#### COL106: Data Structures and Algorithms

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#### Data Structures: Balanced Binary Search Trees

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• Consider the following implementation:

#### Code class Node{ public int key; public String value; public Node leftChild; public Node rightChild; public Node parent; public class BST{ public int size; public Node root; public BST(){ size = 0;root = null; public boolean isLeaf(Node N){//To be written} public String get(int k){//To be written} public void put(int k, String v){//To be written} public void remove(int k){//To be written}

- What is the worst case running time of each of the following operations?
  - get(k):
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  - get(k): *O*(*n*)
  - put(k, v): O(n)
  - remove(k): *O*(*n*)

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- What is the worst case running time of each of the following operations when the BST is balanced?
  - get(k):
  - put(k, v):
  - remove(k):
- A BST is perfectly balanced if for every internal node, there are equal number of nodes in its left and right sub-trees.

- What is the worst case running time of each of the following operations when the BST is balanced?
  - get(k):  $O(\log n)$
  - put(k, v):  $O(\log n)$
  - remove(k):  $O(\log n)$
- So, our next goal shall be to build balanced binary search trees.

• Suppose we start with an empty BST and insert the keys 1, 2, 3, 4, then the BST obtained is shown below.



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- This tree is not balanced. Can you think of a way to balance this tree?



#### Data Structures Balanced Binary Search Trees

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• Rotation for tree balancing.



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- <u>AVL Tree</u>: An AVL tree is a binary search tree that satisfies the following property: <u>Height balance property</u>: For every internal node of the tree, the heights of its children differ by at most 1.
- <u>Claim</u>: The height of any AVL tree storing *n* nodes is  $O(\log n)$ .

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Height balance property: For every internal node of the tree, the heights of its children differ by at most 1.

- <u>Claim</u>: The height of any AVL tree storing n nodes is  $O(\log n)$ .
  - Let n(h) denote the minimum number of internal nodes in an AVL tree with height h.
  - Try writing a recurrence relation for *n*(*h*) and solving it to get a lower bound.

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- Question: How do we perform put(k, v) operation on an AVL tree?
  - Same as in BST. However, you also have to make sure that after insertion, the height balance property is maintained.
  - Consider inserting an entry with key 32 in the Tree below.



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  - Consider inserting an entry with key 32 in the Tree below.



• Question: How do we perform put(k, v) operation on an AVL tree?

#### Algorithm

//p denotes the node that is inserted.

BalanceAfterPut(Node p)

- While going up from *p*, let *z* denote the first node for which the height balance property is not satisfied.
- Let y be the child of z with greater height.
- Let x be the child of y with greater height.
- Perform a tri-node restructuring w.r.t. nodes x, y, z.

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