

Finite-state methods for morphology

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Today's lecture

What are words? How many words are there?

What is the **structure of words**?
(in English, Chinese, Arabic,...)

Morphology: the area of linguistics that deals with this.

How can we identify the structure of words?

We need to build a **morphological analyzer** (parser).

We will use **finite-state transducers** for this task.

Finite-State Automata and Regular Languages
(Review)

Morphology: What is a word?

A Turkish word

uygarlaştıramadıklarımızdanmışsınızcasına
uygar_laş_tır_ama_dık_lar_ımız_dan_mış_sınız_casına

“as if you are among those whom we were not able to civilize (=cause to become civilized)”

uygar: civilized

_laş: become

_tır: cause somebody to do something

_ama: not able

_dık: past participle

_lar: plural

_ımız: 1st person plural possessive (our)

_dan: among (ablative case)

_miş: past

_sınız: 2nd person plural (you)

_casına: as if (forms an adverb from a verb) *K. Oflazer pc to J&M*

Basic word classes (parts of speech)

Content words (open-class):

- Nouns: student, university, knowledge,...
- Verbs: write, learn, teach,...
- Adjectives: difficult, boring, hard,
- Adverbs: easily, repeatedly,...

Function words (closed-class):

- Prepositions: in, with, under,...
- Conjunctions: and, or,...
- Determiners: a, the, every,...

How many words are there?

The Unix command "***wc -w***" counts the words in a file.

```
> cat example.txt
```

```
This company isn't New York-based anymore  
We moved to Chicago
```

```
> wc -w example.txt  
    10 example.txt
```

"***wc -w***" uses blanks to identify words:

```
This1 company2 isn't3 New4 York-based5 anymore6  
We7 moved8 to9 Chicago10
```

Words aren't just defined by blanks

Problem 1: Compounding

“ice cream”, “website”, “web site”, “New York-based”

Problem 2: Other writing systems have no blanks

Chinese: 我开始写小说 = 我 开始 写 小说
I start(ed) writing novel(s)

Problem 3: Clitics

English: “doesn't”, “I'm”,

Italian: “dirglielo” = dir + gli(e) + lo
tell + him + it

How many words are there?

Of course he wants to take the advanced course too.
He already took two beginners' courses.

This is a bad question. Did I mean:

How many **word tokens** are there?

(16 to 19, depending on how we count punctuation)

How many **word types** are there?

(i.e. How many different words are there?)

Again, this depends on how you count, but it's usually much less than the number of tokens)

How many words are there?

Of **course** he wants to **take** the advanced **course** **too**.
He already **took** **two** beginners' **courses**.

The same (underlying) word can take different forms:
course/courses, take/took

We distinguish concrete **word forms** (take, taking)
from abstract **lemmas** or dictionary forms (take)

Different words may be spelled/pronounced the same:
of course vs. advanced course
two vs. too

How many different words are there?

Inflection creates different forms of the same word:

Verbs: to be, being, I am, you are, he is, I was,

Nouns: one book, two books

Derivation creates different words from the same lemma:

grace \Rightarrow disgrace \Rightarrow disgraceful \Rightarrow disgracefully

Compounding combines two words into a new word:

cream \Rightarrow ice cream \Rightarrow ice cream cone \Rightarrow ice cream cone bakery

Word formation is productive:

New words are subject to all of these processes:

Google \Rightarrow Googler, to google, to ungoogle, to misgoogle, googlification, ungooglification, googlified, Google Maps, Google Maps service,...

Inflectional morphology in English

Verbs:

- Infinitive/present tense: walk, go
- 3rd person singular present tense (s-form): walks, goes
- Simple past: walked, went
- Past participle (ed-form): walked, gone
- Present participle (ing-form): walking, going

Nouns:

- Number: singular (book) vs. plural (books)
- Plural: books
- Possessive (~ genitive case): book's, books
- Personal pronouns inflect for person, number, gender, case:
I saw him; he saw me; you saw her; we saw them; they saw us.

Derivational morphology

Nominalization:

V + -ation: computerization

V+ -er: killer

Adj + -ness: fuzziness

Negation:

un-: undo, unseen, ...

mis-: mistake,...

Adjectivization:

V+ -able: doable

N + -al: national

Morphemes: stems, affixes

dis-grace-ful-ly
prefix-stem-suffix-suffix

Many word forms consist of a **stem** plus a number of **affixes** (*prefixes* or *suffixes*)

Infixes are inserted inside the stem.

Circumfixes (German gesehen) surround the stem

Morphemes: the smallest (meaningful/grammatical) parts of words.

Stems (grace) are often **free morphemes**.

Free morphemes can occur by themselves as words.

Affixes (dis-, -ful, -ly) are usually **bound morphemes**.

Bound morphemes have to combine with others to form words.

