# CSL 356: Analysis and Design of Algorithms

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<u>Median Finding</u>: Finding the  $k^{\text{th}}$  smallest number in an unsorted array.

- <u>Problem(Median Finding)</u>: Given an array A of unsorted numbers and an integer k. Give an algorithm that finds the  $k^{th}$  smallest number in the array. Assume A contains distinct numbers.
- Divide and Conquer:
  - Pick an number p as pivot. Partition the numbers in A into  $A_L$  (all numbers < p) and  $A_R$  (all numbers > p).

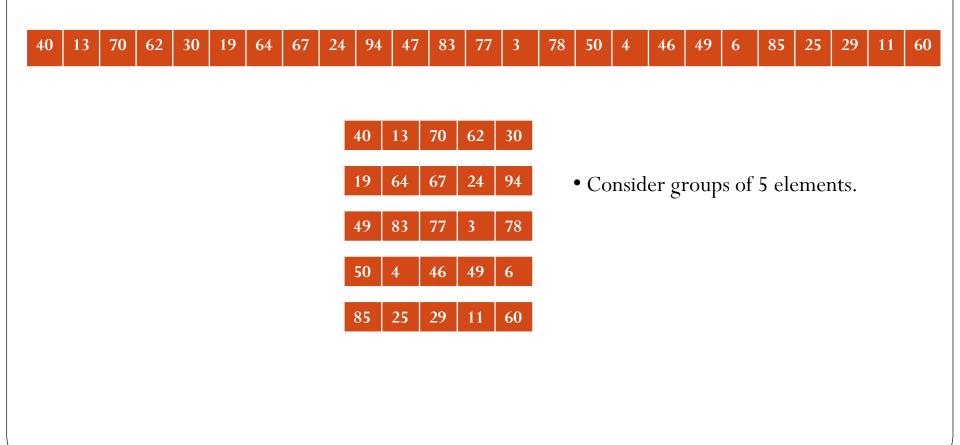
• If 
$$|A_L| = k - 1$$
, then output  $p$ .

- If  $|A_L| > k 1$ , then recursively find the  $k^{th}$  smallest number in  $A_L$
- If  $|A_L| < k 1$ , then recursively find the  $(k |A_L| 1)^{th}$  smallest number in  $A_R$ .

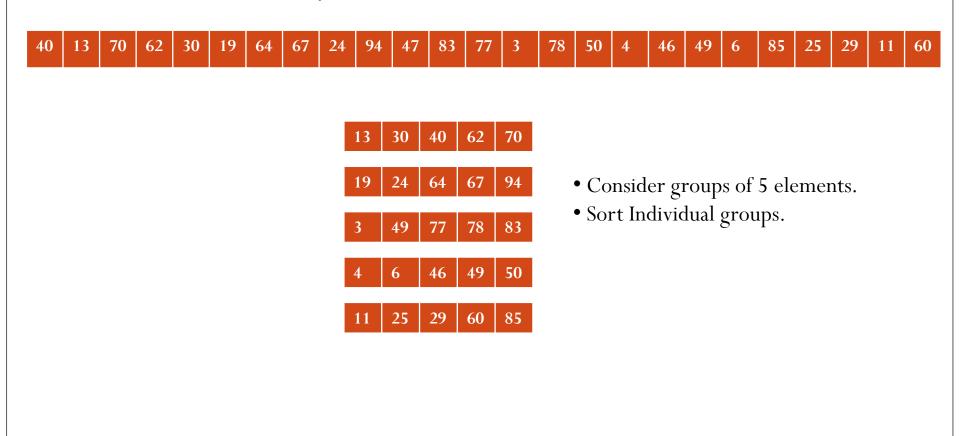
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- What is the running time of this algorithm?
  - If we pick a bad pivot each time, then the running time can be as bad as  $O(n^2)$ .

