



**COL333/671: Introduction to AI**  
Semester I, 2021

**Local Search Algorithms**

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# Outline

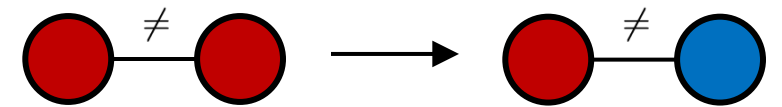
- Last Class
  - Constraint Satisfaction Problems
- This Class
  - Local Search Algorithms
- Reference Material
  - AIMA Ch. 4.1

# Acknowledgement

**These slides are intended for teaching purposes only. Some material has been used/adapted from web sources and from slides by Doina Precup, Dorsa Sadigh, Percy Liang, Mausam, Dan Klein, Nicholas Roy and others.**

# Iterative Approaches to Solving CSPs

- Local search methods
  - Keep track of “complete” states, i.e., all variables assigned.
  - If there are conflicts, then try to improve the assignment iteratively till a solution is obtained.
- Local search for CSPs:
  - Take an assignment with unsatisfied constraints
  - Reassign variable values
  - Repeat: Till the CSP does not have a solution,
    - Variable selection: randomly select any conflicted variable
    - Value selection: min-conflicts heuristic:
      - Choose a value for the variable that violates the fewest constraints
      - I.e., “hill climb” with  $h(x)$  = total number of violated constraints
  - Hill climbing (a general technique for search)
    - Start at a state
    - Repeat: move to the best neighboring state
    - If no neighbors better than current, quit