Morphemes and morphs

There are many *irregular word forms*:

- Plural nouns add -s to singular: book-books,
but: box-boxes, fly-flies, child-children
- Past tense verbs add -ed to infinitive: walk-walked,
but: like-liked, leap-leapt

Morphemes are abstract categories

Examples: plural morpheme, past tense morpheme

The same morpheme (e.g. for plural nouns) can be realized as *different surface forms (morphs)*:

-s/-es/-ren

Allomorphs: two different realizations (-s/-es/-ren)
of the same underlying morpheme (plural)

Morphological parsing and generation

Morphological parsing

disgracefully
dis *grace* *ful* *ly*
prefix *stem* *suffix* *suffix*
NEG *grace+N* *+ADJ* *+ADV*

Morphological generation

Generate possible English words:

grace, graceful, gracefully
disgrace, disgraceful, disgracefully,
ungraceful, ungracefully,
undisgraceful, undisgracefully,...

Don't generate impossible English words:

*gracelyful, *gracefuly, *disungracefully,...

Review: Finite-State Automata and Regular Languages

Formal languages

An **alphabet** Σ is a **set of symbols**:

e.g. $\Sigma = \{a, b, c\}$

A **string** ω is a **sequence of symbols**, e.g. $\omega = abcb$.

The **empty string** ε consists of zero symbols.

The Kleene closure Σ^* (**'sigma star'**) is the **(infinite) set of all strings** that can be formed from Σ :

$\Sigma^* = \{\varepsilon, a, b, c, aa, ab, ba, aaa, \dots\}$

A **language** $L \subseteq \Sigma^*$ over Σ is also a set of strings.

Typically we only care about **proper subsets of Σ^*** ($L \subset \Sigma^*$).

Automata and languages

An **automaton** is an abstract model of a computer which reads an **input string**, and **changes its internal state** depending on the current input symbol. It can either **accept or reject** the input string.

Every automaton defines a language

(the set of strings it accepts).

Different automata define different language classes:

- **Finite-state** automata define **regular** languages
- **Pushdown** automata define **context-free** languages
- **Turing machines** define **recursively enumerable** languages

Finite State Automata (FSAs)

A finite-state automaton $M = \langle Q, \Sigma, q_0, F, \delta \rangle$ consists of:

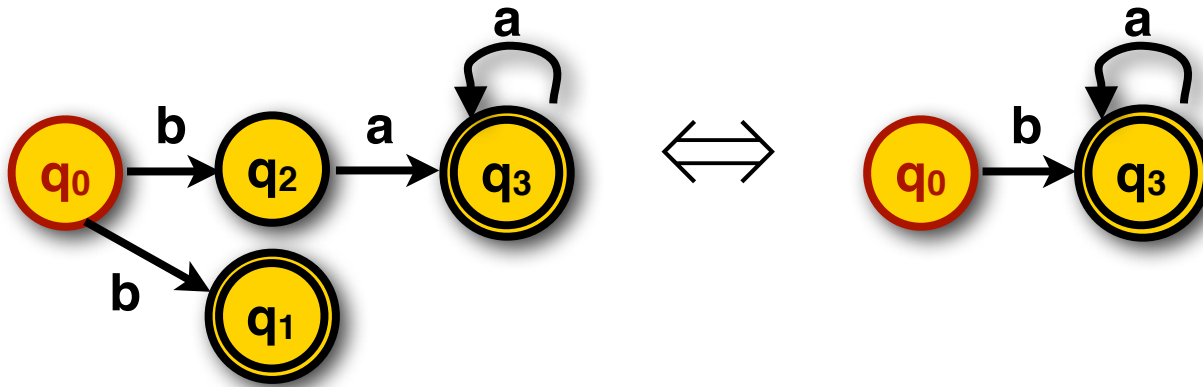
- A finite set of states $Q = \{q_0, q_1, \dots, q_n\}$
- A finite alphabet Σ of input symbols (e.g. $\Sigma = \{a, b, c, \dots\}$)
- A designated start state $q_0 \in Q$
- A set of final states $F \subseteq Q$
- A transition function δ :
 - The transition function for a **deterministic (D)FSA**: $Q \times \Sigma \rightarrow Q$
$$\delta(q, w) = q' \quad \text{for } q, q' \in Q, w \in \Sigma$$

If the current state is q and the current input is w , go to q'
 - The transition function for a **nondeterministic (N)FSA**: $Q \times \Sigma \rightarrow 2^Q$
$$\delta(q, w) = Q' \quad \text{for } q \in Q, Q' \subseteq Q, w \in \Sigma$$

If the current state is q and the current input is w , go to any $q' \in Q'$

Finite State Automata (FSAs)

Every NFA can be transformed into an equivalent DFA:



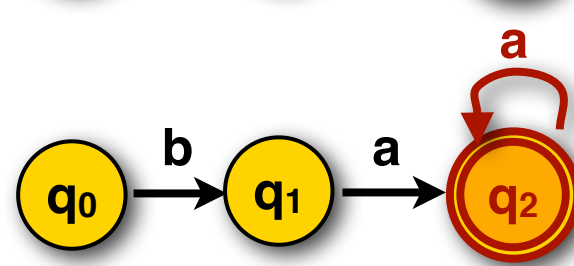
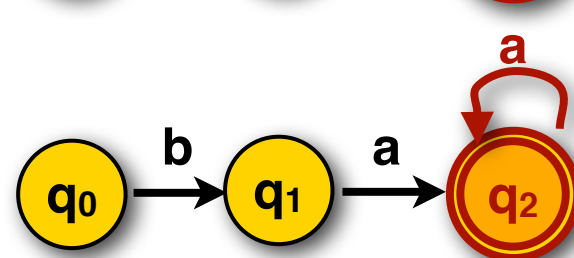
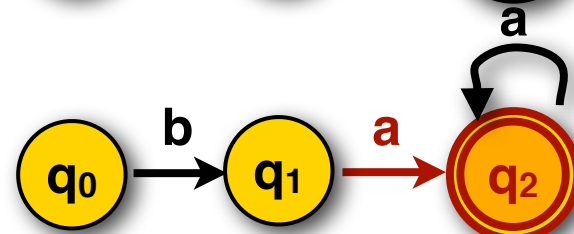
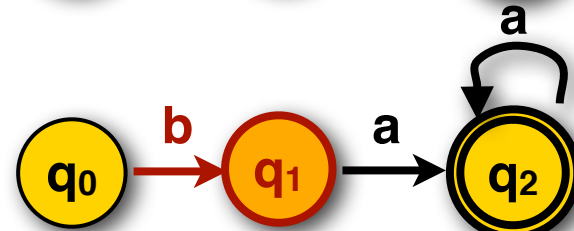
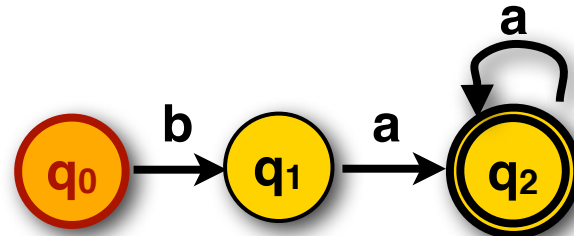
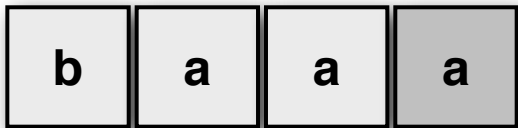
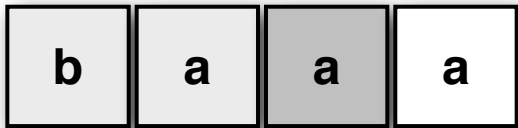
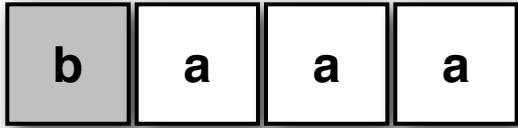
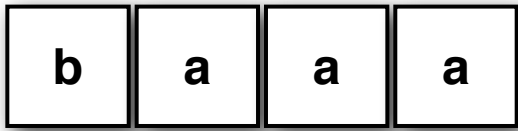
Recognition of a string w with a DFA is linear in the length of w

Finite-state automata define the class of **regular languages**

$L_1 = \{ a^n b^m \} = \{ ab, aab, abb, aaab, abb, \dots \}$ is a regular language,

$L_2 = \{ a^n b^n \} = \{ ab, aabb, aaabbb, \dots \}$ is not (it's context-free).

You cannot construct an FSA that accepts all the strings in L_2 and nothing else.



Regular Expressions

Simple patterns:

- **Standard characters** match themselves: ‘a’, ‘1’
- **Character classes**: ‘[abc]’, ‘[0-9]’, **negation**: ‘[^aeiou]’
(Predefined: \s (whitespace), \w (alphanumeric), etc.)
- **Any character** (except newline) is matched by ‘.’

Complex patterns: (e.g. `^[A-Z]([a-z])+\s`)

- **Group**: ‘(...)’
- **Repetition**: 0 or more times: ‘*’, 1 or more times: ‘+’
- **Disjunction**: ‘...|...’
- **Beginning of line** ‘^’ and **end of line** ‘\$’

