

# COL202: Discrete Mathematical Structures

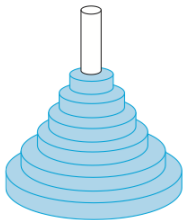
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## Advanced Counting Techniques

# Advanced Counting Techniques

## Recurrence relations

- Tower of Hanoi: Let  $H_n$  denote the number of moves needed to solve the Tower of Hanoi problem with  $n$  disks. Set up a recurrence relation for the sequence  $\{H_n\}$ .



Peg 1



Peg 2



Peg 3

# Advanced Counting Techniques

## Recurrence relations

- Find a recurrence relation and give initial conditions for the number of bit strings of length  $n$  that do not have two consecutive 0s. How many such bit strings are there of length five?

# Advanced Counting Techniques

## Recurrence relations

- Dynamic Programming: This is an algorithmic technique where a problem is recursively broken down into simpler overlapping subproblems, and the solution is computed using the solutions of the subproblems.
- Problem: Given a sequence of integers, find the length of the *longest increasing subsequence* of the given sequence.
  - Example: The longest increasing subsequence of the sequence  $(7, 2, 8, 10, 3, 6, 9, 7)$  is  $(2, 3, 6, 7)$  and its length is 4.

# Advanced Counting Techniques

## Solving recurrence relations

### Definition (Linear homogeneous recurrence)

A *linear homogeneous* recurrence relation of degree  $k$  with constant coefficients is a recurrence relation of the form

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k},$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$ .

- *Linear* means that that RHS is a sum of linear terms of the previous elements of the sequence.
  - $a_n = a_{n-1} + a_{n-2}$  is a linear recurrence relation whereas  $a_n = a_{n-1} + a_{n-2}^2$  is not.

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where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$ .

- *Linear* means that that RHS is a sum of linear terms of the previous elements of the sequence.
- *Homogeneous* means that there are no terms in the RHS that are not multiples of  $a_j$ 's.
  - $a_n = a_{n-1} + a_{n-2}$  is homogeneous whereas  $a_n = a_{n-1} + a_{n-2} + 2$  is not.

# Advanced Counting Techniques

## Solving recurrence relations

### Definition (Linear homogeneous recurrence)

A *linear homogeneous* recurrence relation of degree  $k$  with constant coefficients is a recurrence relation of the form

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where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$ .

- *Linear* means that that RHS is a sum of linear terms of the previous elements of the sequence.
- *Homogeneous* means that there are no terms in the RHS that are not multiples of  $a_j$ 's.
- The coefficients of all the terms on the RHS are constants.
- The degree is  $k$  since  $a_n$  is expressed as the previous  $k$  terms of the sequence.



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## Solving recurrence relations

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where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$ .

- $a_n = r^n$  is a solution of the recurrence if and only if

$$r^k - c_1 r^{k-1} - \dots - c_k = 0. \quad (1)$$

- (1) is called the *characteristic equation* of the recurrence relation.
- The solutions of the characteristic equation are called the *characteristic roots* of the recurrence relation.

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Let  $c_1$  and  $c_2$  be real numbers. Suppose  $r^2 - c_1r - c_2 = 0$  has two distinct roots  $r_1$  and  $r_2$ . Then the sequence  $\{a_n\}$  is a solution of the linear homogeneous recurrence relation  $a_n = c_1a_{n-1} + c_2a_{n-2}$  if and only if  $a_n = \alpha_1r_1^n + \alpha_2r_2^n$  for all  $n = 0, 1, 2, \dots$ , where  $\alpha_1$  and  $\alpha_2$  are constants.

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## Solving recurrence relations

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- What is the solution of the recurrence relation  $a_n = a_{n-1} + 2 \cdot a_{n-2}$  with  $a_0 = 2$  and  $a_1 = 7$ ?

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Let  $c_1$  and  $c_2$  be real numbers. Suppose  $r^2 - c_1r - c_2 = 0$  has two distinct roots  $r_1$  and  $r_2$ . Then the sequence  $\{a_n\}$  is a solution of the linear homogeneous recurrence relation  $a_n = c_1a_{n-1} + c_2a_{n-2}$  if and only if  $a_n = \alpha_1r_1^n + \alpha_2r_2^n$  for all  $n = 0, 1, 2, \dots$ , where  $\alpha_1$  and  $\alpha_2$  are constants.

