

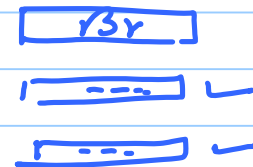
Oct. 18.

## Branches: Delayed Branch

Normal  
Branch.



Delayed  
Branch



+ Higher IPC if delay slots are filled

- Over-emphasis on the compiler.

Predict the direction of the branch:

	Taken (T) (NPC $\neq$ PC+4)	Not Taken (NT) (NPC = PC+4)
--	-----------------------------------	-----------------------------------

IF	ID	EX	MEM	WB
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.

.

.

~~X~~

~~X~~

Correct  
Wrong

branch  
branch

prediction  
"

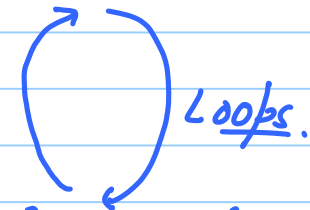
→ do nothing

→ cancel the incorrectly  
fetched instructions  
fetch from correct address

+ compiler independent

How do we predict branches?

Typical program



Predict : always taken (mostly correct)  
(70% acc.)



1) Loop loss  $\rightarrow$  50%.

2) always taken  $\rightarrow$  70%.

2-bit saturating counter (c)  
(implemented in HW)

inc()      dec()  
            ↓  
get value (c)  
            0, 1, 2, 3

```
C. inc() {
```

```
    value = min(value + 1, 3);
```

```
}
```

```
C. dec() {
```

```
    value = max(value - 1, 0);
```

```
}
```

1) Simple Predictor:

Branch prediction (bp)

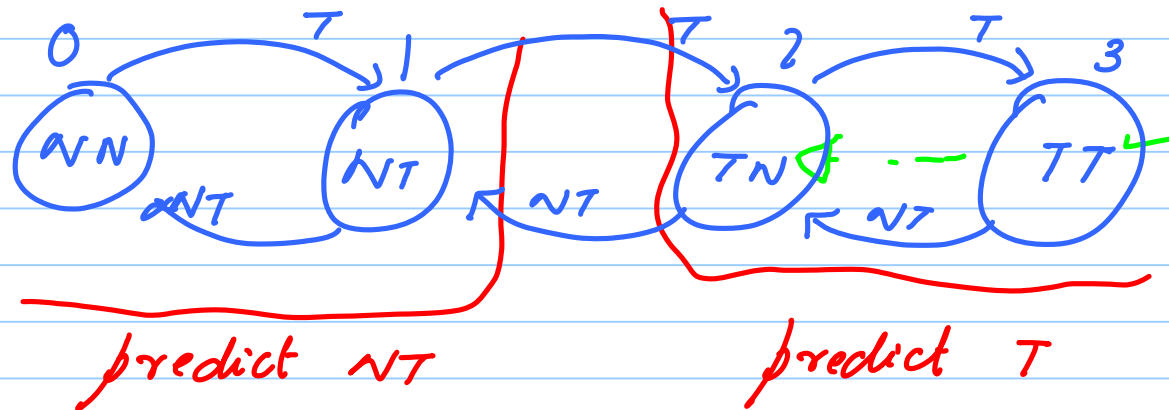
bp = same output as last branch.



1st pred: x  
 .  
 .  
 (N-1)<sup>th</sup> pred }  
 N<sup>th</sup> pred x } NT

1st pred → x } NT  
 .  
 .  
 .  
 .  
 nth pred → x

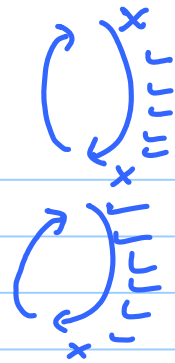
predictor, which does not take exceptions very seriously.



