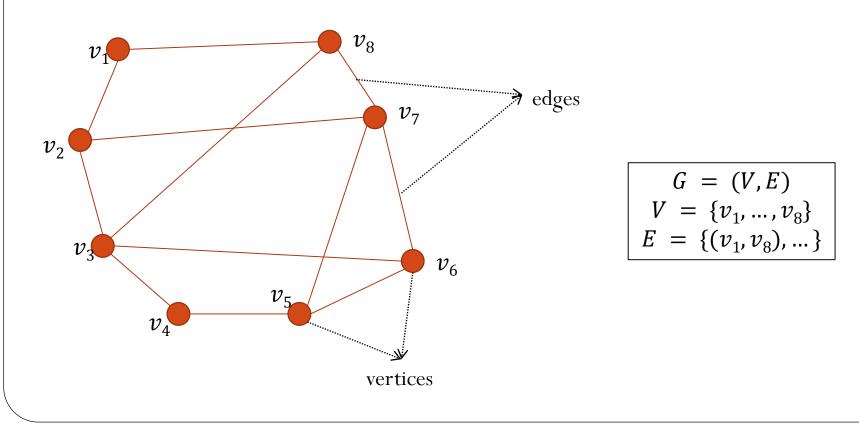
CSL 356: Analysis and Design of Algorithms

Ragesh Jaiswal

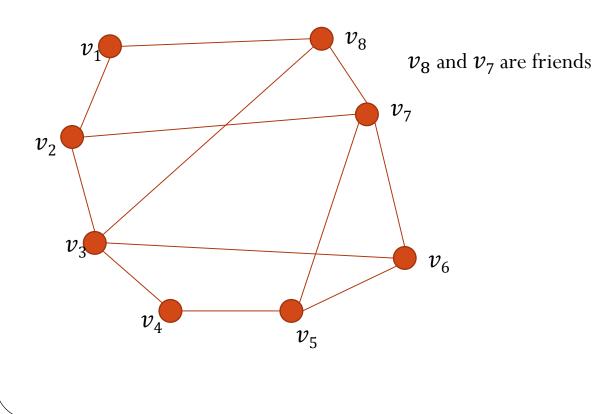
CSE, IIT Delhi

Graphs

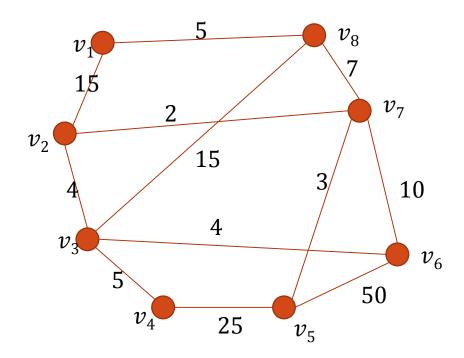
- A way to represent a set of objects with pair-wise relationships among them.
- The objects are represented as vertices and the relationships are represented as edges.



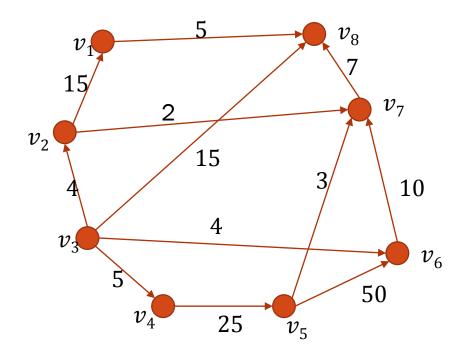
- Examples:
 - Social networks
 - Communication networks
 - Transportation networks
 - Dependency networks



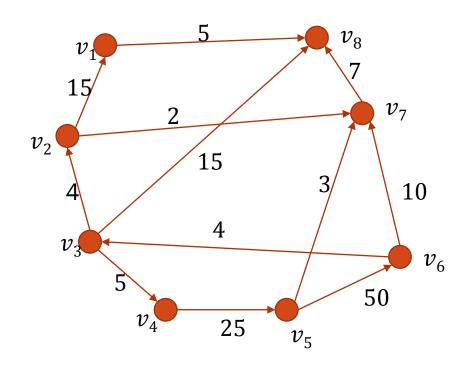
• <u>Weighted Graphs</u>: There are weights associated with each edge quantifying the relationship. For example, delay in communication network.