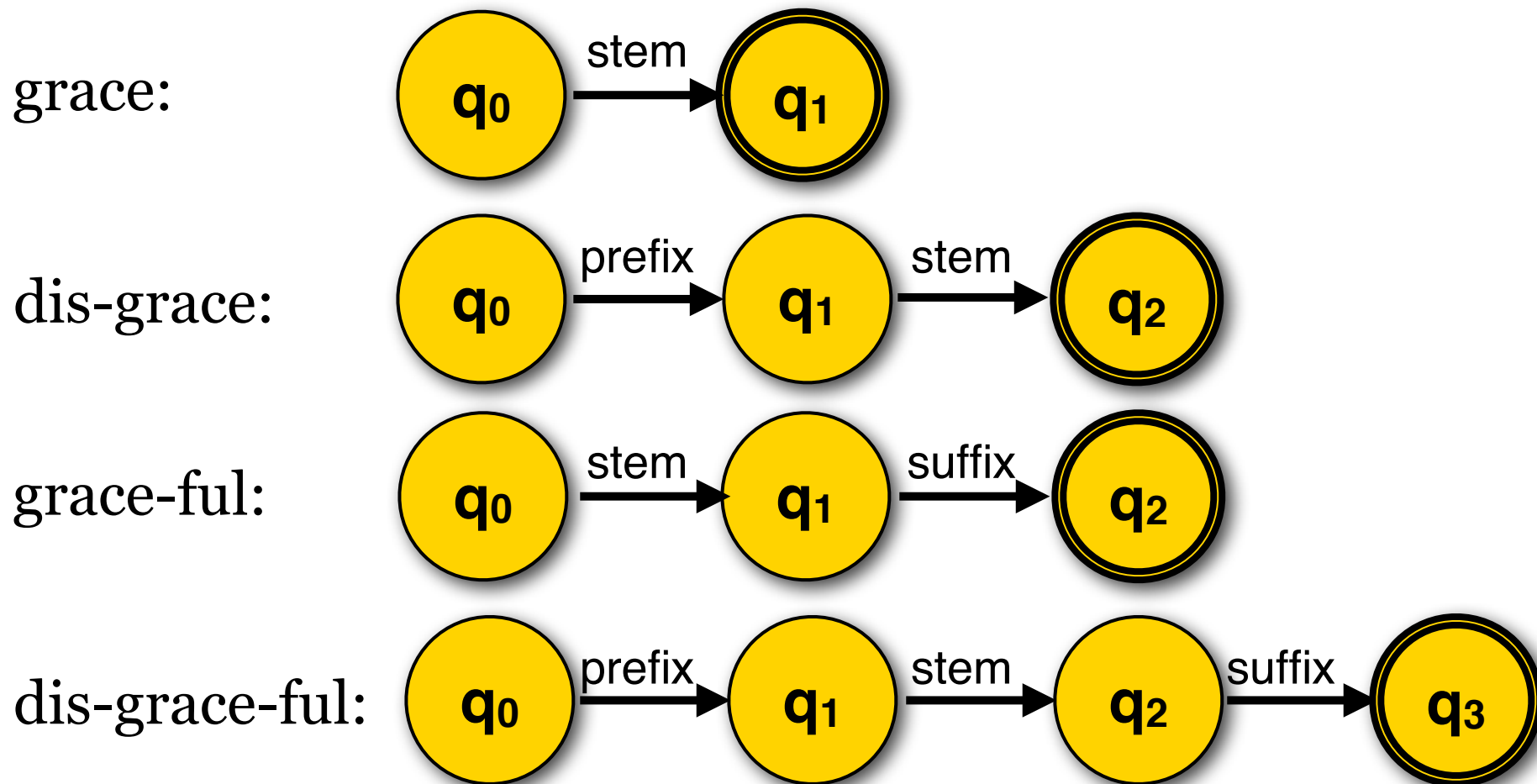
Python: Regular Expressions

```
>>> import re                                %% Import re package
>>> ex = re.compile('a.c')                   %% '...' : reg.expression
>>> m = ex.search('ab')                       %% Does 'ab' contain ex?
>>> print m                                   %% No.
None
>>> m = ex.search('abc')                     %% Does 'abc' contain ex?
>>> print m                                   %% Yes.
<_sre.SRE_Match object at 0x70640>
```

<http://docs.python.org/dev/howto/regex.html>
<http://docs.python.org/lib/module-re.html>