### Theorem

Let  $c_1$  and  $c_2$  be real numbers with  $c_2 \neq 0$ . Suppose that  $r^2 - c_1r - c_2 = 0$  has only one root  $r_0$ . A sequence  $\{a_n\}$  is a solution of the recurrence relation  $a_n = c_1a_{n-1} + c_2a_{n-2}$  if and only if  $a_n = \alpha_1r_0^n + \alpha_2nr_0^n$ , for  $n = 0, 1, 2, \dots$ , where  $\alpha_1$  and  $\alpha_2$  are constants.

- What is the solution of the recurrence relation  $a_n = 6a_{n-1} - 9 \cdot a_{n-2}$  with  $a_0 = 1$  and  $a_1 = 6$ ?

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Let  $c_1, c_2, \dots, c_k$  be real numbers. Consider the linear homogeneous recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ . Suppose the characteristic equation of the recurrence relation has  $k$  distinct characteristic roots  $r_1, r_2, \dots, r_k$ . Then  $\{a_n\}$  is a solution of the recurrence relation if and only if  $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n + \dots + \alpha_k r_k^n$  for  $n = 0, 1, 2, \dots$ , where  $\alpha_1, \alpha_2, \dots, \alpha_k$  are constants.

- What is the solution of the recurrence relation

$$a_n = 6a_{n-1} - 11 \cdot a_{n-2} + 6a_{n-3} \text{ with } a_0 = 2, a_1 = 5, \text{ and } a_2 = 15?$$

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Let  $c_1, c_2, \dots, c_k$  be real numbers. Consider the linear homogeneous recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ . Suppose the characteristic equation of the recurrence relation has  $t \leq k$  distinct characteristic roots  $r_1, r_2, \dots, r_t$  with multiplicities  $m_1, m_2, \dots, m_t$ , respectively, so that  $m_i \geq 1$  for  $i = 1, 2, \dots, t$  and  $m_1 + m_2 + \dots + m_t = k$ . Then  $\{a_n\}$  is a solution of the recurrence relation if and only if

$$\begin{aligned} a_n = & (\alpha_{1,0} + \alpha_{1,1}n + \dots + \alpha_{1,m_1-1}n^{m_1-1})r_1^n \\ & + (\alpha_{2,0} + \alpha_{2,1}n + \dots + \alpha_{2,m_2-1}n^{m_2-1})r_2^n \\ & + \dots + (\alpha_{t,0} + \alpha_{t,1}n + \dots + \alpha_{t,m_t-1}n^{m_t-1})r_t^n \end{aligned}$$

for  $n = 0, 1, 2, \dots$ , where  $\alpha_{i,j}$  are constants for  $1 \leq i \leq t$  and  $0 \leq j \leq m_i - 1$ .

- What is the solution of the recurrence relation

$$a_n = -3a_{n-1} - 3 \cdot a_{n-2} - a_{n-3} \text{ with } a_0 = 1, a_1 = -2, \text{ and } a_2 = -1?$$

# Advanced Counting Techniques

## Solving recurrence relations

- A *linear non-homogeneous recurrence relation with constant coefficients* is a recurrence of the form:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

where  $F(n)$  is a function not identically equal to zero and depending only on  $n$ .

- The recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$  is called the *associated homogeneous recurrence relation*.

### Theorem

If  $\{a_n^{(p)}\}$  is a particular solution of the non-homogeneous linear recurrence relation with constant coefficients

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

then every solution is of the form  $\{a_n^{(p)} + a_n^{(h)}\}$ , where  $\{a_n^{(h)}\}$  is a solution of the associated homogeneous recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}.$$

# Advanced Counting Techniques

## Solving recurrence relations

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$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}.$$

- Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ . What is the solution with  $a_1 = 3$ ?



# Advanced Counting Techniques

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- Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ . What is the solution with  $a_1 = 3$ ?
- Find all solutions if the recurrence relation  $a_n = 5a_{n-1} - 6a_{n-2} + 7^n$ .