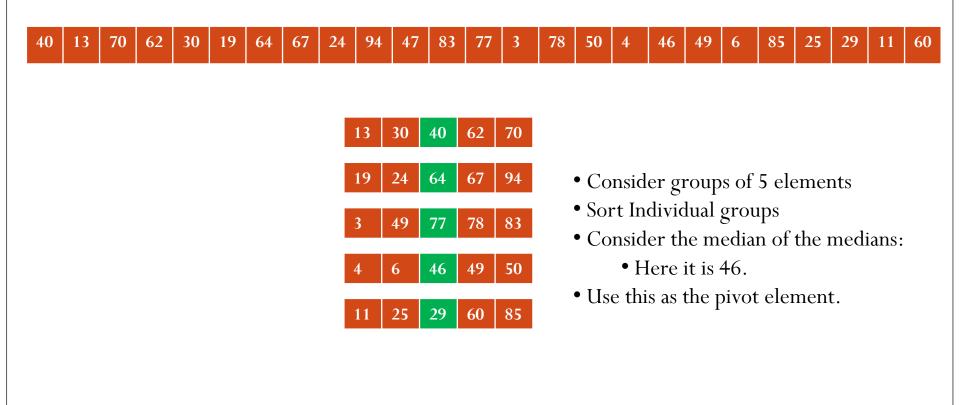
- How do we pick a good pivot number?
  - Randomly: We will look at this a bit later.
  - Deterministically: ?



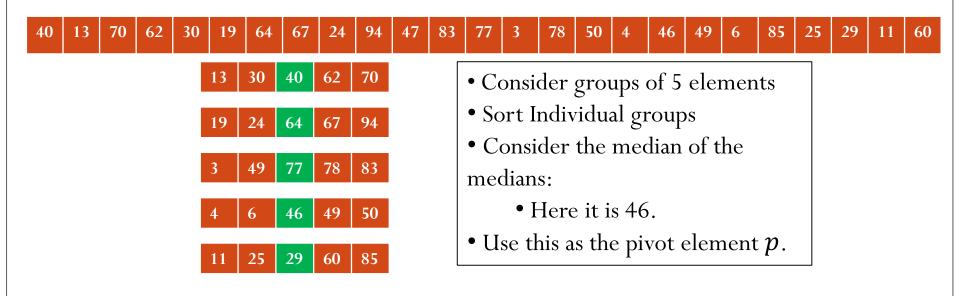
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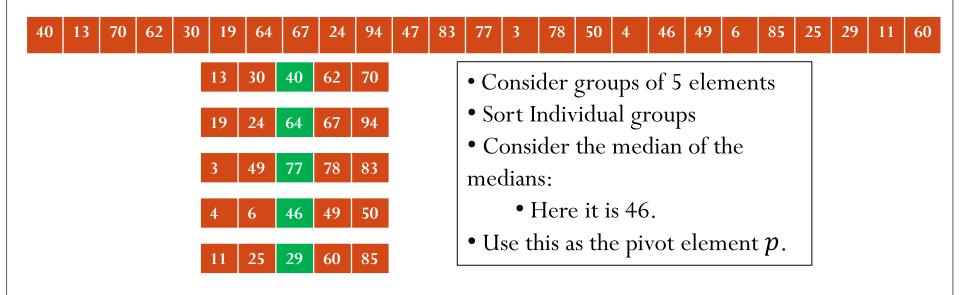


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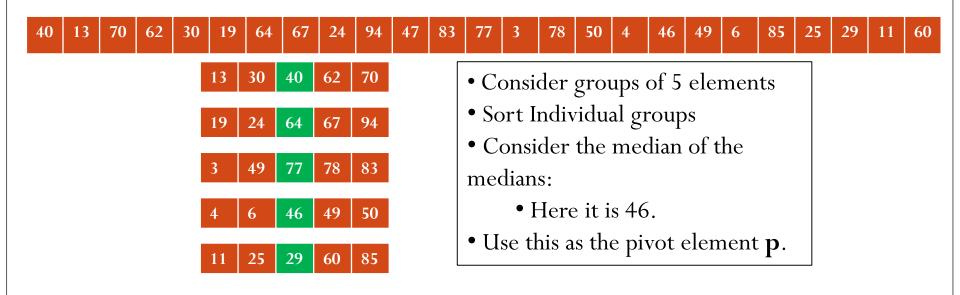
• How many elements in *A* are larger than *p*?

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  - Deterministically: ?



- How many elements in *A* are larger than *p*?
  - <u>Claim</u>: There are at least (3n/10 6) numbers in *A* that are larger than *p*.

- How do we pick a good pivot number?
  - Randomly: We will look at this a bit later.
  - Deterministically: ?



- How many elements in *A* are smaller than *p*?
  - <u>Claim</u>: There are at least (3n/10 6) numbers in A that are smaller than p.

