
Problem Set 2

Due: Monday October 11, 2010, in class.

See course information section (on course web page) for instructions and rules on working on problem sets and turning them in.

Problem 1. [30 points] Let K be a 56-bit DES key and L a 64-bit auxiliary key. For any 64-bit plaintext M let

$$\text{DESY}(K \parallel L, M) = \text{DES}(K, L \oplus M).$$

This defines a family of functions $\text{DESY}: \{0, 1\}^{120} \times \{0, 1\}^{64} \rightarrow \{0, 1\}^{64}$.

- (a) **[8 points]** Show that DESY is a block cipher.
- (b) **[22 points]** Let $(M_1, C_1), (M_2, C_2)$ be input-output examples of DESY under a random 120-bit target key $K \parallel L$. Present an attack that given $(M_1, C_1), (M_2, C_2)$ recovers the target key using at most 2^{57} computations of DES or DES^{-1} . (As usual, the job is actually only to recover a key consistent with the input-output examples, but in practice this is typically the target key.)

Problem 2. [50 points] Let $F: \{0, 1\}^k \times \{0, 1\}^l \rightarrow \{0, 1\}^l$ be a family of functions and let $r \geq 1$ be an integer. The r -round Feistel cipher associated to F is the family of functions $F^{(r)}: \{0, 1\}^k \times \{0, 1\}^{2l} \rightarrow \{0, 1\}^{2l}$, defined as follows for any key $K \in \{0, 1\}^k$ and input $x \in \{0, 1\}^{2l}$ —

Function $F^{(r)}(K, x)$

Parse x as L_0R_0 with $|L_0| = |R_0| = l$

For $i = 1, \dots, r$ do

$L_i \leftarrow R_{i-1}; R_i \leftarrow F(K, R_{i-1}) \oplus L_{i-1}$

Return L_rR_r

- [20 points]** Show that $F^{(1)}$ is not a secure PRF by presenting a practical adversary A such that $\text{Adv}_{F^{(1)}}^{\text{prf}}(A)$ is close to one.
- [30 points]** Show that $F^{(2)}$ is not a secure PRF by presenting a practical adversary A such that $\text{Adv}_{F^{(2)}}^{\text{prf}}(A)$ is close to one.

For both **(1)** and **(2)** above, say what is the advantage achieved by your adversary. Also say what is its running time and the number of oracle queries it makes.