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Suppose  $\{a_n\}$  satisfies the linear non-homogeneous recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and

$$F(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n,$$

where  $b_0, b_1, \dots, b_t$  are  $s$  real numbers. When  $s$  is not a root of the characteristic equation of the associated linear homogeneous recurrence relation, there is a particular solution of the form

$$(p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n.$$

When  $s$  is a root of this characteristic equation and its multiplicity is  $m$ , there is a particular solution of the form

$$n^m (p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n.$$

# Advanced Counting Techniques

## Solving recurrence relations

### Theorem

Suppose  $\{a_n\}$  satisfies the linear non-homogeneous recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

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# Advanced Counting Techniques

## Divide-and-conquer recurrence relations

### Theorem

Let  $f$  be an increasing function that satisfies the recurrence relation

$$f(n) = a \cdot f(n/b) + c$$

whenever  $n$  is divisible by  $b$ , where  $a \geq 1$ ,  $b$  is an integer greater than 1, and  $c$  is a positive real number. Then

$$f(n) \text{ is } \begin{cases} O(n^{\log_b a}) & \text{if } a > 1 \\ O(\log n) & \text{if } a = 1 \end{cases}$$

Furthermore, when  $n = b^k$  and  $a \neq 1$ , where  $k$  is a positive integer,  $f(n) = C_1 n^{\log_b a} + C_2$ , where  $C_1 = f(1) + c/(a - 1)$  and  $C_2 = -c/(a - 1)$ .

# Advanced Counting Techniques

## Divide-and-conquer recurrence relations

### Theorem (Master Theorem)

Let  $f$  be an increasing function that satisfies the recurrence relation

$$f(n) = a \cdot f(n/b) + cn^d$$

whenever  $n = b^k$ , where  $k$  is a positive integer,  $a \geq 1$ ,  $b$  is an integer greater than 1, and  $c$  and  $d$  are real numbers with  $c$  positive and  $d$  nonnegative. Then

$$f(n) \text{ is } \begin{cases} O(n^d) & \text{if } a < b^d \\ O(n^d \log n) & \text{if } a = b^d \\ O(n^{\log_b a}) & \text{if } a > b^d. \end{cases}$$

## Advanced Counting Techniques: Generating Functions

# Advanced Counting Techniques

## Generating functions

### Theorem (Generating function)

*The generating function for the sequence  $a_0, a_1, \dots, a_k, \dots$  of real numbers is the infinite series*

$$G(x) = a_0 + a_1x + \dots + a_kx^k + \dots = \sum_{k=0}^{\infty} a_kx^k$$

- We can define generating functions for finite sequences of real numbers by extending a finite sequence  $a_0, a_1, \dots, a_n$  into an infinite sequence by setting  $a_{n+1} = 0, a_{n+2} = 0$ , and so on.
- Examples:
  - What is the generating function for the sequence 1, 1, 1, 1, 1, 1?
  - Let  $m$  be a positive integer and let  $a_k = \binom{m}{k}$ , for  $k = 0, 1, \dots, m$ . What is the generating function for  $a_0, a_1, \dots, a_m$ ?

# Advanced Counting Techniques

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- Examples:
  - What is the generating function for the sequence 1, 1, 1, 1, 1, 1?
  - Let  $m$  be a positive integer and let  $a_k = \binom{m}{k}$ , for  $k = 0, 1, \dots, m$ . What is the generating function for  $a_0, a_1, \dots, a_m$ ?
  - The function  $f(x) = \frac{1}{1-x}$  is the generating function of the sequence 1, 1, ..., because  $\frac{1}{1-x} = 1 + x + x^2 + \dots$  for  $|x| < 1$ .



# Advanced Counting Techniques

## Generating functions

### Theorem

Let  $f(x) = \sum_{k=0}^{\infty} a_k x^k$  and  $g(x) = \sum_{k=0}^{\infty} b_k x^k$ . Then

$$f(x) + g(x) = \sum_{k=0}^{\infty} (a_k + b_k) x^k \text{ and}$$

$$f(x)g(x) = \sum_{k=0}^{\infty} \left( \sum_{j=0}^k a_j b_{k-j} \right) x^k.$$

- Let  $f(x) = \frac{1}{(1-x)^2}$ . Find coefficients  $a_0, a_1, \dots$  in the expansion  $f(x) = \sum_{k=0}^{\infty} a_k x^k$ .