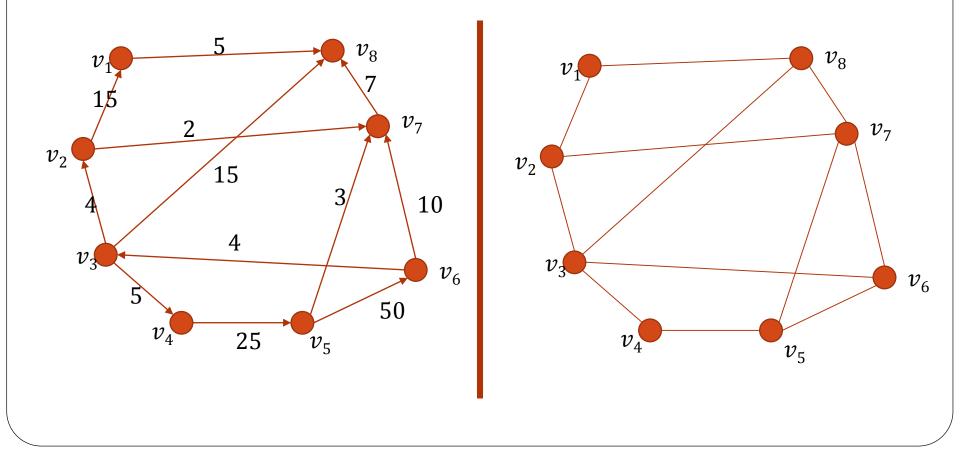
• <u>Directed graphs</u>: Asymmetric relationships between the objects. For example, one way streets.



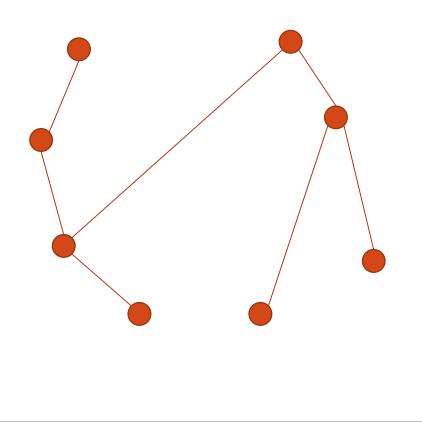
• <u>Path</u>: A sequence of vertices $v_1, v_2, ..., v_k$ such that for any consecutive pair of vertices $v_i, v_{i+1}, (v_i, v_{i+1})$ is an edge in the graph. It is called a path from v_1 to v_k . A cycle is a path where $v_1 = v_k$ and $v_1, ..., v_{k-1}$ are distinct vertices.



• <u>Strongly connected</u>: A graph is called strongly connected if for any pair of vertices u, v, there is a path from u to v and a path from v to u.



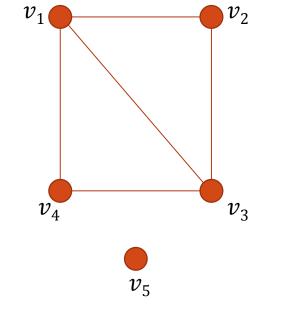
- <u>Tree</u>: A strongly connected, undirected graph is called a tree if it has no cycles.
- How many edged does a tree have?



