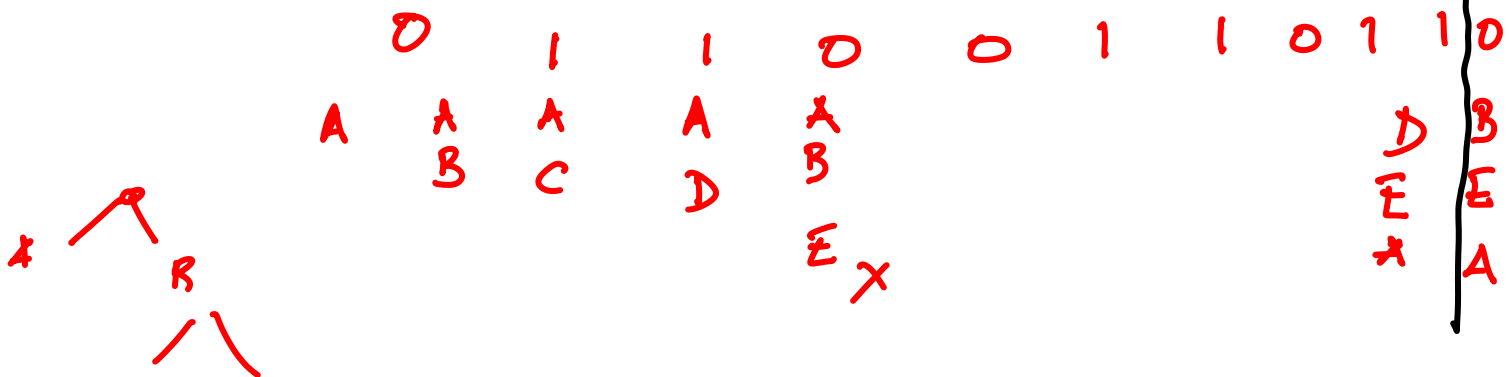


From state A, on input 0, the machine arbitrarily "chooses" one of the possible transitions i.e. state A or state B

Input :



A string w is accepted iff one of the possible terminating states is a Final state

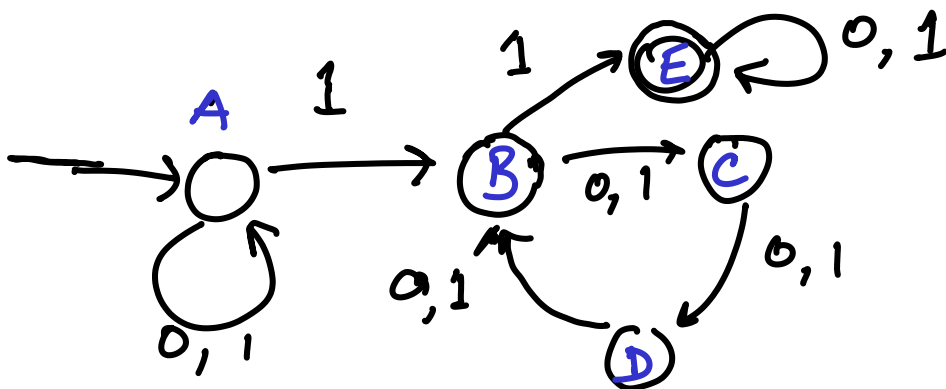
The above automaton accepts
 $w \in \Sigma^*$, iff w ends with 011

This automaton that can have multiple transitions is known as
 Non deterministic Finite Automaton
 (NFA)

Design an NFA that accepts the
 language
 $w \in \Sigma^*$, such that some pair of
 1's is separated by $3i$ symbols
 $i \geq 0$

0 0 1 1 0 1 0 0 1 1 ✓
 ↑ ↑

0 1 0 1 0 0 ✗



NFA N :

Q : set of states

$F \subseteq Q$ set of final states

$q_0 \in Q$: initial state

$\delta: Q \times \Sigma \rightarrow 2^Q - \emptyset$ (multiple successors, states)

$$L(N) = \{ \omega \mid \hat{\delta}(q_0, \omega) \cap F \neq \emptyset \}$$

$$\hat{\delta}(\{q\}, a) = \delta(q, a) \quad \begin{array}{l} q \in Q \\ a \in \Sigma \end{array}$$

$$\hat{\delta}(P, a) = \bigcup_{q \in P} \hat{\delta}(q, a)$$

$P \subseteq Q$

$$\hat{\delta}(P, a\omega) = \hat{\delta}(\hat{\delta}(P, a), \omega)$$

$a \in \Sigma, \omega \in \Sigma^*$