

# Introduction to the Fundamentals of Blockchain

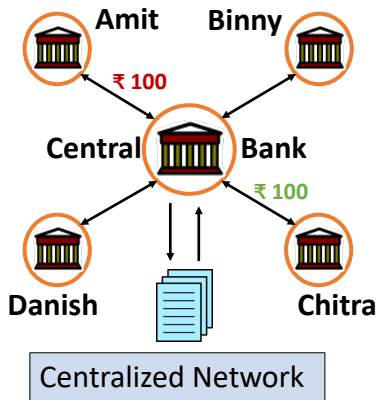
Subodh Sharma

Subhashis Banerjee



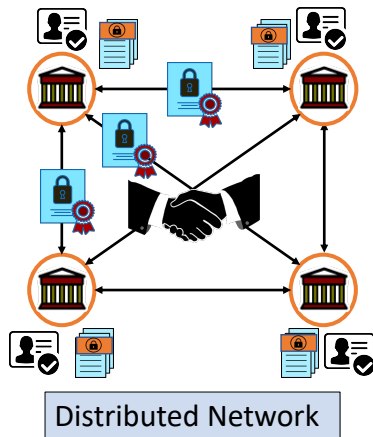
IIT Delhi, Computer Science Department

# What is Blockchain?



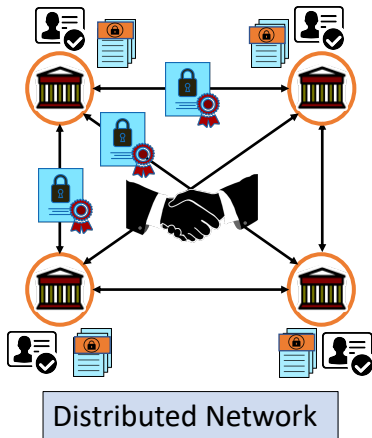
- ▶ Central authority oversee the transaction; implicit trust placed on the overseer.
- ▶ Maintain a ledger of records, balance information per account etc.

# What is Blockchain?



- ▶ No central authority; register of transactions and other meta-information is *replicated* and *distributed*.
- ▶ Every transaction and account information is visible to everyone
- ▶ The records are *untamperable*
- ▶ How do we agree on the transaction or the order on them?
- ▶ How do we stop double-spending?
- ▶ What problems in digital systems can blockchain solve?

# What is Blockchain?



- ▶ Blockchain is a data-structure that contains an ordered sequence of transaction records and other meta-information
- ▶ Each participant in the network can have a fully copy of the blockchain
- ▶ The records are chained via *hash pointers*
- ▶ All updates to Blockchain via *distributed consensus*

# Properties satisfied by Blockchain?



- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place

# Properties satisfied by Blockchain?



- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place
- ▶ Every participant having visibility of the entire global state of the chain lends *transparency* and *local verifiability*

# Properties satisfied by Blockchain?



- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place
- ▶ Every participant having visibility of the entire global state of the chain lends *transparency* and *local verifiability*
- ▶ Hash chains (through hash pointers) provides *data integrity*

# Properties satisfied by Blockchain?



- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place
- ▶ Every participant having visibility of the entire global state of the chain lends *transparency* and *local verifiability*
- ▶ Hash chains (through hash pointers) provides *data integrity*
- ▶ Hash chains also provide chronology of data's existence. This gives *traceability*



# Properties satisfied by Blockchain?



- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place
- ▶ Every participant having visibility of the entire global state of the chain lends *transparency* and *local verifiability*
- ▶ Hash chains (through hash pointers) provides *data integrity*
- ▶ Hash chains also provide chronology of data's existence. This gives *traceability*
- ▶ Digitally signed transactions provide *accountability*

# Properties satisfied by Blockchain?



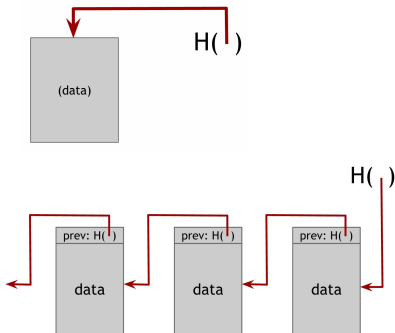
- ▶ Replicated storage of the chain makes data *available* even when n/w faults take place
- ▶ Every participant having visibility of the entire global state of the chain lends *transparency* and *local verifiability*
- ▶ Hash chains (through hash pointers) provides *data integrity*
- ▶ Hash chains also provide chronology of data's existence. This gives *traceability*
- ▶ Digitally signed transactions provide *accountability*
- ▶ *Distributed consensus* provides trust and, consequently, reliability



Cryptographic Hash Functions (CHF) have three properties:

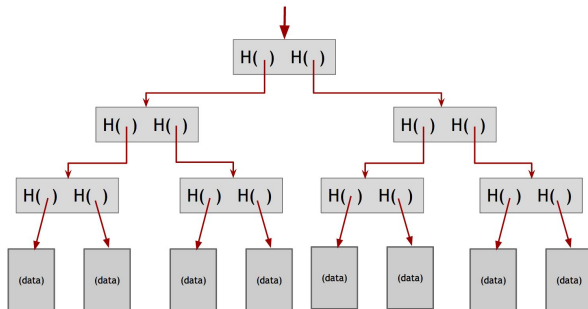
- ▶ Collision-resistance: Infeasible to find two values  $x$  and  $y$  s.t.  $x \neq y$ , yet  $H(x) = H(y)$ .
- ▶ Hiding: Given  $y = H(x)$ , there is no feasible way to figure out the value of  $x$ . If  $x$  is not drawn from a domain that is *spread out*, then choose a secret value  $r$  from a probability distribution that has *high min-entropy* s.t. the hiding property holds for  $H(r||x)$ .
  - ▶ Application in creating **commitment**.
- ▶ Puzzle-friendliness : If for every possible  $n$ -bit output value  $y$ , if  $k$  is chosen from a distribution with high min-entropy, then it is infeasible to find  $x$  such that  $H(k||x) = y$  in time significantly less than  $2^n$ .
  - ▶ Application in **search puzzles**. Given a nonce  $n$ , the hash function  $H$ , and the output target set  $Y$  find  $x$  s.t.  $H(n||x) \in Y$

# Components in Blockchain: Hash Chains



- ▶ HashPtr: Simply a pointer to where the information is stored together with the hash of the information.
- ▶ HashChain: List of HashPtrs. As long as the head ptr is stored securely (i.e. an adversary can't access it), we will have a **tamper-evident** log.

# Components in Blockchain: Merkle Trees



- ▶ Hashptrs organised in a binary tree
- ▶ Property: provides concise proof of *membership*

# Components in Blockchain: Digital Signatures



- ▶ Properties: Valid signatures must verify and infeasible to forge signatures
- ▶ Public keys as digital identities: decentralized identity management
- ▶ Consequence: one can make many identities



## The consensus Problem

- ▶ **Agreement:** All honest processes must agree on the *same* value
- ▶ **Validity:** If all the honest processes have the same initial value, then the agreed upon value must all be that same value
- ▶ **Termination:** Every honest process must *eventually* decide on a value.

# Results from Distributed Computing



Failure Mode	Synchronous System	Asynchronous system
No failure	Agreement	Agreement
Crash failure	Agreement $f < n$	No agreement
Byzantine failure	Agreement $f \leq n/3$	No agreement

**Table:** Results on Agreement.  $n$  is the total number of processes and  $f$  is the number of failure-prone processes.

- ▶ Impossibility of distributed consensus with one faulty processor. M. Fischer, N. Lynch, M. Paterson. *Journal of ACM*, 1985.
- ▶ Consensus in the presence of partial synchrony. C. Dwork and N. Lynch. *Journal of ACM*, 1988.





- ▶ Rely on **hash puzzle**.
- ▶ The node proposing a block is required to find a number, or nonce, s.t.  $H(\text{nonce} || \text{prev}_{hash} || tx_1 \cdots tx_n) < \text{tgt}$
- ▶ Target space is quite small in comparison to the output space of the hash function  $H$ .
- ▶ Fixed protocol to assign the target space

Many other consensus protocols: PBFT, Proof-of-stake, Algorand (cryptographic sortition), Hashgraph, etc.

# Putting it all together



## Simplified protocol

- ▶ New transaction are bcast to all the nodes; each node selects and collects transactions into a block.
- ▶ In each round a random node (vsi proof-of-XYZ) is chosen who gets to bcast its block.
- ▶ Other nodes decide on the block (accept: if all transactions in it are valid)
- ▶ Implicit acceptance: node express acceptance by attaching the block in their local copies of the chain and including the hash of the accepted block in the next block they propose.

# Addressing Denial of Service Attack



- ▶ Say  $A$  dislikes  $B$  and decides to not include any transaction from  $B$  in any block that she proposes.
- ▶  $B$ 's transactions may genuinely not get included in a block in a round where  $A$  is proposing.
- ▶ However due to random node selection,  $B$ 's transactions will get eventually added to a block.

# Can Blockchain solve the privacy problem?



- ▶ While blockchains can support data minimisation, can they support purpose limitation?
- ▶ How is distributed yet regulated access control will be implemented in Blockchain?
- ▶ While distributed consensus may ensure safety, can it guarantee no private information leaks through insider attacks?
- ▶ How are private keys secured from privileged software?



- ▶ Bitcoin and Cryptocurrency Technologies. Arvind Narayanan, Joseph Bonneau, Edward Felten, Andrew Miller, Steven Goldfeder.