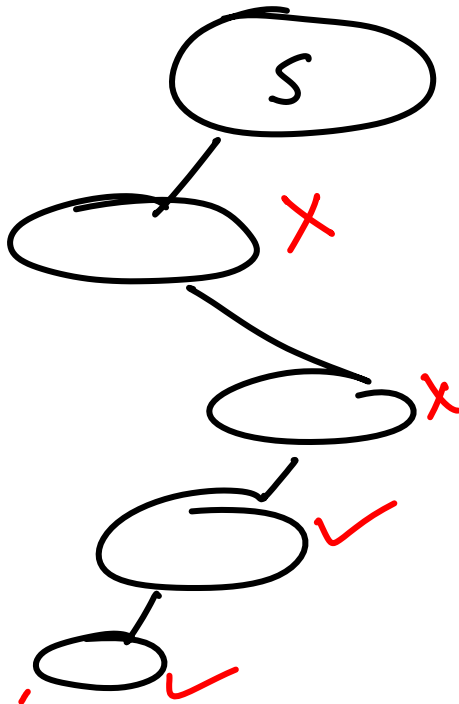
using linearity of expectation

$$E[X+Y] = E[X] + E[Y]$$

for "any" n.v.  $X, Y$  (not necessarily independent)

$$E[X] \text{ is } O(cn)$$

No t repeating sampling



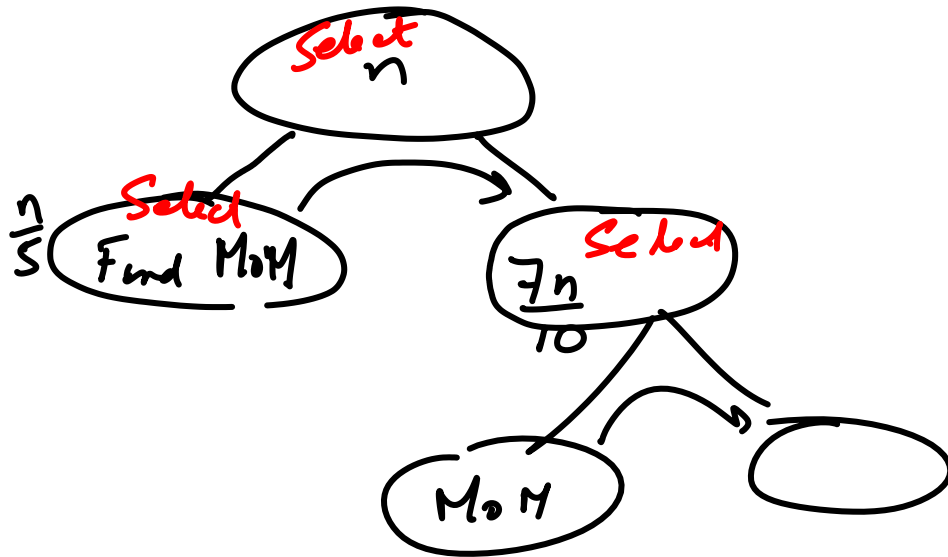
Median-of-medians

for approximate median

$n$

			→			
$G_1$	x	x		x	x	x
$G_2$	x	x		x	x	x
$\vdots$						
$G_{\frac{n}{5}}$	x	x		x	x	x

Choose the median of medians and that is  $n$



$$T(n) \leq T\left(\frac{n}{5}\right) + T\left(\frac{7n}{10}\right) + O(n)$$

Solve by guessing  $T = cn$  for some  
 appropriate constant  $c$

$\sim c$

Analyse the space complexity of this algorithm