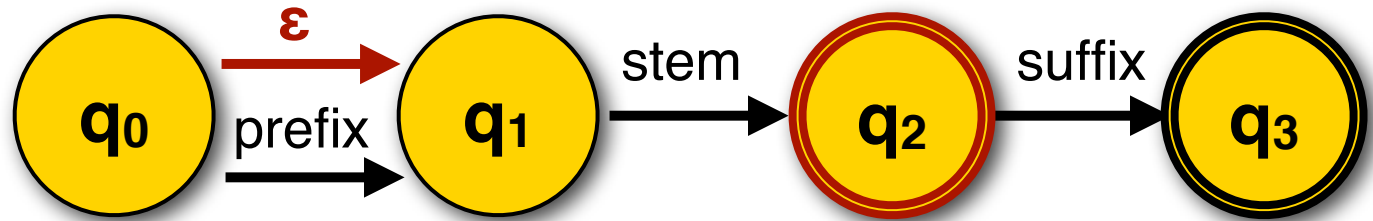
Finite-state methods for morphology

Finite state automata for morphology



Union: merging automata

grace,
dis-grace,
grace-ful,
dis-grace-ful



Stem changes

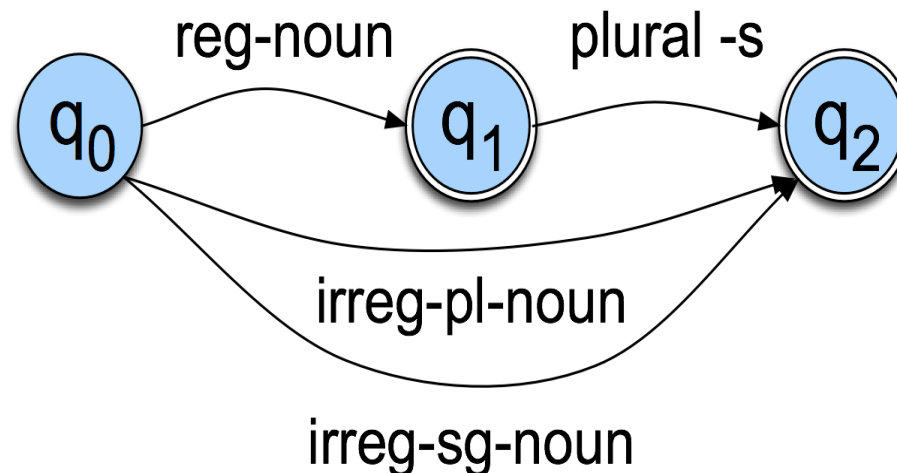
Some irregular words require stem changes:

Past tense verbs:

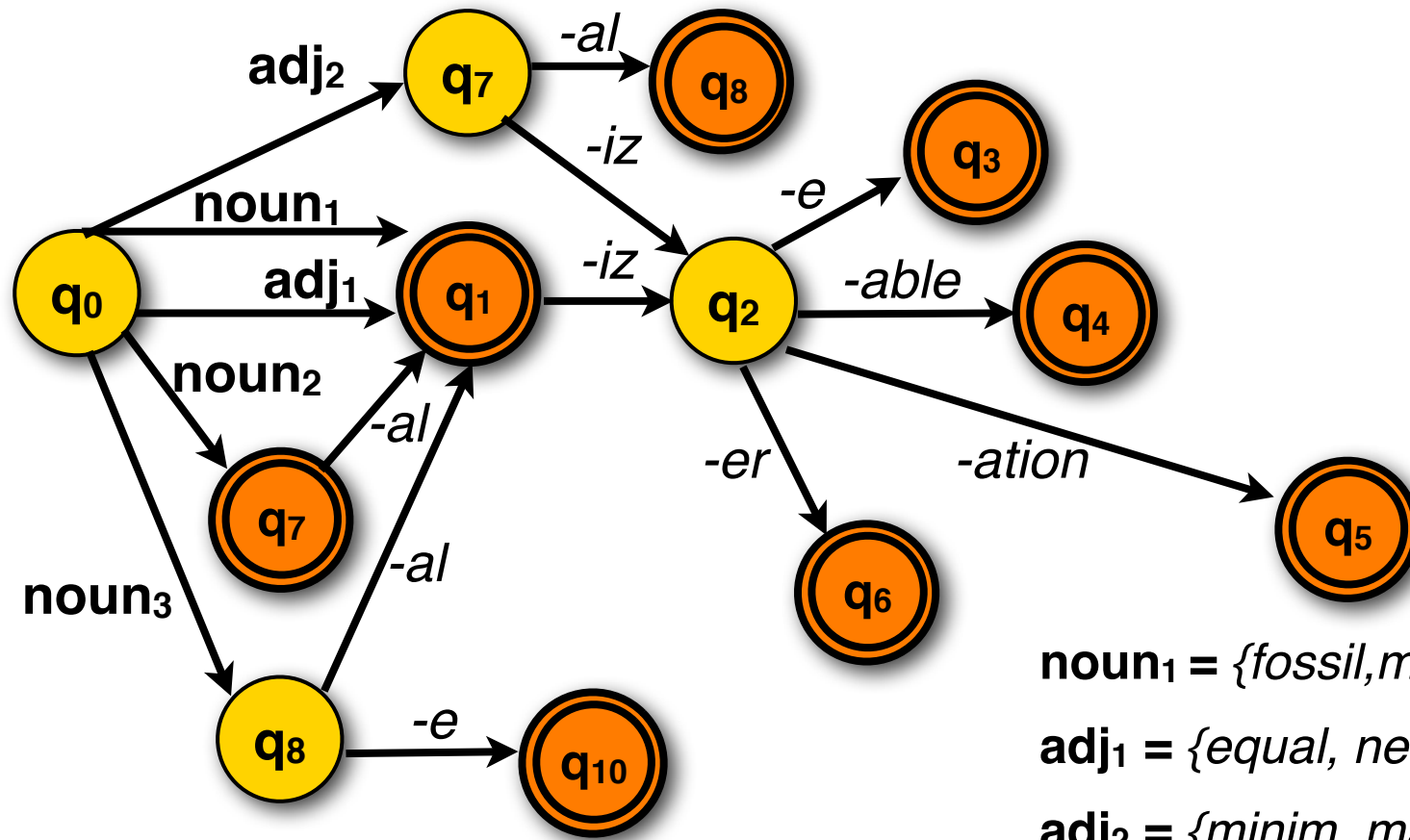
teach-taught, go-went, write-wrote

Plural nouns:

mouse-mice, foot-feet, wife-wives



FSAs for derivational morphology



noun₁ = {*fossil, mineral, ...*}

adj₁ = {*equal, neutral*}

adj₂ = {*minim, maxim*}

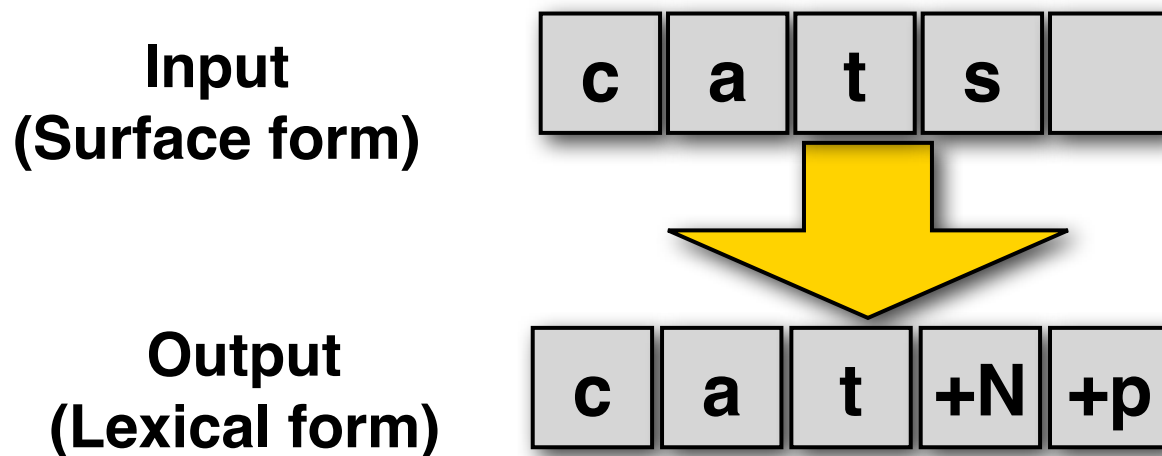
noun₂ = {*nation, form, ...*}

noun₃ = {*natur, structur, ...*}

Recognition vs. Analysis

FSAs can recognize (**accept**) a string, but they don't tell us its internal structure.

We need is a machine that maps (**transduces**) the input string into an output string that encodes its structure:



Finite-state transducers

A **finite-state transducer** $T = \langle Q, \Sigma, \Delta, q_0, F, \delta, \sigma \rangle$ consists of:

- A finite **set of states** $Q = \{q_0, q_1, \dots, q_n\}$
 - A finite alphabet Σ of **input symbols** (e.g. $\Sigma = \{a, b, c, \dots\}$)
 - A finite alphabet Δ of **output symbols** (e.g. $\Delta = \{+N, +pl, \dots\}$)
 - A designated **start state** $q_0 \in Q$
 - A set of **final states** $F \subseteq Q$
 - A **transition function** $\delta: Q \times \Sigma \rightarrow 2^Q$
 $\delta(q, w) = Q'$ for $q \in Q, Q' \subseteq Q, w \in \Sigma$
 - **An output function** $\sigma: Q \times \Sigma \rightarrow \Delta^*$
 $\sigma(q, w) = \omega$ for $q \in Q, w \in \Sigma, \omega \in \Delta^*$
- If the current state is q and the current input is w , write ω .

Finite-state transducers

An FST $T = L_{in} \times L_{out}$ defines a relation between two regular languages L_{in} and L_{out} :

$L_{in} = \{\mathbf{cat}, \mathbf{cats}, \mathbf{fox}, \mathbf{foxes}, \dots\}$

$L_{out} = \{cat+N+sg, cat+N+pl, fox+N+sg, fox+N+PL \dots\}$



$T = \{ \langle \mathbf{cat}, cat+N+sg \rangle, \langle \mathbf{cats}, cat+N+pl \rangle, \langle \mathbf{fox}, fox+N+sg \rangle, \langle \mathbf{foxes}, fox+N+pl \rangle \}$

Some FST operations

Inversion T^{-1} :

The inversion (T^{-1}) of a transducer switches input and output labels.

