

Computational Geometry Lecture 21

Topic: Randomized Incremental
Construction :- contd.
(RIC)

Quicksort:

We visualised the choice of
pivots as RIC

→ Main tool: Backward analysis

Useful to analyse the probability
that an arbitrary element
is involved in the partition
involving the i^{th} pivot.

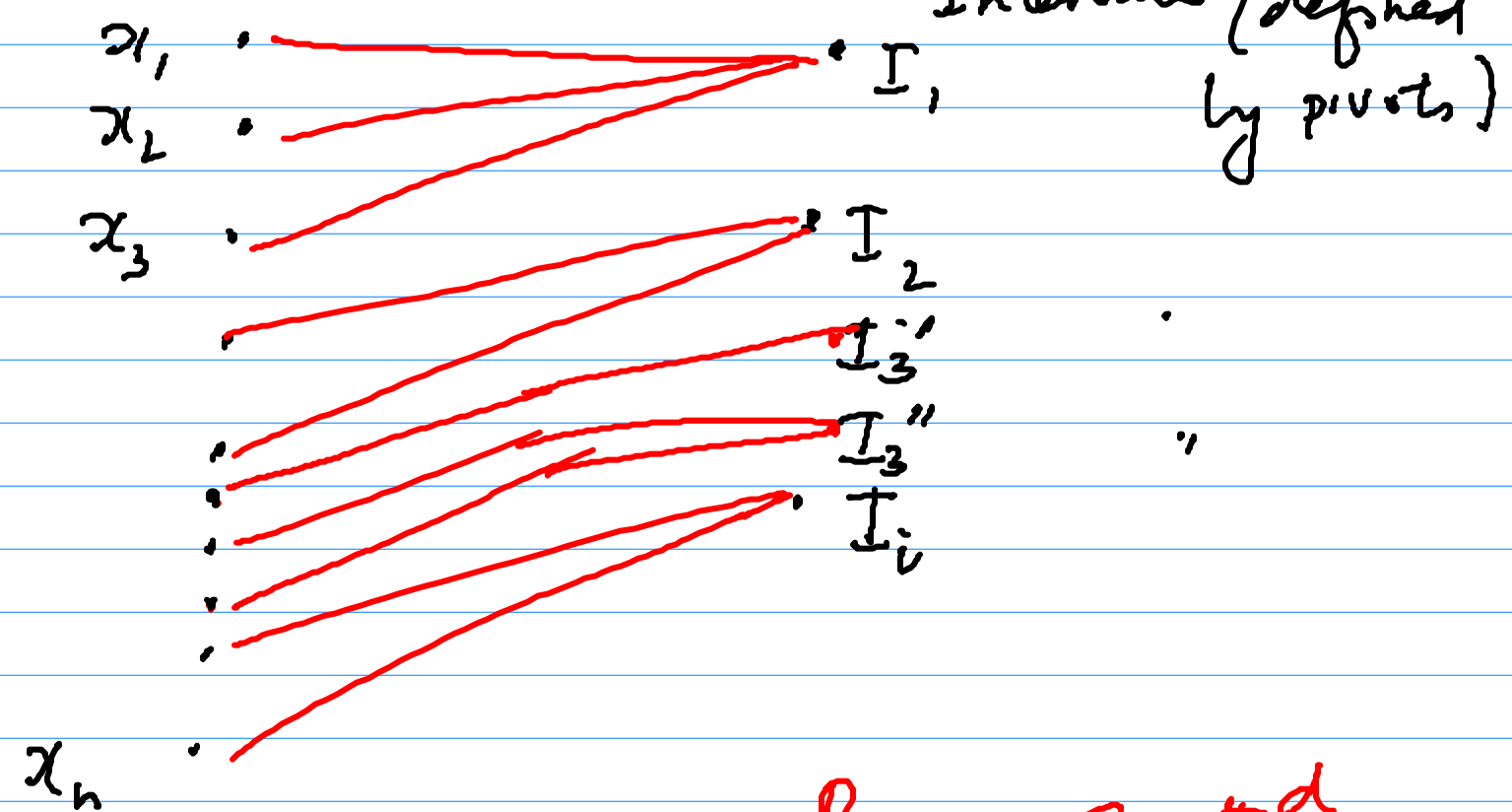
$$\sim \frac{2}{i}$$

Total cost of Quicksort per element
 $= \sum_i \frac{2}{i} \sim O(\log n)$ (Expected)

⇒ $O(n \log n)$ expected for all elements.