Branch Predictor  $\rightarrow$  predict (pc)  $\rightarrow$  [T, NT] (IF)  
 $\rightarrow$  Train (pc, [T/NT]) (EX)

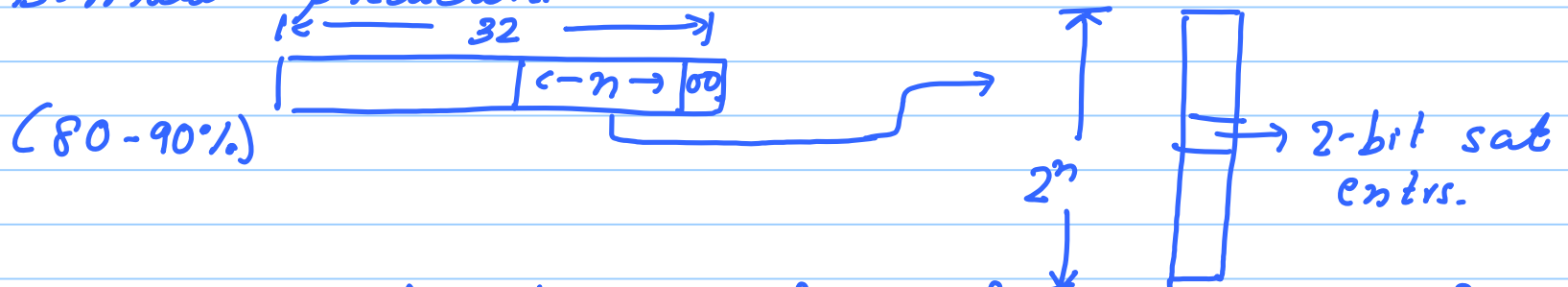
2-bit sat entr. adds some memory to the branch predictor.

For the example with two loops, makes one less mistake.



✗ You cannot have only one  
 2-bit sat entrs. (collisions)  
 ✗ Ideally,  
 1 2-bit sat entrs per  
 branch (Too much  
 space)

### 1) Bimodal predictor.



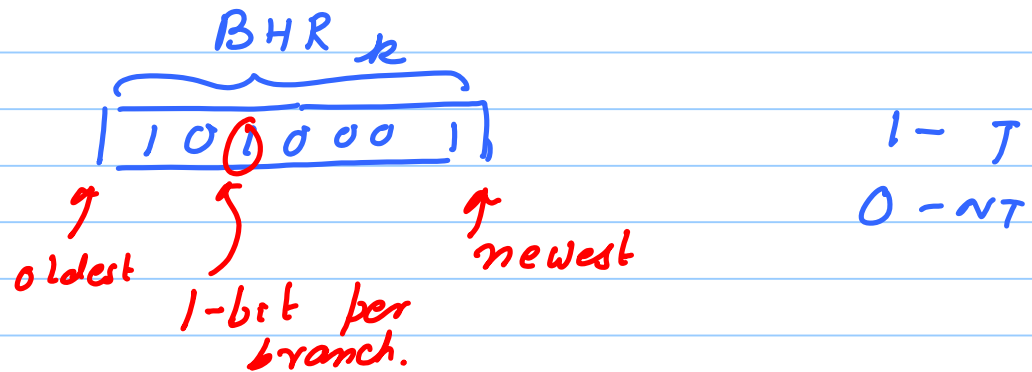
- Taking a look at the past behavior of a branch PC
- Saving the result in the 2-bit entrs.
- Using it for pred.

+ reduces collisions.