#### Find-kth-smallest(A, k)

- ... / / Base cases
- Consider groups of 5 numbers, sort each group and create another array B containing the median numbers from each group.
- $p \leftarrow \text{Find-kth-smallest}(B, \text{floor}(|B|/2))$
- Partition the array A into  $A_L$  and  $A_R$  using p as the pivot.
- If  $(|A_L| = k 1)$  then output(p)
- If  $(|A_L| > k 1)$  then output(Find-kth-smallest( $A_L$ , k))
- If  $(|A_L| < k 1)$  then output(Find-kth-smallest( $A_R, k |A_L| 1$ ))
- <u>Running time</u>:?

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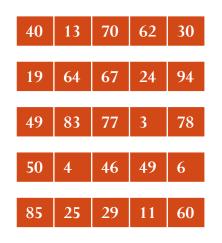
- T(n) = T([n/5]) + T(7n/10 + 6) + O(n)
- T(n) = ?

• Suppose we want to find the 12<sup>th</sup> smallest element in the array?



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• Consider groups of 5 elements.

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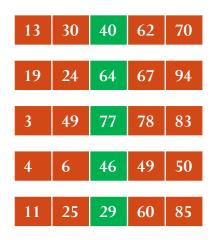


13	30	40	62	70
19	24	64	67	94
3	49	77	78	83
4	6	46	49	50
11	25	29	60	85

- Consider groups of 5 elements.
- Sort individual groups

• Suppose we want to find the 12<sup>th</sup> smallest element in the array?





- Consider groups of 5 elements.
- Sort individual groups.
- Consider medians of each group.

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- Consider groups of 5 elements.
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- Consider medians of each group.
- Make a recursive call to find median element of medians.

29

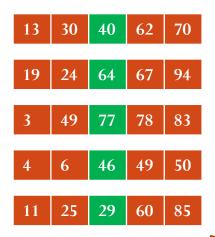
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64

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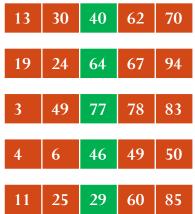




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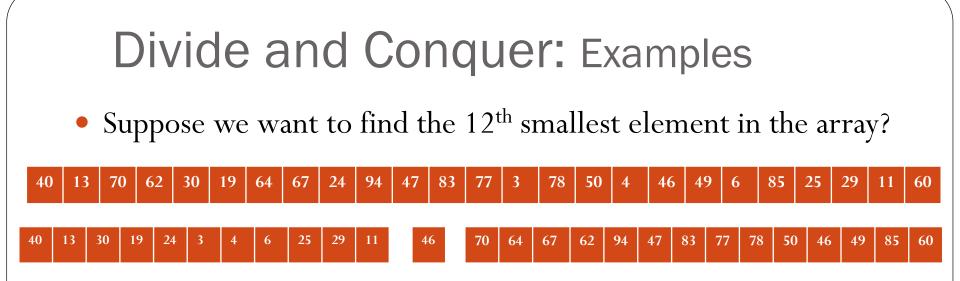
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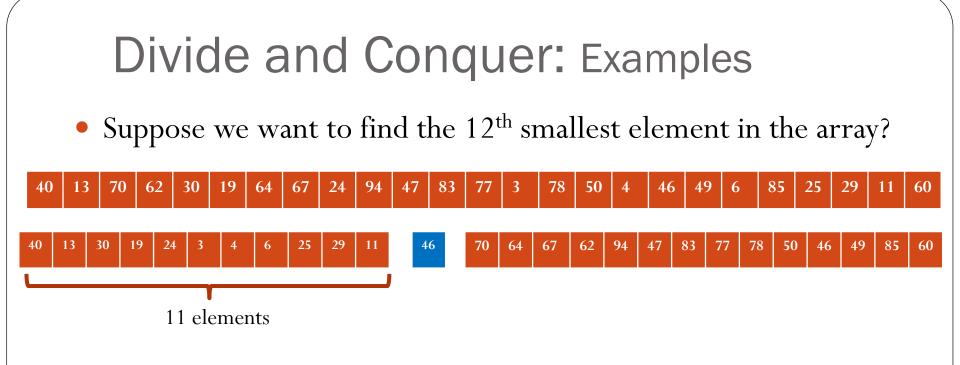
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Fast Fourier Transform (FFT)

- <u>Problem</u>: Given two polynomials  $A(x) = a_0 + a_1 \cdot x + a_2 \cdot x^2 + \dots + a_{n-1} \cdot x^{n-1},$ and  $B(x) = b_0 + b_1 \cdot x + b_2 \cdot x^2 + \dots + b_{n-1} \cdot x^{n-1}$ multiply them.
- We have to obtain the polynomial  $C(x) = A(x) \cdot B(x)$  $C(x) = c_0 + c_1 \cdot x + c_2 \cdot x^2 + \dots + c_{2n-2} * x^{2n-2}$
- What is  $c_i$  in terms of coefficients of A and B?

