## Digital Design Basics

- Two voltage levels - high and low (1 and 0, true and false) Hence, the use of binary arithmetic/logic in all computers
- A transistor is a 3-terminal device that acts as a switch



## Logic Blocks

- A logic block has a number of binary inputs and produces a number of binary outputs - the simplest logic block is composed of a few transistors
- A logic block is termed combinational if the output is only a function of the inputs
- A logic block is termed sequential if the block has some internal memory (state) that also influences the output
- A basic logic block is termed a gate (AND, OR, NOT, etc.)

We will only deal with combinational circuits today

## Truth Table

- A truth table defines the outputs of a logic block for each set of inputs
- Consider a block with 3 inputs $\mathrm{A}, \mathrm{B}, \mathrm{C}$ and an output E that is true only if exactly 2 inputs are true



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| $A$ | $B$ | $C$ | $E$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |

Can be compressed by only representing cases that have an output of 1

## Boolean Algebra

- Equations involving two values and three primary operators:
- OR : symbol + , $X=A+B \rightarrow X$ is true if at least one of A or B is true
- AND : symbol. , $X=A . B \rightarrow X$ is true if both $A$ and $B$ are true
- NOT : symbol ${ }^{-}, X=\bar{A} \rightarrow X$ is the inverted value of $A$


## Boolean Algebra Rules

- Identity law : A + 0 = A ; A. $1=\mathrm{A}$
- Zero and One laws: A + $1=1$; A. $0=0$
- Inverse laws: $A \cdot \bar{A}=0 ; A+\bar{A}=1$
- Commutative laws : $\mathrm{A}+\mathrm{B}=\mathrm{B}+\mathrm{A}$; $\mathrm{A} . \mathrm{B}=\mathrm{B} . \mathrm{A}$
- Associative laws : $A+(B+C)=(A+B)+C$

$$
A \cdot(B \cdot C)=(A \cdot B) \cdot C
$$

- Distributive laws : A. $(B+C)=(A . B)+(A . C)$

$$
A+(B \cdot C)=(A+B) \cdot(A+C)
$$

## DeMorgan's Laws

- $\overline{A+B}=\bar{A} \cdot \bar{B}$
- $\overline{A \cdot B}=\bar{A}+\bar{B}$
- Confirm that these are indeed true


## Pictorial Representations



What logic function is this?


Source: H\&P textbook

## Boolean Equation

- Consider the logic block that has an output E that is true only if exactly two of the three inputs $\mathrm{A}, \mathrm{B}, \mathrm{C}$ are true

Multiple correct equations:

Two must be true, but all three cannot be true: $E=((A \cdot B)+(B \cdot C)+(A \cdot C)) \cdot \overline{(A \cdot B \cdot C)}$

Identify the three cases where it is true:
$E=(A \cdot B \cdot \bar{C})+(A \cdot C \cdot \bar{B})+(C \cdot B \cdot \bar{A})$

## Sum of Products

- Can represent any logic block with the AND, OR, NOT operators
- Draw the truth table
- For each true output, represent the corresponding inputs as a product
- The final equation is a sum of these products

| A | B | C | E |  |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |  |
| 0 | 0 | 1 | 0 | $(A \cdot B \cdot \bar{C})+(A \cdot C \cdot \bar{B})+(C \cdot B \cdot \bar{A})$ |
| 0 | 1 | 0 | 0 |  |
| 0 | 1 | 1 | 1 | - Can also use "product of sums" |
| 1 | 0 | 0 | 0 | - Any equation can be implemented with an array of ANDs, followed by |
| 1 | 0 | 1 | 1 | an array of ORs |
| 1 | 1 | 0 | 1 |  |
| 1 | 1 | 1 | 0 | 10 |

## NAND and NOR

- NAND : NOT of AND : A nand $B=\overline{A . B}$
- NOR : NOT of OR: $A$ nor $B=A+B$
- NAND and NOR are universal gates, i.e., they can be used to construct any complex logical function