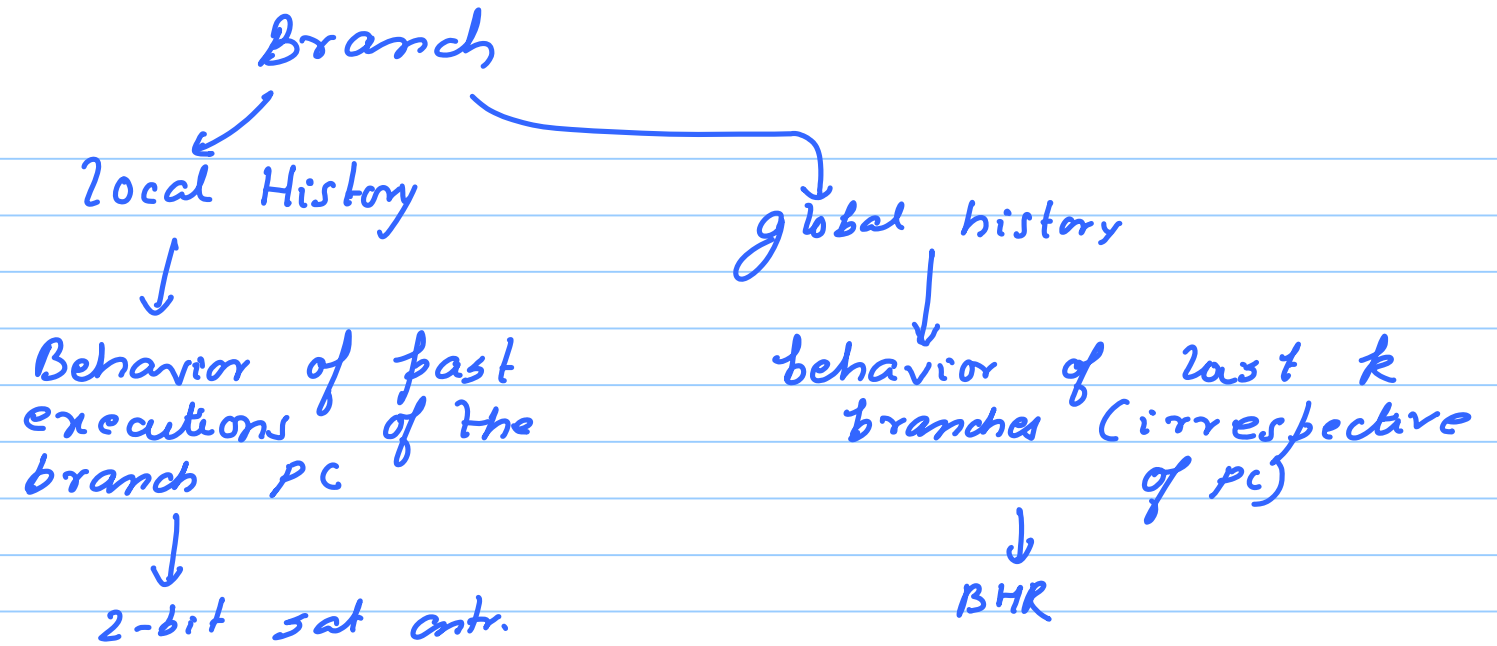
```
{ if (flag1 == 0) (B1)
  flag2 = 1;
  if (flag2 == 1) (B2)
    printf("great");
```

(B1 is NT)  $\Rightarrow$  (B2 is NT)

Branch History Register:



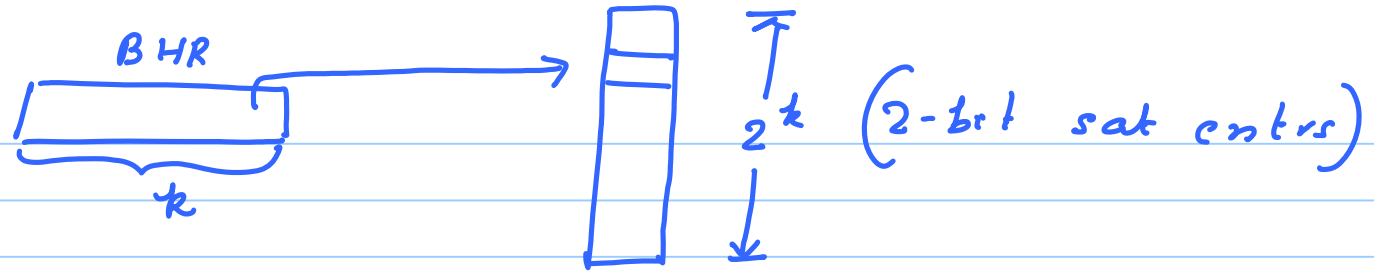




Family of branch preds.

