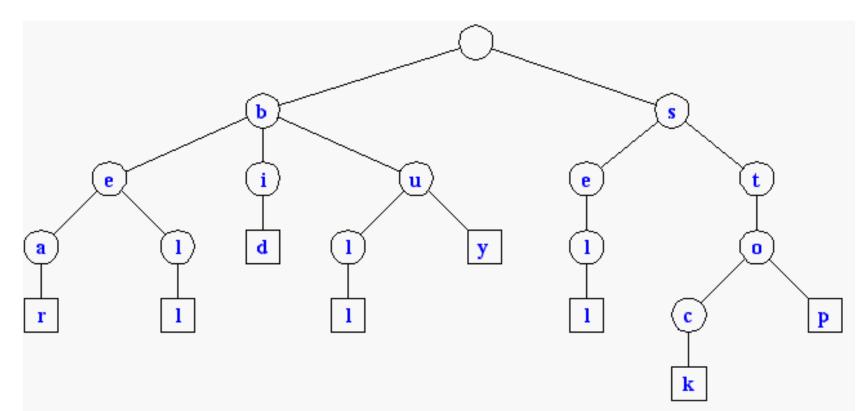
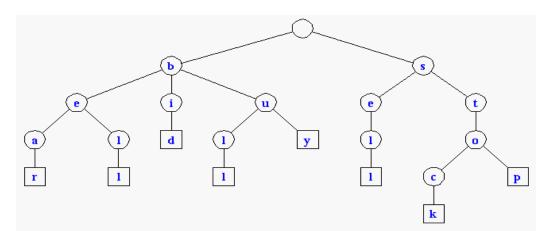
Tries

- Standard Tries
- Compressed Tries
- Suffix Tries



Standard Tries

- The *standard trie* for a set of strings S is an ordered tree such that:
 - each node but the root is labeled with a character
 - the children of a node are alphabetically ordered
 - the paths from the external nodes to the root yield the strings of S
- Example: standard trie for the set of strings
 - S = { bear, bell, bid, bull, buy, sell, stock, stop }



•A standard trie uses O(n) space. Operations (find, insert, remove) take time O(dm) each, where:

-n = total size of the strings in S,

-m =size of the string parameter of the operation

-d =alphabet size,

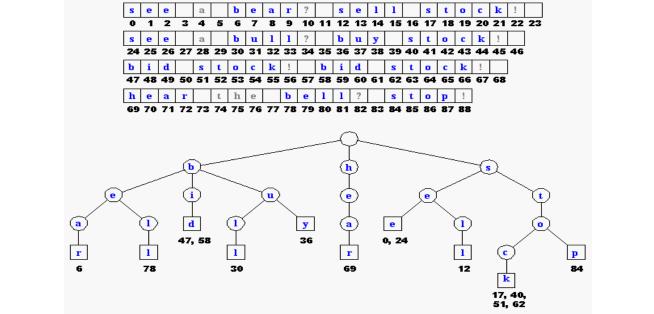
Applications of Tries

• A standard trie supports the following operations on a preprocessed text in time O(m), where m = |X|

-word matching: find the first occurence of word X in the text

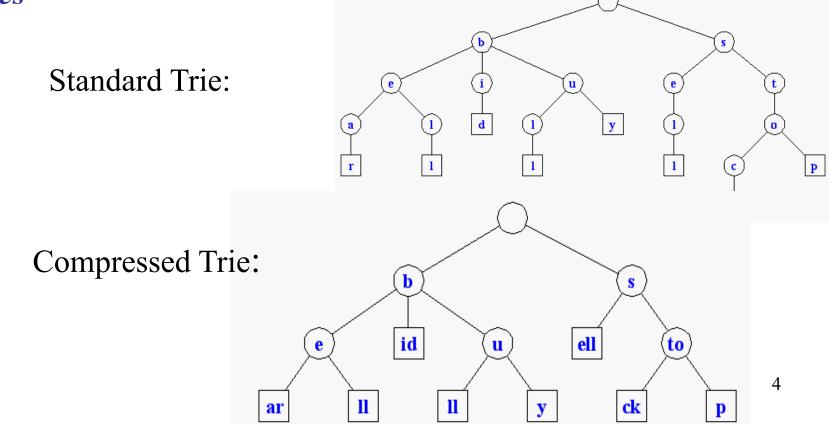
-prefix matching: find the first occurrence of the longest prefix of word X in the text

• Each operation is performed by tracing a path in the trie starting at the root



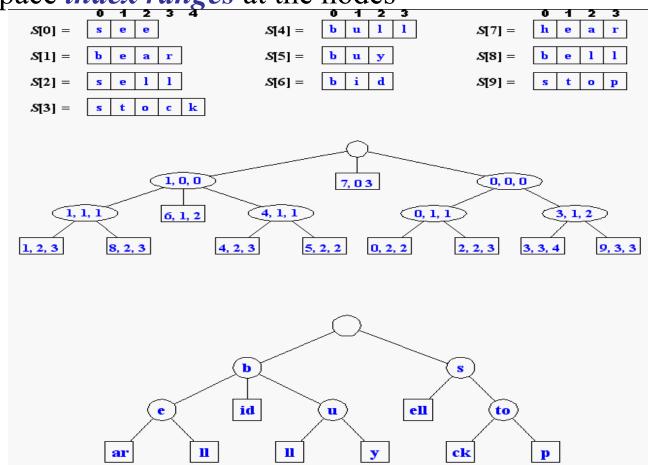
Compressed Tries

- Trie with nodes of degree at least 2
- Obtained from standard trie by compressing chains of *redundant* nodes