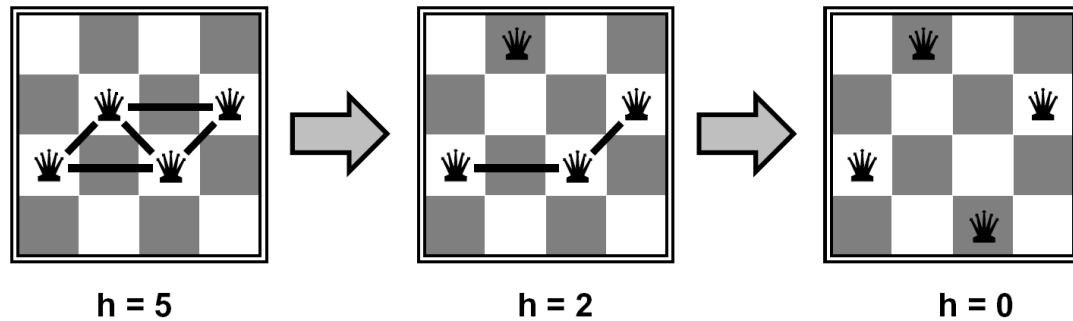


*Keep track of complete assignments and explore at local changes.*

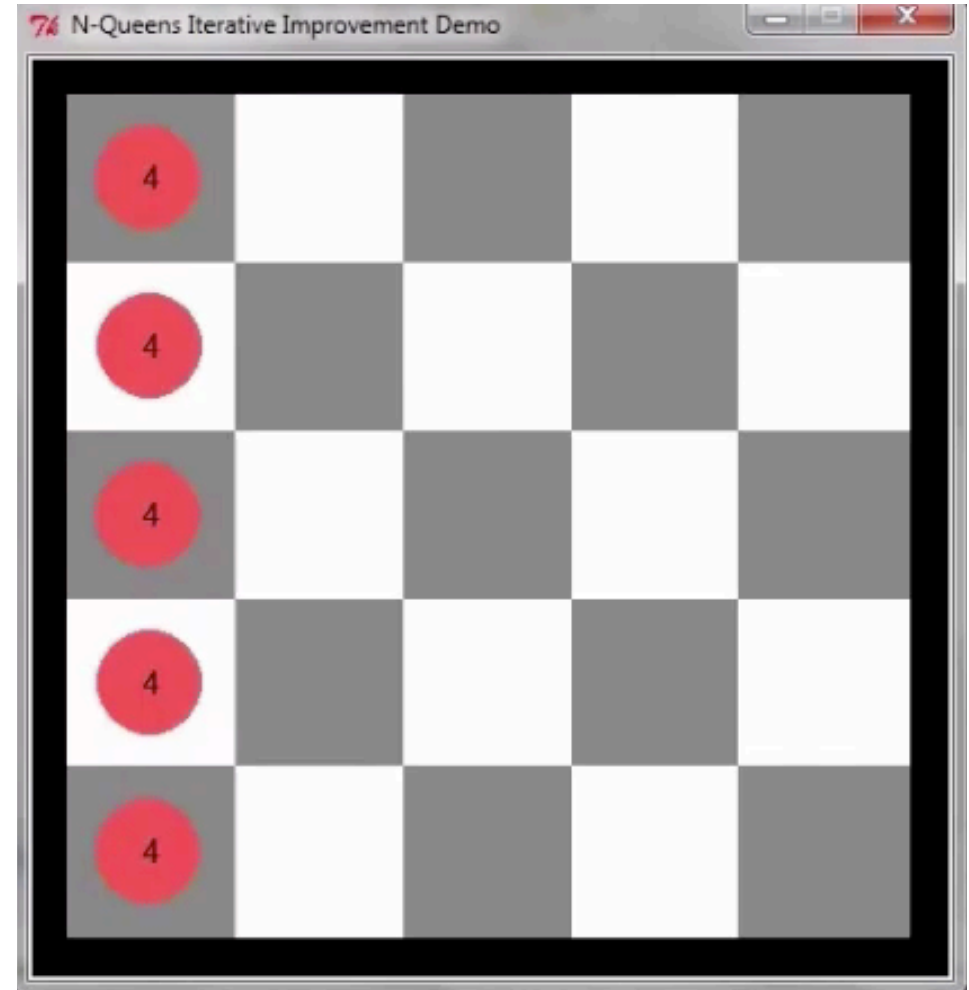


*Hill climbing*

# Example: 4 Queens

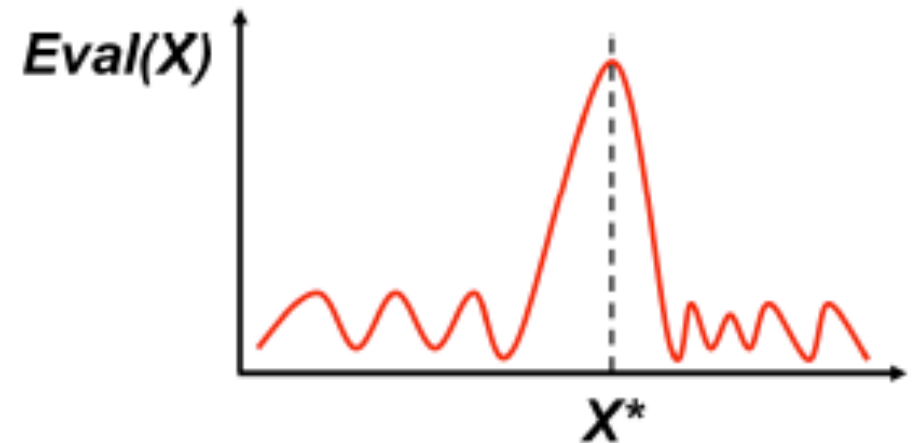


- States: 4 queens in 4 columns ( $4^4 = 256$  states)
- Operators: move queen in column
- Goal test: no attacks
- Evaluation:  $h(x)$  = number of attacks (number of violated binary constraints)