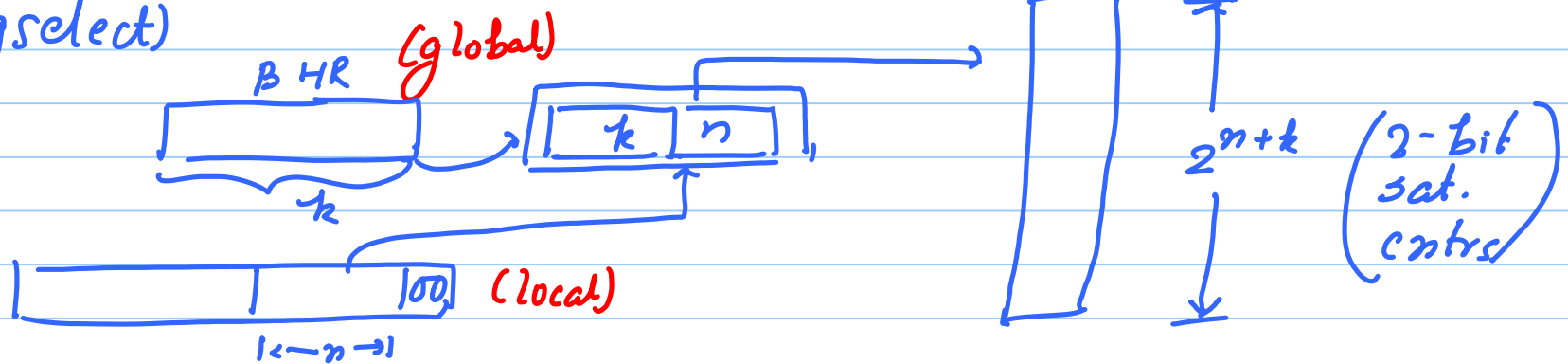
G → global  
P → pattern (local)

Gag:  
(90%)