Graph

Data Structures for representing graphs

Graph: Data structures



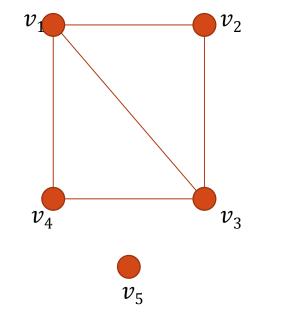
• Adjacency matrix:

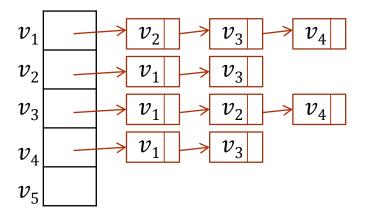
	v_1	v_2	v_3	v_4	v_5
v_1	0	1	1	1	0
v_2	1	0	1	0	0
v_3	1	1	0	1	0
v_4	1	0	1	0	0
v_5	0	0	0	0	0

• Space: $O(n^2)$

Graph: Data structures

• Adjacency list: For each vertex store its neighbors





• Space: O(n + m)

Graph

Graph algorithms

Graph Algorithms: s-t connectivity

• <u>Problem</u>: Given an (undirected) graph G = (V, E) and two vertices S, t, check if there is a path between S and t.

Graph Algorithms: s-t connectivity

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- There is a path between *S* and *t* iff *S* and *t* are in the same connected component.