*This can be used to switch from **parsing** words to **generating** words.*

Composition ($T \circ T'$): (*Cascade*)

Two transducers $T = L_1 \times L_2$ and $T' = L_2 \times L_3$ can be composed into a third transducer $T'' = L_1 \times L_3$.

*Sometimes **intermediate representations** are useful*

English spelling rules

English spelling (orthography) is funny:

The underlying morphemes (*plural-s*, etc.) can have different orthographic surface realizations (-s, -es)

Spelling changes at morpheme boundaries:

- E-insertion: fox +s = fox**e**s
- E-deletion: make**e** +ing = making

Intermediate representations

English plural -s: **cat** \Rightarrow **cats** **dog** \Rightarrow **dogs**
but: **fox** \Rightarrow **foxes**, **bus** \Rightarrow **buses** **buzz** \Rightarrow **buzzes**

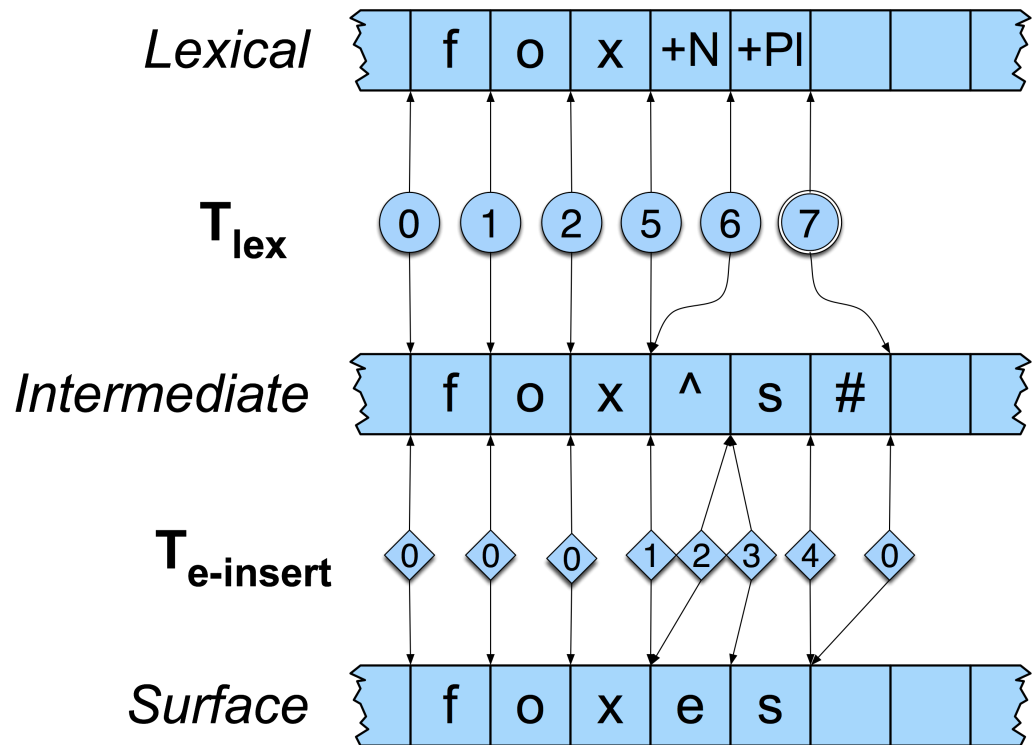
We define an **intermediate representation** which captures morpheme boundaries (^) and word boundaries (#):

<i>Lexicon:</i>	cat+N+PL	fox+N+PL
\Rightarrow <i>Intermediate representation:</i>	cat^s#	fox^s#
\Rightarrow <i>Surface string:</i>	cats	foxes

Intermediate-to-Surface Spelling Rule:

