

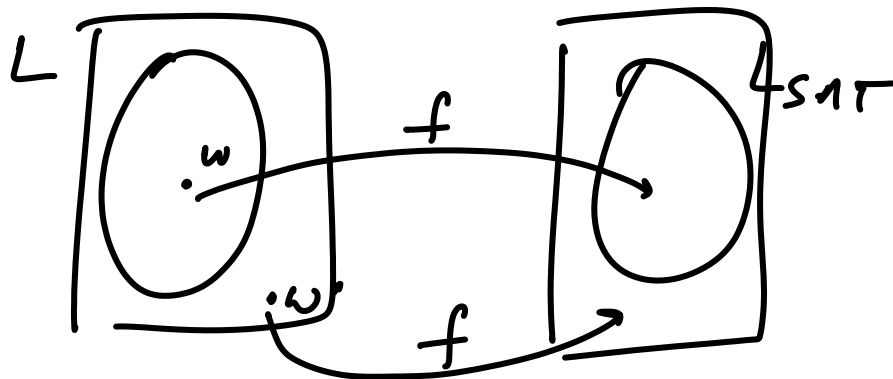
Cook-Levin Theorem : The satisfiability problem of boolean expression is NP-complete

Part 1 : $L_{SAT} \in NP$: easy

Part 2 : for any $L \in NP$ $L \leq_{poly} L_{SAT}$

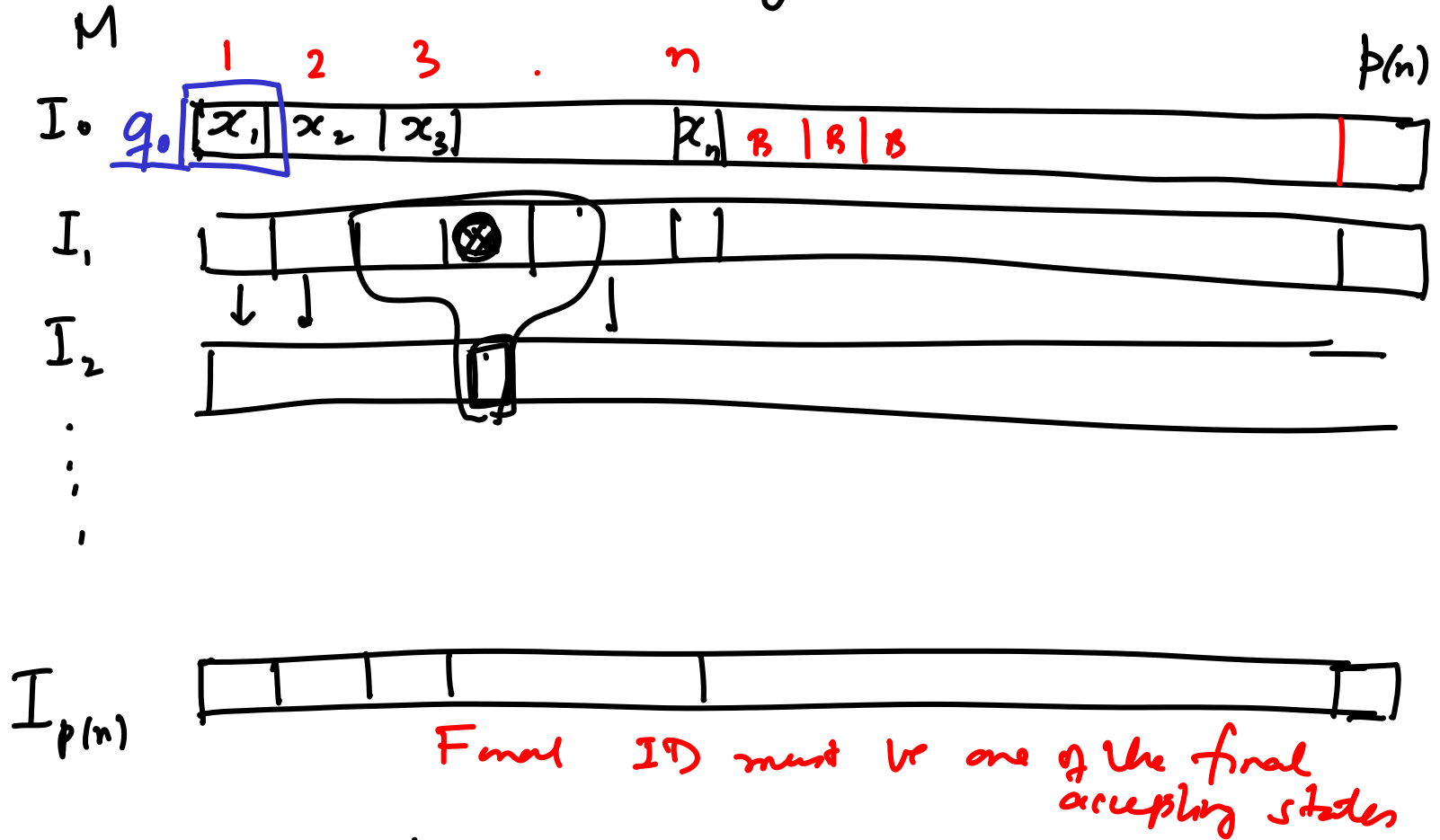
$\Rightarrow \exists$ a Nondet M s.t. $L(M) = L$

