## CSL 356: Analysis and Design of Algorithms

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## Graphs

## Graphs: Introduction

- A way to represent a set of objects with pair-wise relationships among them.
- The objects are represented as vertices and the relationships are represented as edges.


$$
\begin{gathered}
G=(V, E) \\
V=\left\{v_{1}, \ldots, v_{8}\right\} \\
E=\left\{\left(v_{1}, v_{8}\right), \ldots\right\}
\end{gathered}
$$

## Graphs: Introduction

- Examples:
- Social networks
- Communication networks
- Transportation networks
- Dependency networks



## Graphs: Introduction

- Weighted Graphs: There are weights associated with each edge quantifying the relationship. For example, delay in communication network.



## Graphs: Introduction

- Directed graphs: Asymmetric relationships between the objects. For example, one way streets.



## Graphs: Introduction

- Path: A sequence of vertices $v_{1}, v_{2}, \ldots, v_{k}$ such that for any consecutive pair of vertices $v_{i}, v_{i+1},\left(v_{i}, v_{i+1}\right)$ is an edge in the graph. It is called a path from $v_{1}$ to $v_{k}$. A cycle is a path where $v_{1}=v_{k}$ and $v_{1}, \ldots, v_{k-1}$ are distinct vertices.



## Graphs: Introduction

- Strongly connected: A graph is called strongly connected if for any pair of vertices $u, v$, there is a path from $u$ to $v$ and a path from $v$ to $u$.



## Graphs: Introduction

- Tree: A strongly connected, undirected graph is called a tree if it has no cycles.
- How many edged does a tree have?



## Graph

Data Structures for representing graphs

## Graph: Data structures

- Adjacency matrix:


| $\boldsymbol{v}_{1}$ | $\boldsymbol{v}_{2}$ |  | $\boldsymbol{v}_{3}$ | $\boldsymbol{v}_{4}$ | $\boldsymbol{v}_{5}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
|  | $v_{1}$ | 0 | 1 | 1 | 1 |
| $v_{2}$ | 1 | 0 | 1 | 0 | 0 |
|  | $v_{3}$ | 1 | 1 | 0 | 1 |
| $v_{4}$ | 1 | 0 | 1 | 0 | 0 |
|  | $v_{5}$ | 0 | 0 | 0 | 0 |

- Space: $O\left(n^{2}\right)$


## Graph: Data structures

- Adjacency list: For each vertex store its neighbors

- Space: $O(n+m)$


## Graph

Graph algorithms

## Graph Algorithms: s-t connectivity

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## Graph Algorithms: s-t connectivity

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- There is a path between $S$ and $t$ iff $S$ and $t$ are in the same connected component.
- Alternate problem: What are the vertices which are reachable from $s$. Is $t$ among these reachable vertices.
- Graph exploration: Explore all the vertices reachable from $S$.


## Graph Algorithms: BFS

- Breadth First Search (BFS):

```
BFS}(G,s)
    - Layer(0) = {s}
    - i=1
    - while(true){
        - visit all new nodes that have an edge to a vertex in Layer(i-1)
        - put these nodes in the set Layer(i)
        - if Layer(i) is empty then end
        -i=i+1
    }
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Show: The shortest path from $s$ to any vertex in $L(i)$ is equal to $i$.

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Running time: $O(n+m)$

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- The BFS algorithm defines the following "BFS Tree" rooted at $S$ :
- Vertex $u$ is the parent of vertex $v$ if $u$ caused the immediate discovery of $v$.


## Graph Algorithms: BFS

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- The BFS algorithm defines the following "BFSTree" rooted at $S$ :
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## Graph Algorithms: BFS

- Problem: Given a graph $G=(V, E)$ check if the graph is bipartite.

- A graph is bipartite if the vertices can be partitioned into two sets such that there is no edge between any pair of vertices in the same set.


## Graph Algorithms: BFS

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- Consider BFS:
- Is it possible that there is an edge between vertices which belong to sets $L(i)$ and $L(j)$ such that $(j-i)>1$ ?


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- Can you now use BFS to check if the graph is bipartite?


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- Can you now use BFS to check if the graph is bipartite?
- What is the running time of your algorithm?


## Graph Algorithms: BFS

- Problem: Given a graph $G=(V, E)$ check if the graph is bipartite.
- Suppose the given graph contains a cycle of odd length. Can this graph be bipartite?
- Can you now use BFS to check if the graph is bipartite?
- What is the running time of your algorithm?
- Suppose a graph does not have an odd cycle. Does that mean that the graph is bipartite?


## End

## Problems to think about:

1. The BFS algorithm gives the shortest path from s to any vertex. Suppose we are given a weighted graph where the weights are numbers between 1 and 10. Can you use the BFS algorithm to find the shortest length path from s to any vertex in the graph. The length of a path is the sum of weights of edges in the path.