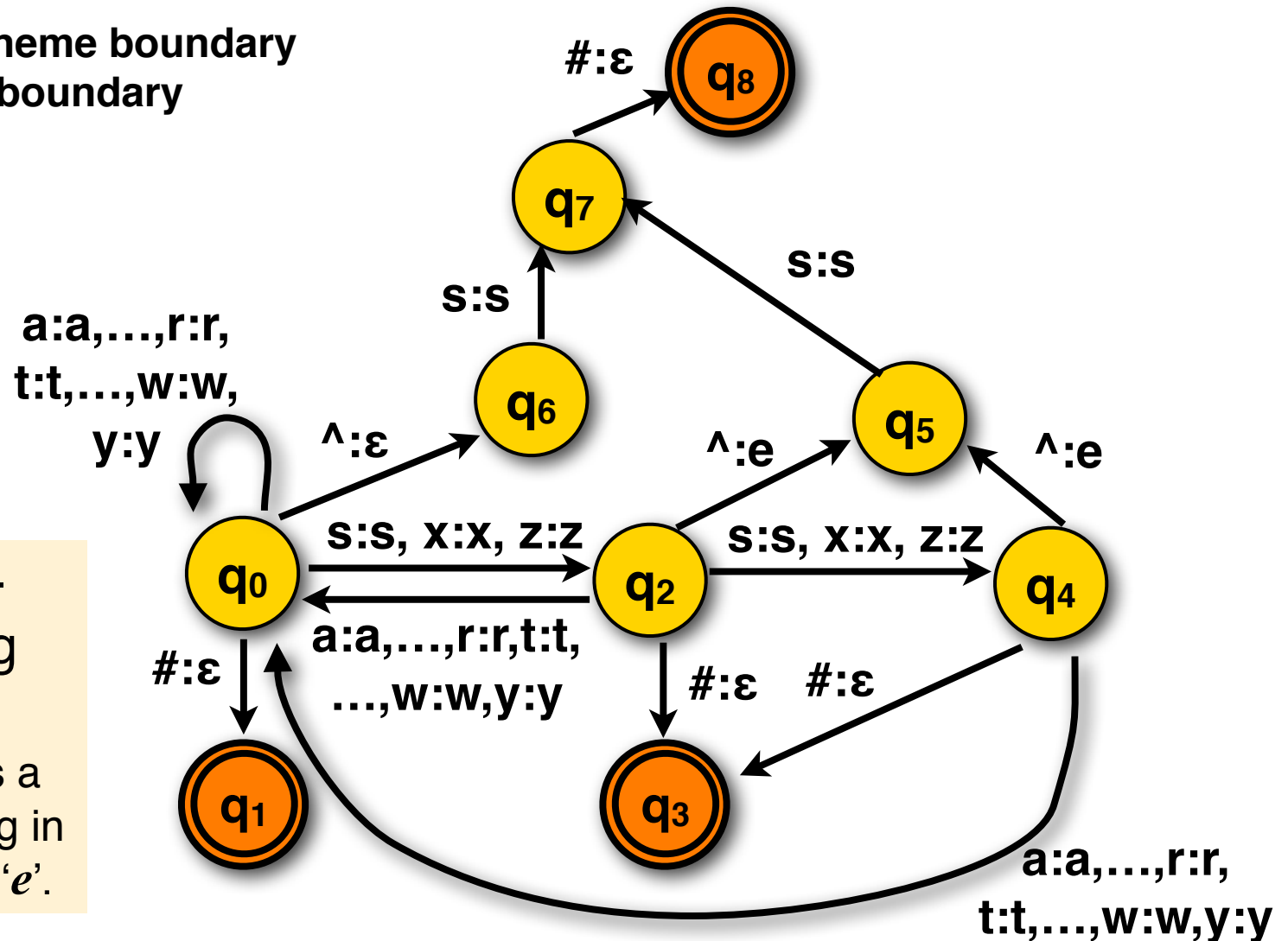
If plural '**s**' follows a morpheme ending in '**x**', '**z**' or '**s**', insert '**e**'.

FST composition/cascade:



FST: intermediate to surface level

^ = morpheme boundary
= word boundary



Intermediate-to-Surface Spelling Rule:
If plural 's' follows a morpheme ending in 'x', 'z' or 's', insert 'e'.

Dealing with ambiguity

book: book +N +sg or book +V?

Generating words is generally unambiguous, but **analyzing** words often requires disambiguation.

Efficiency problem:

Not every nondeterministic FST can be translated into a deterministic one!

A computer therapist?

Computer: TELL ME MORE ABOUT YOUR FAMILY

Human: My mother takes care of me.

Computer: WHO ELSE IN YOUR FAMILY TAKES CARE OF YOU

Human: My father.

Computer: YOUR FATHER

Human: You are like my father in some ways.

Computer: WHAT RESEMBLANCE DO YOU SEE

Human: You are not very aggressive but I think you don't want me to notice that.

Computer: WHAT MAKES YOU THINK I AM NOT VERY AGGRESSIVE

Human: You don't argue with me.

Computer: WHY DO YOU THINK I DON'T ARGUE WITH YOU

Human: You are afraid of me.

Weizenbaum (1966), ELIZA.

ELIZA as a FST cascade

Human: *You don't argue with me.*

Computer: *WHY DO YOU THINK I DON'T ARGUE WITH YOU*

1. Replace **you** with *I* and **me** with **you**:

I don't argue with you.

2. Replace **<...>** with **Why do you think <...>**:

Why do you think I don't argue with you.

What about compounds?

Compounds have hierarchical structure:

((ice cream) cone) bakery)
not (ice ((cream cone) bakery))

((computer science) (graduate student))
not (computer ((science graduate) student))

We will need context-free grammars to capture this underlying structure.