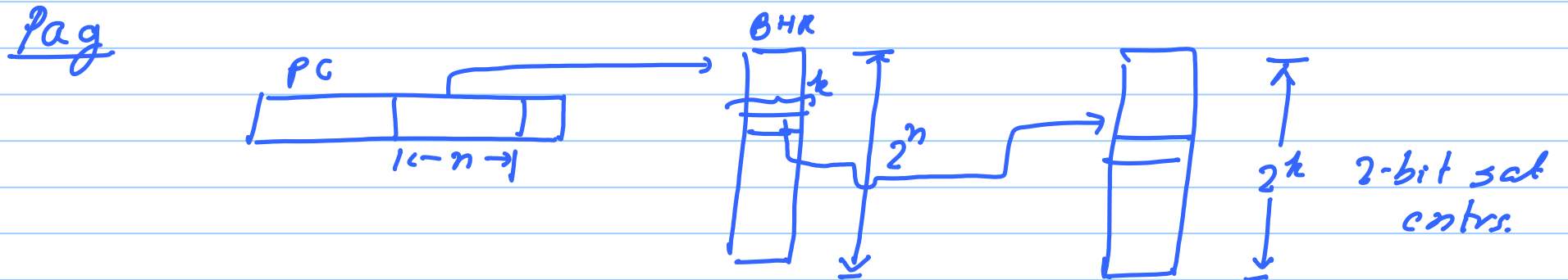
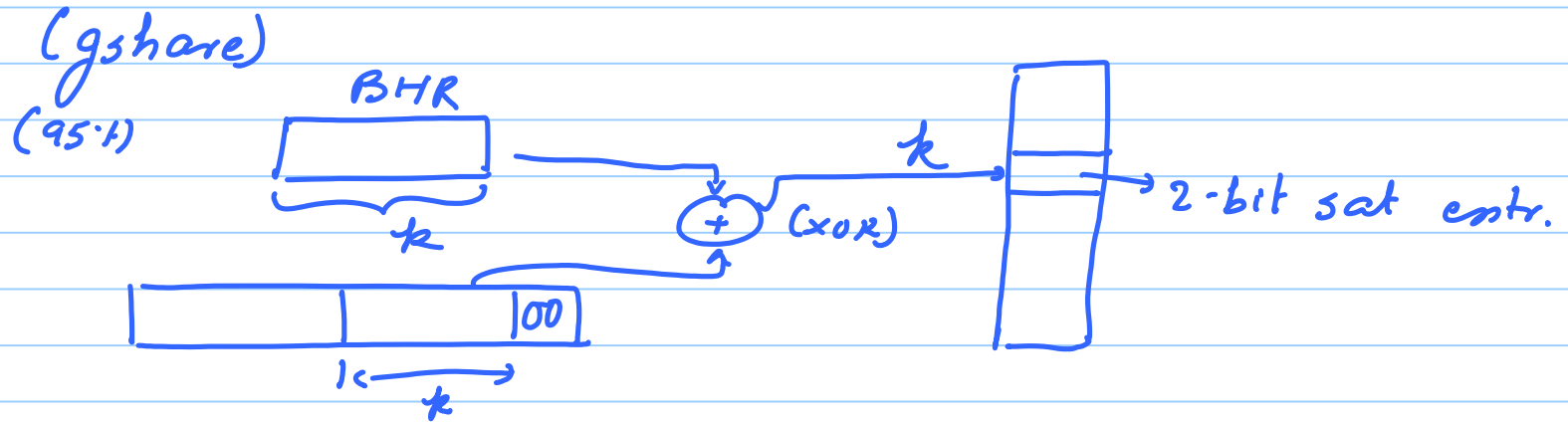


- + simple
- not taking local history into account

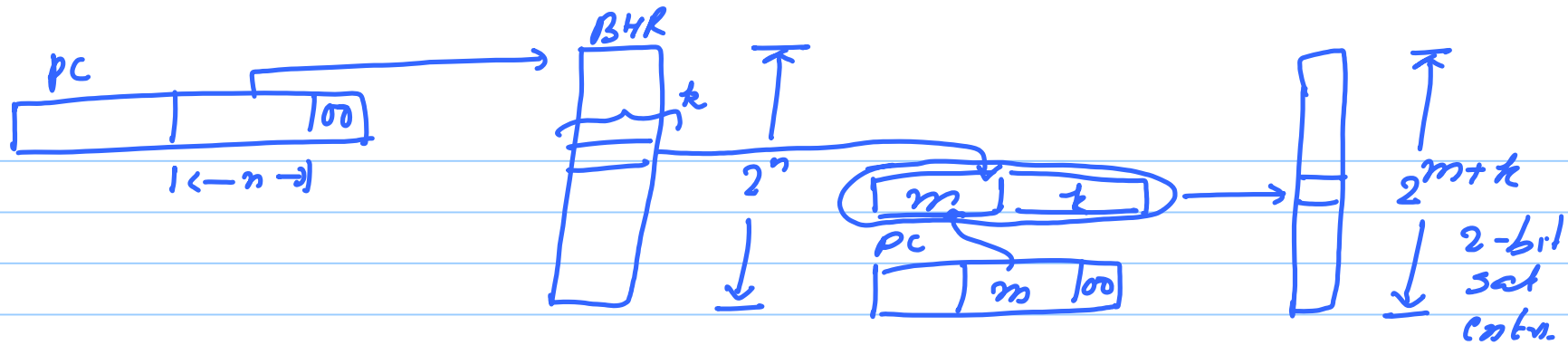
Gag (gselect)  
(95%)