## Common Logic Blocks - Decoder

Takes in N inputs and activates one of $2^{\mathrm{N}}$ outputs

| $\mathrm{I}_{0}$ | $\mathrm{I}_{1}$ | $\mathrm{I}_{2}$ | $\mathrm{O}_{0}$ | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ | $\mathrm{O}_{6}$ | $\mathrm{O}_{7}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |



## Common Logic Blocks - Multiplexor

- Multiplexor or selector: one of $N$ inputs is reflected on the output depending on the value of the $\log _{2} \mathrm{~N}$ selector bits


2-input mux

## Adder Algorithm



Truth Table for the above operations:

| A | B | Cin | Sum Cout |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 |  |
| 0 | 0 | 1 |  |
| 0 | 1 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |
| 1 | 1 | 1 |  |

## Adder Algorithm



Truth Table for the above operations:

| A | B | Cin | Sum | Cout |
| :--- | :--- | :--- | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

Equations:

$$
\begin{aligned}
\text { Sum }= & \text { Cin } \cdot \bar{A} \cdot \bar{B}+ \\
& \text { B } \cdot \overline{\operatorname{Cin}} \cdot \overline{\mathrm{A}}+ \\
& \text { A } \cdot \overline{\mathrm{Cin}} \cdot \overline{\mathrm{~B}}+ \\
& \text { A } \cdot \mathrm{B} \cdot \mathrm{Cin}
\end{aligned}
$$

Cout $=$ A. B. Cin + A. B. $\overline{\mathrm{Cin}}+$
A. Cin. $\bar{B}+$
B. Cin. $\bar{A}$
= A. B +
A. Cin +
B. Cin

## Carry Out Logic



Equations:

$$
\begin{aligned}
\text { Sum }= & \operatorname{Cin} \cdot \bar{A} \cdot \bar{B} \cdot+ \\
& \text { B } \cdot \operatorname{Cin} \cdot \bar{A}+ \\
& \text { A } \cdot \overline{\operatorname{Cin}} \cdot \bar{B}+ \\
& \text { A } \cdot
\end{aligned}
$$

Cout $=$ A. B. Cin + A. B. $\overline{\mathrm{Cin}}+$ A. Cin. $\bar{B}+$ B. Cin. $\overline{\mathrm{A}}$
= A. B +
A. Cin +
B. Cin

CarryOut

## 1-Bit ALU with Add, Or, And

- Multiplexor selects between Add, Or, And operations



## 32-bit Ripple Carry Adder

1-bit ALUs are connected
"in series" with the carry-out of 1 box going into the carry-in of the next box


## Incorporating Subtraction

Must invert bits of $B$ and add a 1

- Include an inverter
- Carryln for the first bit is 1
- The Carryln signal (for the first bit) can be the same as the Binvert signal


Source: H\&P textbook

## Incorporating NOR and NAND



## Control Lines

What are the values of the control lines and what operations do they correspond to?

|  | Ai | Bn | Op |
| :--- | :---: | :---: | :---: |
| AND | 0 | 0 | 00 |
| OR | 0 | 0 | 01 |
| Add | 0 | 0 | 10 |
| Sub | 0 | 1 | 10 |
| NAND | 1 | 1 | 01 |
| NOR | 1 | 1 | 00 |



## Incorporating slt

- Perform a - b and check the sign
- New signal (Less) that is zero for ALU boxes 1-31
- The $31^{\text {st }}$ box has a unit to detect overflow and sign - the sign bit serves as the Less signal for the $0^{\text {th }}$ box



## Incorporating beq

- Perform a - b and confirm that the result is all zero's



## Control Lines



## Control Lines

What are the values of the control lines and what operations do they correspond to?

|  | Ai | Bn | Op |
| :--- | :---: | :---: | :---: |
| AND | 0 | 0 | 00 |
| OR | 0 | 0 | 01 |
| Add | 0 | 0 | 10 |
| Sub | 0 | 1 | 10 |
| NOR | 1 | 1 | 00 |
| NAND | 1 | 1 | 01 |
| SLT | 0 | 1 | 11 |
| BEQ | 0 | 1 | 10 |



