## CSL 356: Analysis and Design of Algorithms

Instructor: Ragesh Jaiswal

## 1. Another algorithm for max-flow

Consider the following slightly changed version of the Ford-Fulkerson max-flow algorithm. This algorithm is also due to Jack Edmonds and Richard Karp.

## Max-Flow

- Start with a flow $f$ such that $\forall e \in E, f(e)=0$.
- While there is an $s-t$ path in $G_{f}$
- Find an $s-t$ path $P$ in $G_{f}$ with largest bottleneck value.
- Augment along $P$ to obtain $f^{\prime}$.
- Update $f$ to $f^{\prime}$ and $G_{f}$ to $G_{f^{\prime}}$
return(f).
(a) Think of an algorithm to find the largest bottleneck path from $s$ to $t$ in a given graph. A bottleneck path is a path such that the bottleneck edge has maximum weight. Discuss its running time.
(Hint: Try ideas from Dijkstra's Algorithm.)
(b) Let $f$ be any $s-t$ flow and $t$ be the value of maximum flow in the residual graph $G_{f}$. Let $f^{\prime}$ be the new flow after one augmentation and $t^{\prime}$ be the value of the new maximum flow in the residual graph $G_{f^{\prime}}$. Argue that $t^{\prime} \leq(1-1 / m) \cdot t$.
(c) Use the properties you showed above to argue that for a graph with integer capacities, the algorithm runs in time $O\left(m^{2} \cdot \log m \cdot \log f^{*}\right)$, where $f^{*}$ is the value of the max-flow in the original graph $G$.

