# COL106: Data Structures and Algorithms

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- Note that a Max-Heap can be defined in a similar manner as a Min-Heap.

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- Pointer manipulations can sometimes be intricate (recall the incrementLastNode methods in the homework).
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  - Array elements are typically stored in contiguous locations in the memory and this has its advantages.

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  - <u>Claim 1</u>: Nodes of level *i* are labeled  $2^i, ..., 2^{i+1} 1$ .



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- Main idea: Store node labeled *i* at index *i* of the array.



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  - Which array location should we put this entry in order to maintain a complete binary tree with (*size* + 1) nodes?



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• Consider an array based implementation of min-heap.

Implementation
class Entry{
public int key;
public String value;
<pre>public Entry(int k, String v){</pre>
key = k;
value = v;
}
}
public class MinHeapArray{
final int MAX_HEAP_SIZE = 1000;
public Entry[] A;
public int size;
public MinHeapArray(){
size $= 0;$
$A = new Entry[MAX_HEAP_SIZE];$
}
<pre>public void upHeapBubble(int i){//To be written}</pre>
<pre>public void downHeapBubble(int i){//To be written}</pre>
<pre>public String min(){//To be written}</pre>
<pre>public void insert(int k, String v){//To be written}</pre>
<pre>public String removeMin(){//To be written}</pre>
}

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- <u>Method 1</u>: Perform *n* insert operations in  $O(n \log n)$  time.
- <u>Method 2</u>: Bottom-up heap construction
  - <u>Question</u>: Suppose we have a min-heap H<sub>1</sub> and H<sub>2</sub> both containing 2<sup>h</sup> 1 entries and an entry E. Can you construct a min-heap for all entries in H<sub>1</sub>, H<sub>2</sub> and E combined? What is the running time for your combination algorithm?

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  - Suppose this construction is performed on an array with  $n = 2^{h+1} 1$  entries. What is the running time?

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