- Always try to give algorithm with best possible running time. The points that you obtain will depend on the running time of your algorithm. For example, a student who gives an $O(n)$ algorithm will receive more points than a student who gives an $O\left(n^{2}\right)$ algorithm.
- You are required to give proofs of correctness whenever needed. For example, if you give a greedy algorithm for some problem, then you should also give a proof why this algorithm outputs optimal solution.
- Use of unfair means will be severely penalized.

There are 3 questions for a total of 50 points.
(15) 1. Recall the critical vertex problem discussed in class. Consider two vertices $s$ and $t$ in a given graph. A pair of vertices $(u, v)$ (different from $s$ and $t$ ) are called bi-critical with respect to $s$ and $t$ if the removal of $u$ and $v$ from the graph disconnects $s$ and $t$. Suppose in a given graph the shortest distance between $s$ and $t$ is strictly greater than $\lceil n / 3\rceil$. Prove or disprove the following statement:

There exists a pair of vertices that are bi-critical with respect to $s$ and $t$.
Give an algorithm for finding this pair of vertices in case there exists one.
(15) 2. A directed graph $G=(V, E)$ is called one-way-connected if for all pair of vertices $u$ and $v$ there is a path from vertex $u$ to $v$ or there is a path from vertex $v$ to $u$. Give an algorithm to check if a given graph is one-way-connected.
(20) 3. You are a party organizer and you need to solve the following problem. There are $n$ people and you know their friendship network. Your job is to decide a subset $S$ of people who will be invited to the party. The constraint that you need to satisfy is that every person in the subset $S$, is friends with at least five other people in $S$ and not friends with at least five other people in $S$. Assume that you are given the friendship network as a graph (assume adjacency list representation) where the edges denote friendships. Design an algorithm that maximizes the size of the set $S$.