# Advanced Counting Techniques

## Generating functions

### Definition (Extended binomial coefficient)

Let  $u$  be a real number and  $k$  a nonnegative integer. Then the extended binomial coefficient  $\binom{u}{k}$  is defined by

$$\binom{u}{k} = \begin{cases} \frac{u(u-1)\dots(u-k+1)}{k!} & \text{if } k > 0 \\ 1 & \text{if } k = 0 \end{cases}$$

- Find the value of the extended binomial coefficient  $\binom{1/2}{3}$ .
- Find the value of the extended binomial coefficient  $\binom{-n}{r}$ .

# Advanced Counting Techniques

## Generating functions

### Definition (Extended binomial coefficient)

Let  $u$  be a real number and  $k$  a nonnegative integer. Then the extended binomial coefficient  $\binom{u}{k}$  is defined by

$$\binom{u}{k} = \begin{cases} \frac{u(u-1)\dots(u-k+1)}{k!} & \text{if } k > 0 \\ 1 & \text{if } k = 0 \end{cases}$$

### Theorem (Extended binomial theorem)

Let  $x$  be a real number with  $|x| < 1$  and let  $u$  be a real number. Then

$$(1+x)^u = \sum_{k=0}^{\infty} \binom{u}{k} x^k.$$

- What is the expansion of  $(1-x)^{-n}$ ?

# Advanced Counting Techniques

## Generating functions

TABLE 1 Useful Generating Functions.	
$G(x)$	$a_k$
$(1+x)^n = \sum_{k=0}^n C(n,k)x^k$ $= 1 + C(n,1)x + C(n,2)x^2 + \dots + x^n$	$C(n,k)$
$(1+ax)^n = \sum_{k=0}^n C(n,k)a^k x^k$ $= 1 + C(n,1)ax + C(n,2)a^2x^2 + \dots + a^n x^n$	$C(n,k)a^k$
$(1+x^r)^n = \sum_{k=0}^n C(n,k)x^{kr}$ $= 1 + C(n,1)x^r + C(n,2)x^{2r} + \dots + x^{nr}$	$C(n,k/r)$ if $r \mid k$ ; 0 otherwise
$\frac{1-x^{n+1}}{1-x} = \sum_{k=0}^n x^k = 1 + x + x^2 + \dots + x^n$	1 if $k \leq n$ ; 0 otherwise
$\frac{1}{1-x} = \sum_{k=0}^{\infty} x^k = 1 + x + x^2 + \dots$	1
$\frac{1}{1-ax} = \sum_{k=0}^{\infty} a^k x^k = 1 + ax + a^2x^2 + \dots$	$a^k$
$\frac{1}{1-x^r} = \sum_{k=0}^{\infty} x^{kr} = 1 + x^r + x^{2r} + \dots$	1 if $r \mid k$ ; 0 otherwise
$\frac{1}{(1-x)^2} = \sum_{k=0}^{\infty} (k+1)x^k = 1 + 2x + 3x^2 + \dots$	$k+1$
$\frac{1}{(1-x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)x^k$ $= 1 + C(n,1)x + C(n+1,2)x^2 + \dots$	$C(n+k-1, k) = C(n+k-1, n-1)$
$\frac{1}{(1+x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)(-1)^k x^k$ $= 1 - C(n,1)x + C(n+1,2)x^2 - \dots$	$(-1)^k C(n+k-1, k) = (-1)^k C(n+k-1, n-1)$
$\frac{1}{(1-ax)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)a^k x^k$ $= 1 + C(n,1)ax + C(n+1,2)a^2x^2 + \dots$	$C(n+k-1, k)a^k = C(n+k-1, n-1)a^k$
$e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$	$1/k!$
$\ln(1+x) = \sum_{k=1}^{\infty} \frac{(-1)^{k+1}}{k} x^k = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots$	$(-1)^{k+1}/k$

# Advanced Counting Techniques

## Generating functions

- In how many different ways can eight identical cookies be distributed among three distinct children if each child receives at least two cookies and no more than four cookies?
- Use generating functions to determine the number of ways to insert tokens worth \$1, \$2, and \$5 into a vending machine to pay for an item that costs  $r$  dollars in both the cases when the order in which the tokens are inserted does not matter and when the order does matter.
- Use generating functions to find the number of  $r$ -combinations from a set with  $n$  elements when repetition of elements is allowed.

End