Compact Storage of Compressed Tries

• A compressed trie can be stored in space O(s), where s = |S|, by using O(1) space *index ranges* at the nodes

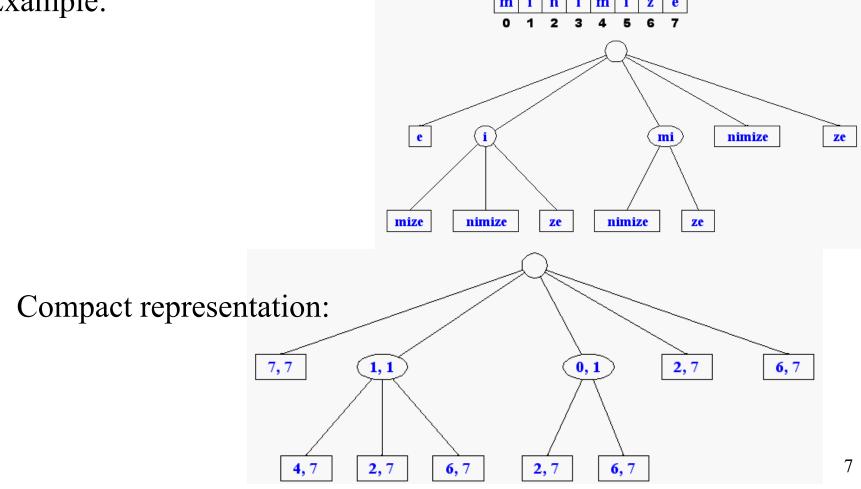


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Insertion and Deletion into/from a Compressed Trie b abbb abab baab b aaa bab search stops here insert(bbaabb) b baab abbb b abab bab aa bb

Suffix Tries

A *suffix trie* is a compressed trie for all the suffixes of a text
Example: <u>m i n i m i z e</u>

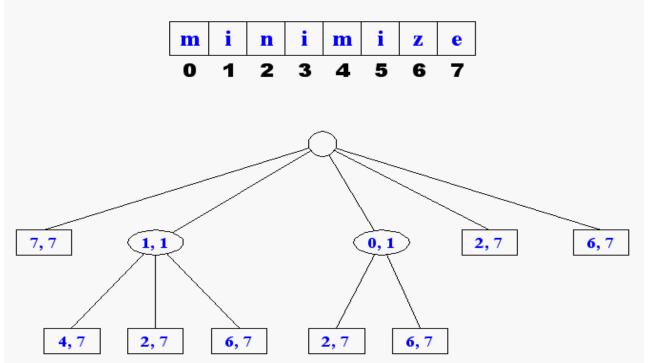


Properties of Suffix Tries

The *suffix trie* for a text X of size *n* from an alphabet of size *d* -stores all the *n(n-1)/2 suffixes* of X in O(n) space
 -supports arbitrary *pattern matching* and prefix matching queries in

O(dm) time, where m is the length of the pattern

-can be constructed in O(dn) time



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Tries and Web Search Engines

- The *index of a search engine* (collection of all searchable words) is stored into a compressed trie
- Each leaf of the trie is associated with a word and has a list of pages (URLs) containing that word, called *occurrence list*
- The trie is kept in internal memory
- The occurrence lists are kept in external memory and are ranked by relevance
- Boolean queries for sets of words (e.g., Java and coffee) correspond to set operations (e.g., intersection) on the occurrence lists
- Additional *information retrieval* techniques are used, such as
 - stopword elimination (e.g., ignore "the" "a" "is")
 - stemming (e.g., identify "add" "adding" "added")
 - link analysis (recognize authoritative pages)

Tries and Internet Routers

- Computers on the internet (hosts) are identified by a unique 32-bit IP (*internet protocol*) addres, usually written in "dotted-quad-decimal" notation
- E.g., www.iitd.ac.in is 103.27.9.24
- Use nslookup on Unix to find out IP addresses
- An organization uses a subset of IP addresses with the same prefix, e.g., IITD uses 103.27.*.*, Yale uses 130.132.*.*
- Data is sent to a host by fragmenting it into packets. Each packet carries the IP address of its destination.
- The internet whose nodes are *routers*, and whose edges are communication links.
- A router forwards packets to its neighbors using IP *prefix matching* rules. E.g., a packet with IP prefix 103.27. should be forwarded to the IITD gateway router.
- Routers use tries on the alphabet 0,1 to do prefix matching.