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- What is  $C_i$  in terms of coefficients of A and B?
  - $c_i = a_i \cdot b_0 + a_{i-1} \cdot b_1 + a_{i-2} \cdot b_2 + \dots + a_0 \cdot b_i$
- The vector  $(c_0, \ldots, c_{2n-2})$  is called the *convolution* of vectors  $(a_0, \ldots, a_{n-1})$  and  $(b_0, \ldots, b_{n-1})$ .

SimpleMultiply(
$$(a_0, ..., a_{n-1}), (b_0, ..., b_{n-1})$$
)  
- For  $i = 0$  to  $2n - 2$   
- For  $j = 0$  to  $i$   
-  $c_i = c_i + a_j \cdot b_{i-j}$   
- return( $(c_0, ..., c_{2n-2})$ )

• What is the running time of the simple algorithm?

•  $O(n^2)$ 

• Is there another way to compute the polynomial C(x)?

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O(n<sup>2</sup>)

• Is there another way to compute the polynomial C(x)?

- Compute  $A(s_1), A(s_2), ..., A(s_{2n})$ .
- Compute  $B(s_1), B(s_2), ..., B(s_{2n})$ .
- Compute

• 
$$C(s_1) = A(s_1) \cdot B(s_1),$$

• 
$$C(s_{2n}) = A(s_{2n}) \cdot B(s_{2n})$$

- Is there another way to compute the polynomial C(x)?
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    - $C(s_1) = A(s_1) \cdot B(s_1)$ ,
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• Interpolate to obtain the polynomial C(x).

• How fast can you compute A(s) for a given value of s?

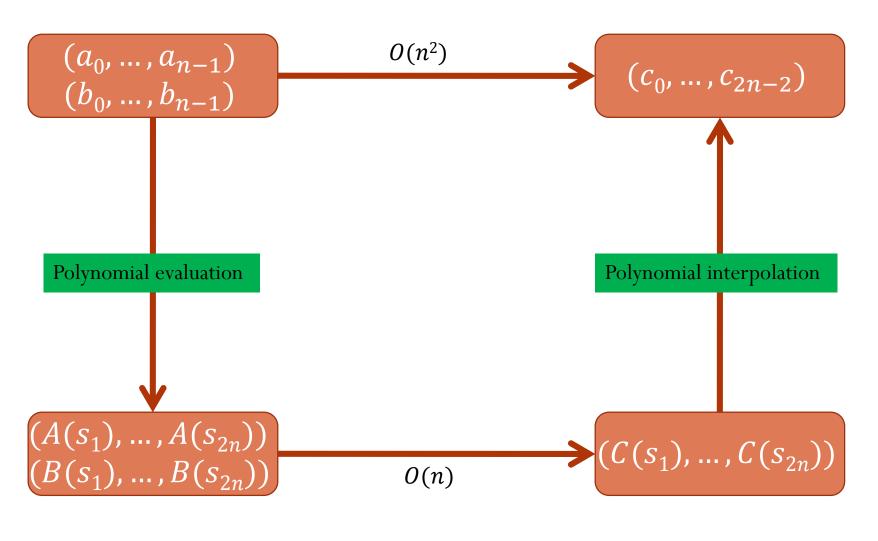
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  - Compute  $A(s_1), A(s_2), ..., A(s_{2n})$ .
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    - $C(s_1) = A(s_1) \cdot B(s_1)$ ,

• ...

•  $C(s_{2n}) = A(s_{2n}) \cdot B(s_{2n}).$ 

• *Interpolate* to obtain the polynomial C(x).

- How fast can you compute A(s) for a given value of s?
  - O(n) arithmetic operations using Horner's rule.
  - $A(s) = a_0 + s \cdot (a_1 + s \cdot (a_2 + \dots + s \cdot (a_{n-2} + s \cdot (a_{n-1})) \dots))$



# End