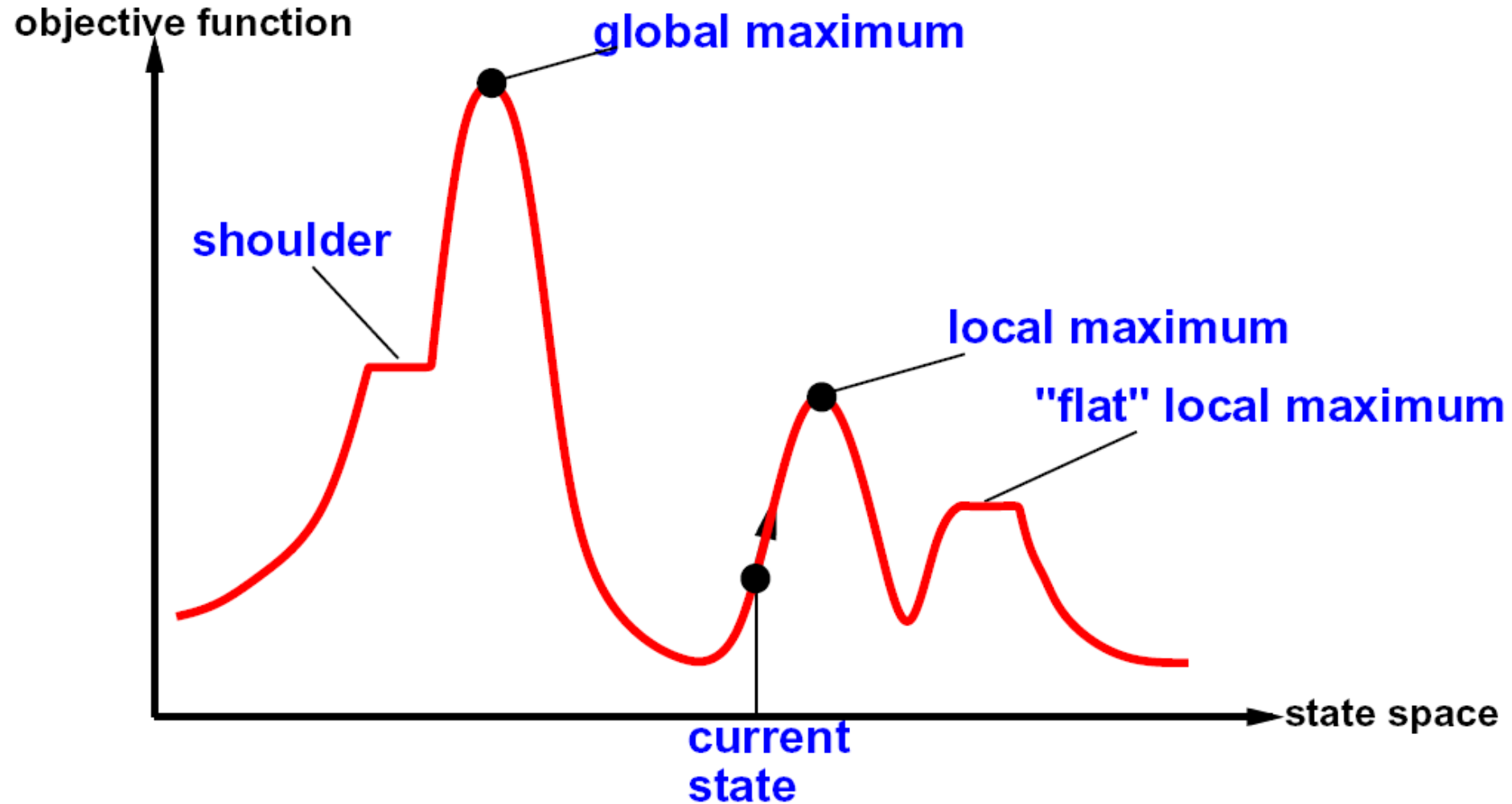


# Optimization Problems: Generic Setup

- There is (discrete) combinatorial structure to the problem (maybe with constraints).
- Cost function, which we want to optimize.
- Searching all possible solutions is likely to be infeasible.
  - At least, we want a “good” solution.



# Optimization Landscape



# Simulated Annealing

- Allows some apparently **bad moves** - to escape local maxima.
- **Decrease** the size and the frequency of bad moves over time.

- Algorithm sketch

1. Start at initial configuration  $X$  of value  $E$  (high is good)

2. Repeat:

- (a) Let  $X_i$  be a *random neighbor* of  $X$  and  $E_i$  be its value

- (b) If  $E < E_i$  then let  $X \leftarrow X_i$  and  $E \leftarrow E_i$

- (c) Else, *with some probability  $p$* , still accept the move:  $X \leftarrow X_i$  and  $E \leftarrow E_i$

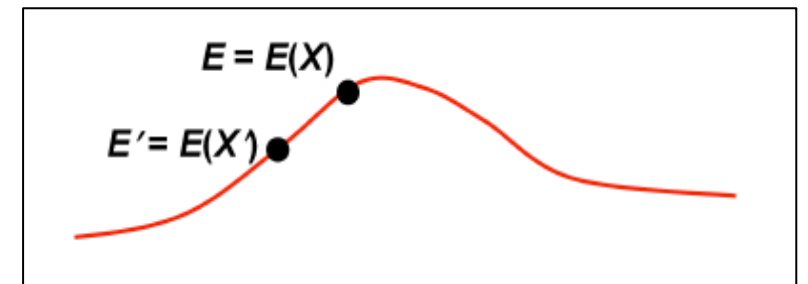
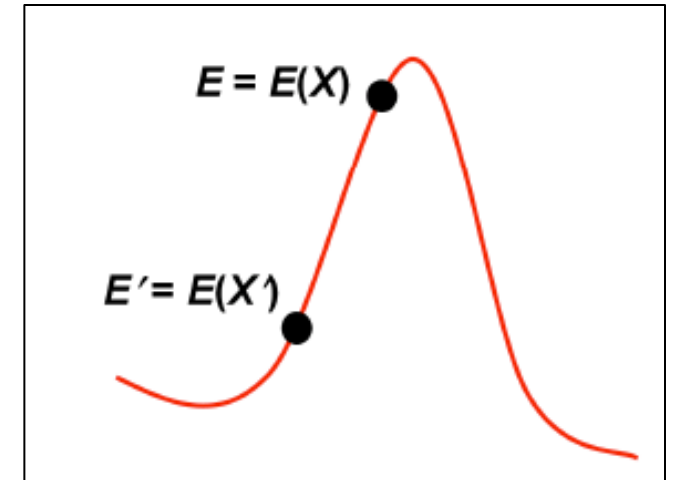
- Best solution ever found is always remembered

A form of Monte-Carlo Search. Move around the environment to explore it instead of systematically sweeping. Powerful technique for large domains.