Given any w , we have to construct a Turing computable polynomial-time function f s.t. $f(w) \in L_{SAT}$ iff $w \in L$



Note : $f(w)$ is a boolean expression (not too long)

N D T M M takes at most $p(n)$ steps
 for any input of length n where
 p is some polynomial.



δ_M is a transition function that may have multiple successors (non-det)

$w \in L$ iff

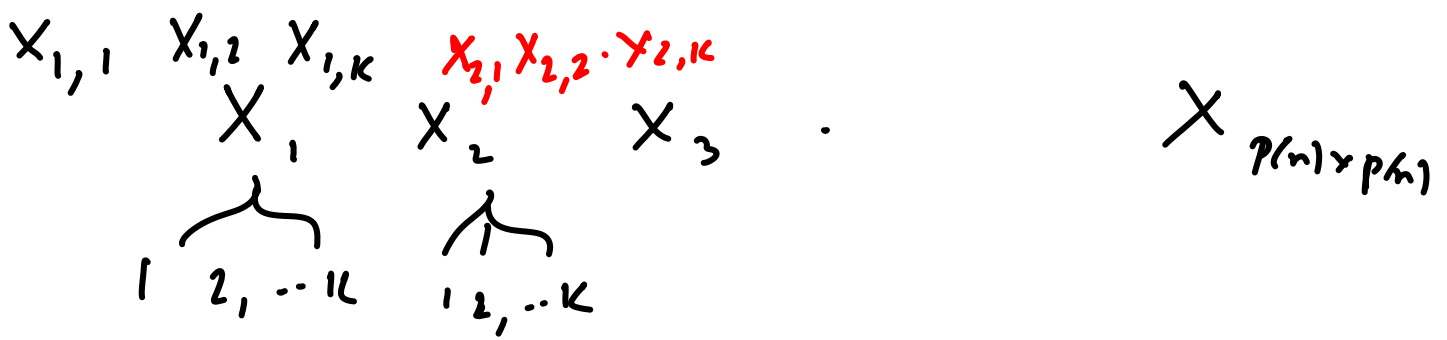
- (1) I_0 contain w as initial input
and
(2) $I_{p(n)}$ contains a final state
and
(3) $\forall j \leq p(n) I_{j+1}$ follows from I_j using a legal transition function of M .

The goal is to represent the above conditions as a boolean formula $F_M(w)$ which is satisfiable iff $w \in L$

redu function f
moreover f is computable in polytime
 $\Rightarrow |F_M(w)|$ is of polynomial length

Idea: To guess each of the $\{1, 2, \dots, k\}$ $p(n) \times p(n)$ symbols and verify conditions 1, 2, 3.

To convert from a k -valued variable to a boolean variable, we can introduce k variables for each variable



$X_{i,j}$ are boolean variables st.

$$X_{i,j} = \begin{cases} \text{true} & \text{if } X_i = j \in \{1, \dots, k\} \\ \text{false} & \text{otherwise} \end{cases}$$

Additional condition (beyond 1, 2, 3)

(4) Exactly one $X_{i,j}$ must be true for each $i \leq p(n) \times p(n)$

$$\left(X_{i,1} \vee X_{i,2} \vee \dots \vee X_{i,k} \right) \wedge_{j \neq j'} \underbrace{\left(X_{i,j} \Rightarrow \overline{X_{i,j'}} \right)}_{\left(\overline{X_{i,j}} \vee \overline{X_{i,j'}} \right)}$$

Total boolean variables : $p(n) \times p(n) \times k$