Graph Algorithms: s-t connectivity

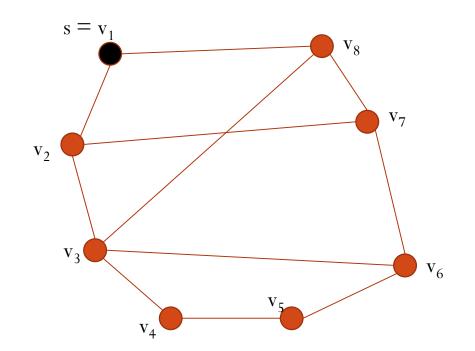
- <u>Problem</u>: Given an (undirected) graph G = (V, E) and two vertices S, t, check if there is a path between S and t.
- There is a path between *S* and *t* iff *S* and *t* are in the same connected component.
- <u>Alternate problem</u>: What are the vertices which are reachable from *S*. Is *t* among these reachable vertices.
 - Graph exploration: Explore all the vertices reachable from *S*.

```
• Breadth First Search (BFS):
```

```
BFS(G, s) \{
- Layer(0) = \{s\}
- i = 1
- while(true) \{
- visit all new nodes that have an edge to a vertex in Layer(i - 1)
- put these nodes in the set Layer(i)
- if Layer(i) is empty then end
- i = i + 1
\}
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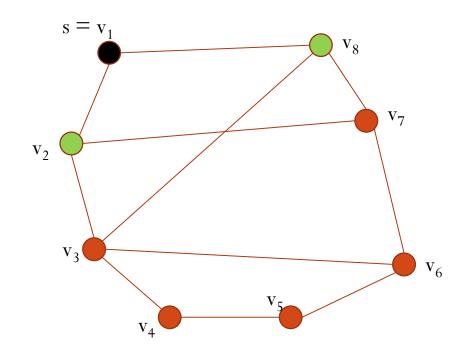
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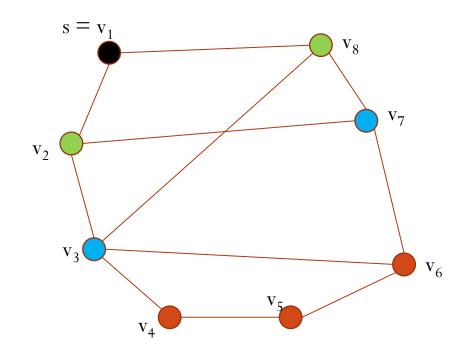
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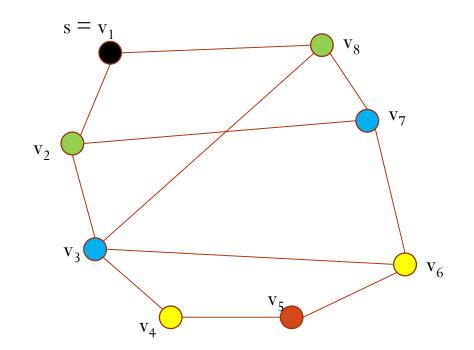
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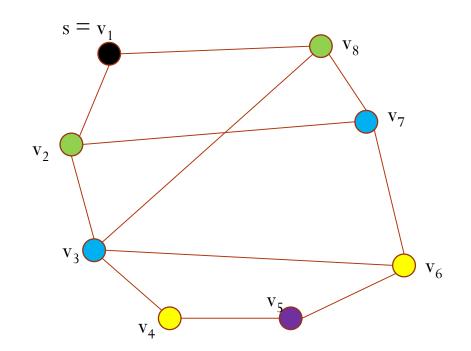
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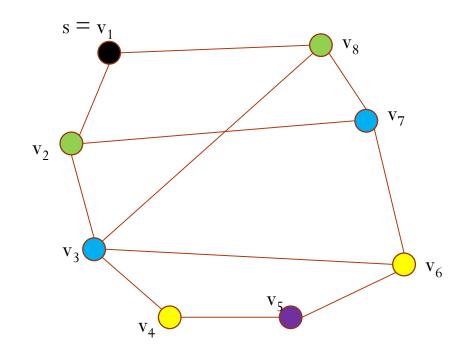
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Show: The shortest path from s to any vertex in L(i) is equal to i.

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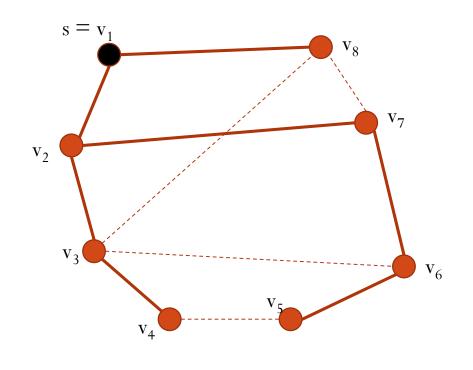
Running time: O(n + m)

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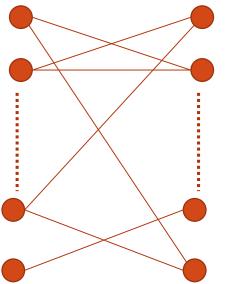
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- The BFS algorithm defines the following "BFS Tree" rooted at *S*:
 - Vertex u is the parent of vertex v if u caused the immediate discovery of v.

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- The BFS algorithm defines the following "BFS Tree" rooted at *S*:
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• Problem: Given a graph G = (V, E) check if the graph is *bipartite*.



• A graph is *bipartite* if the vertices can be partitioned into two sets such that there is no edge between any pair of vertices in the same set.

• Problem: Given a graph G = (V, E) check if the graph is *bipartite*.

BFS(*G*, *s*){

- Layer(
$$0$$
) = {*s*

-i = 1

- while(true) {

- visit all new nodes that have an edge to a vertex in Layer(i-1)

- put these nodes in the set $\mathrm{Layer}(i)$
- if Layer(i) is empty then end

-i = i + 1

• Consider BFS:

Is it possible that there is an edge between vertices which belong to sets L(i) and L(j) such that (j - i) > 1?

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- Suppose the given graph contains a cycle of odd length. Can this graph be bipartite?
- Can you now use BFS to check if the graph is bipartite?
- What is the running time of your algorithm?
- Suppose a graph does not have an odd cycle. Does that mean that the graph is bipartite?

End

Problems to think about:

 The BFS algorithm gives the shortest path from s to any vertex. Suppose we are given a weighted graph where the weights are numbers between 1 and 10. Can you use the BFS algorithm to find the shortest length path from s to any vertex in the graph. The length of a path is the sum of weights of edges in the path.