# Simulated Annealing: How to decide $p$ ?

- Considering a move from state of value  $E$  to a lower valued state of  $E'$ .
- If  $(E - E')$  is large:
  - Likely to be close to a promising maximum.
  - Less inclined to go downhill.
- If  $(E - E')$  is small:
  - The closest maximum may be shallow
  - More inclined to go downhill is not as bad.



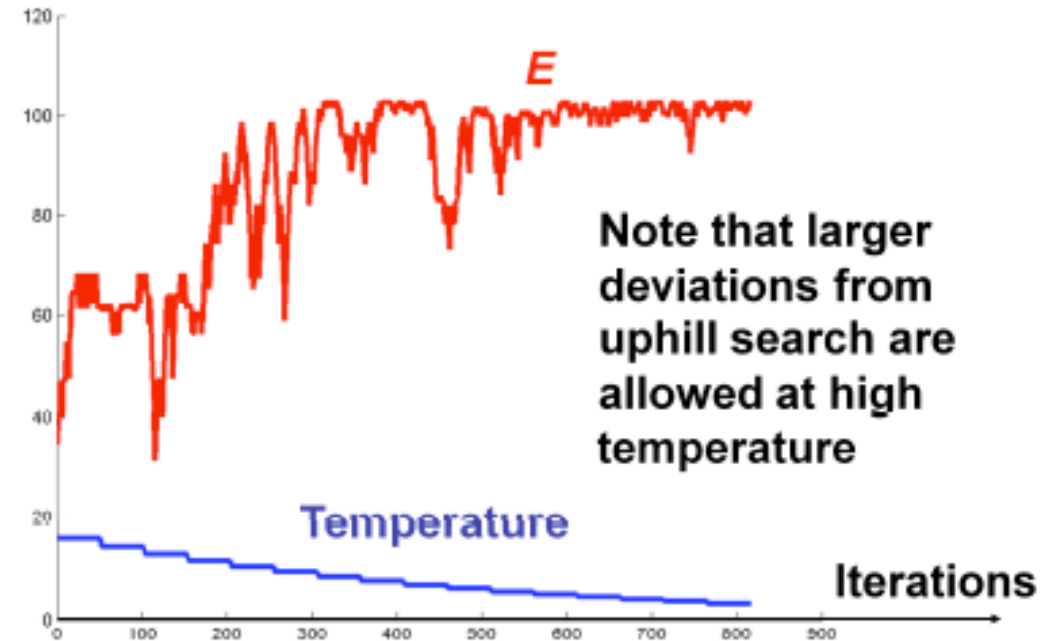
# Simulated Annealing: Selecting Moves

- If the new value  $E_i$  is **better** than the old value  $E$ , move to  $X_i$
- If the new value is **worse** ( $E_i > E$ ) then move to the neighboring solution as per *Boltzmann* distribution.
- Temperature ( $T > 0$ )
  - **T is high**,  $\exp$  is  $\sim 0$ , acceptance probability is  $\sim 1$ , high probability of acceptance of a worse solution.
  - **T is low**, the probability of moving to a worse solution is  $\sim 0$ , low probability of acceptance of a worse solution.
  - Schedule  $T$  to reduce over time.

$$\exp\left(-\frac{E - E_i}{T}\right)$$

# Simulated Annealing: Properties

- **T is high**
  - The algorithm is in an exploratory phase
  - Even bad moves have a high chance of being picked)
- **T is low**
  - The algorithm is in an exploitation phase
  - The “bad” moves have very low probability
- **If T is decreased slowly enough**
  - Simulated annealing is guaranteed to reach the best solution in the limit.



# Local Beam Search

- Look for solutions from multiple points in parallel.
- Algorithm
  - Track  $k$  states (rather than 1).
  - Begin with  $k$  randomly sampled states.
  - Loop
    - Generate successors of each of the  $k$ -states
    - If anyone has the goal, the algorithm halts
    - Otherwise, select only the  $k$ -best successors from the list and repeat.
  - Note:
    - Each run is not independent, information is passed between parallel search threads.
    - Promising states are propagated. Less promising states are not propagated.
    - Problem: states become concentrated in a small region of space.

# Stochastic Beam Search

- Local beam search
  - Problem: states become concentrated in a small region of space
  - Search degenerates to hill climbing
- Stochastic beam search
  - Instead of taking the best k states
  - Sample k states from a distribution
  - Probability of selecting a state *increases* as the *value* of the state.