# COL351: Analysis and Design of Algorithms 

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## Applications of Network Flow

# Network Flow 

- Suppose there are four teams in IPL with their current number of wins:
- Daredevils: 10
- Sunrisers: 10
- Lions: 10
- Supergiants: 8
- There are 7 more games to be played. These are as follows:
- Supergiants plays all other 3 teams.
- Daredevils Vs Sunrisers, Sunrisers Vs Lions, Daredevils Vs Lions, Sunrisers Vs Daredevils


## Network Flow <br> Team Elimination

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- There are 7 more games to be played. These are as follows:
- Supergiants plays all other 3 teams.
- Daredevils Vs Sunrisers, Sunrisers Vs Lions, Daredevils Vs Lions, Sunrisers Vs Daredevils
- A team is said to be eliminated if it cannot end with maximum number of wins.
- Can we say that Supergiants have been eliminated give the current scenario?


# Network Flow <br> Team Elimination 

- Suppose there are four teams in IPL with their current number of wins:
- Daredevils: 10
- Sunrisers: 10
- Lions: 9
- Supergiants: 8
- There are 7 more games to be played. These are as follows:
- Supergiants plays all other 3 teams.
- 4 games between Daredevils and Sunrisers.
- Can we say that Supergiants have been eliminated give the current scenario?


## Network Flow

Team Elimination

## Problem

There are $n$ teams. Each team $i$ has a current number of wins denoted by $w(i)$. There are $G(i, j)$ games yet to be played between team $i$ and $j$. Design an algorithm to determine whether a given team $x$ has been eliminated.

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- Consider the following flow network


Figure: Team $x$ can end with at most $m$ wins, i.e., $m=w(x)+\sum_{j} G(x, j)$

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## Team Elimination

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- Claim 1: Team $x$ has been eliminated iff the maximum flow in the network is $<g^{*}$, where $g^{*}=\sum_{i, j \text { s.t. } x \notin\{i, j\}} G(i, j)$.


Figure: Team $x$ can end with at most $m$ wins, i.e., $m=w(x)+\sum_{j} G(x, j)$

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- Comment: If we can somehow find a subset $T$ of teams (not including $x$ ) such that
$\sum_{i \in T} w(i)+\sum_{i<j \text { and } i, j \in T} G(i, j)>m \cdot|T|$. Then we have a witness to the fact that $x$ has been eliminated.



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- Can we find such a subset $T$ ?



## Network Flow

## Team Elimination

- Claim 1: Team $x$ has been eliminated iff the maximum flow in the network is $<g^{*}$, where $g^{*}=\sum_{i, j \text { s.t. } x \notin\{i, j\}} G(i, j)$.


## Proof.

- Claim 1.1: If $x$ has been eliminated, then the max flow in the network is $<g^{*}$.



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## Team Elimination

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## Proof of Claim 1

- Claim 1.1: If $x$ has been eliminated, then the max flow in the network is $<g^{*}$.
- Claim 1.2: If the max flow is $<g^{*}$, then team $x$ has been eliminated.


## Proof of Claim 1.2

- Consider any s-t min-cut $(A, B)$ in the graph.
- Claim 1.2.1: If $v_{i j}$ is in $A$, then both $v_{i}$ and $v_{j}$ are in $A$.



## Network Flow

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Team Elimination

- Claim 1: Team $x$ has been eliminated iff the maximum flow in the network is $<g^{*}$, where $g^{*}=\sum_{i, j \text { s.t. } x \notin\{i, j\}} G(i, j)$.


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- Claim 1.1: If $x$ has been eliminated, then the max flow in the network is $<g^{*}$.
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## Proof of Claim 1.2

- Consider any $s$ - $t$ min-cut $(A, B)$ in the graph.
- Claim 1.2.1: If $v_{i j}$ is in $A$, then both $v_{i}$ and $v_{j}$ are in $A$.
- Claim 1.2.2: If both $v_{i}$ and $v_{j}$ are in $A$, then $v_{i j}$ is in $A$.
- Let $T$ be the set of teams such that $i \in T$ iff $v_{i} \in A$.



## Network Flow

Team Elimination

- Claim 1: Team $x$ has been eliminated iff the maximum flow in the network is $<g^{*}$, where $g^{*}=\sum_{i, j \text { s.t. } \times \notin\{i, j\}} G(i, j)$.


## Proof of Claim 1

- Claim 1.1: If $x$ has been eliminated, then the max flow in the network is $<\mathrm{g}^{*}$.
- Claim 1.2: If the max flow is $<g^{*}$, then team $x$ has been eliminated.


## Proof of Claim 1.2

- Consider any s-t min-cut $(A, B)$ in the graph.
- Claim 1.2.1: If $v_{i j}$ is in $A$, then both $v_{i}$ and $v_{j}$ are in $A$.
- Claim 1.2.2: If both $v_{i}$ and $v_{j}$ are in $A$, then $v_{i j}$ is in $A$.
- Let $T$ be the set of teams such that $i \in T$ iff $v_{i} \in A$. Then we have:

$$
\begin{aligned}
& C(A, B)=\sum_{i \in T}(m-w(i))+\sum_{\{i, j\} \not \subset T} G(i, j)<g^{*} \\
\Rightarrow \quad & m \cdot|T|-\sum_{i \in T} w(i)+\left(g^{*}-\sum_{\{i, j\} \subset T} G(i, j)\right)<g^{*} \\
\Rightarrow \quad & \sum_{i \in T} w(i)+\sum_{\{i, j\} \subset T} G(i, j)>m \cdot|T|
\end{aligned}
$$

## End