+ combines both global & local  
 - needs more space ( $2^{n+k}$ )

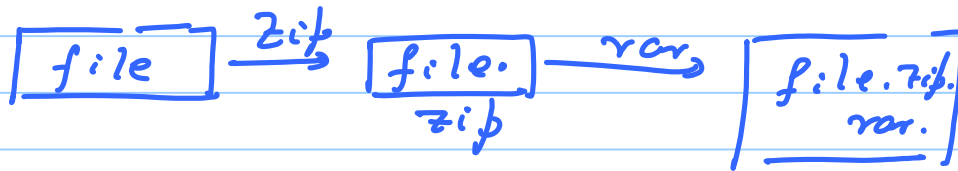


PaP

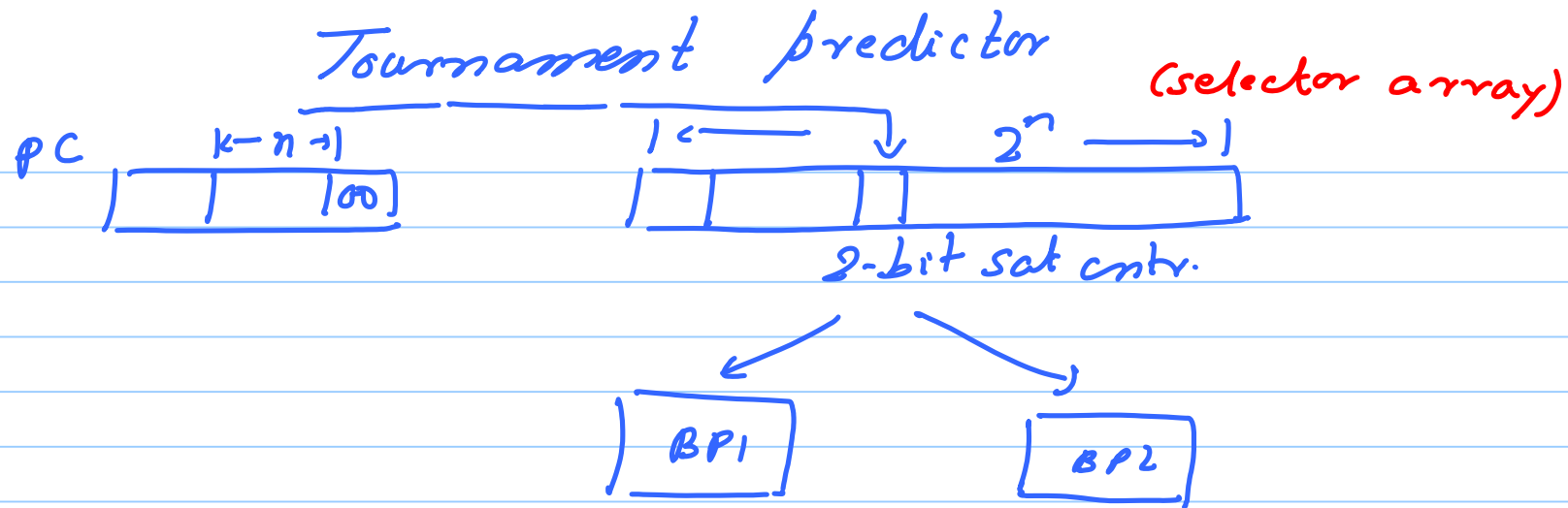


Theory of branch preds:

Prediction is related to compression  
information theory



(Fano's inequality)



Predict:

1) Find 2-bit sat cntr. <sup>(c)</sup> in the selector array.

2) if  $v = \text{value}(c)$

$v = [0, 1]$       predict BP<sub>1</sub>, predict()

$v = [2, 3]$  predict BP<sub>2</sub>. predict()

Train:

1) Train the selector array.

2) Train BP<sub>1</sub> and BP<sub>2</sub>