When the i^{th} pivot is chosen

n elements



Objects

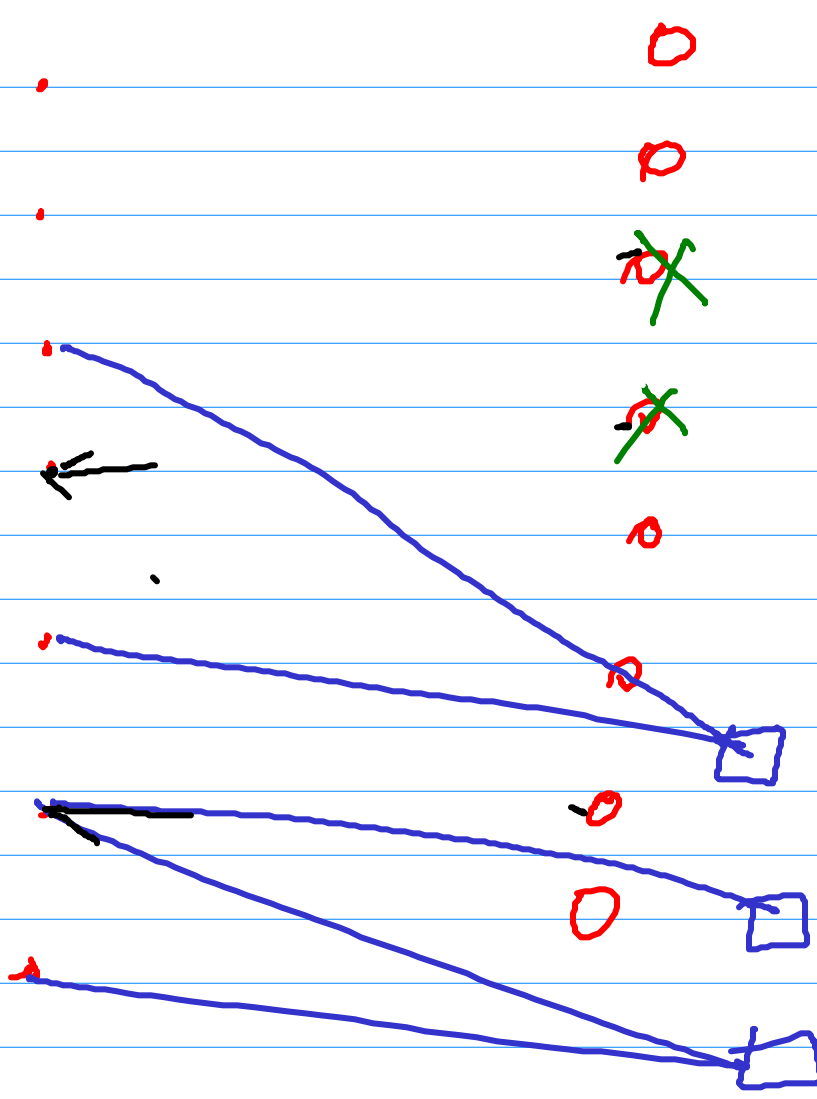
Ranges $\subset \mathbb{F}^d$

↓
defined by the objects

- Kind of ranges: Δ , trapezoids, disks,
- How many objects define the ranges

In the general analysis, as we insert the i^{th} object -

- Some existing ranges have to be deleted
- " new ranges " created
- Objects have to be redistributed



Some edges will be deleted
namely, the ones associated with the ranges being deleted

Some new edges will be created because of new ranges

Structural change in the i^{th} step
 $\# \text{ edges deleted} + \# \text{ edges created}$
 Total str changes $\sum_{i^{\text{th}} \text{ insert.}} (\# \text{ edges deleted} + \text{created})$

Assuming that the algorithmic cost is proportional to the structural changes of the "conflict-graph"
 Total Expected cost = $\sum_i \text{Exp no. of edges created in } i^{\text{th}} \text{ step}$