COL 351

TUTORIAL SHEET 5

- 1. You are given a positive integer N. You want to reach N by starting from 1, and performing one of the following two operations at each step: (i) increment the current number by 1, or (ii) double the current number. For example, if N = 10, you can start with 1, and then go in the sequence 1, 2, 4, 8, 9, 10, or 1, 2, 4, 5, 10. Give an efficient algorithm, which given the number N finds the minimum number of such operations to go from 1 to N.
- 2. Suppose you are given two sets of n points, one set $\{p_1, \ldots, p_n\}$ on the y = 0 line and the other set $\{q_1, \ldots, q_n\}$ on the y = 1 line. Now we draw n line segments given by $[p_i, q_i], i = 1, \ldots, n$. Describe and analyze a divide-and-conquer algorithm to determine how many pairs of these line segments intersect, in $O(n \log n)$ time.
- 3. You are given a rooted binary tree T where each vertex v has an integer val(v) stored in it (you can assume that all the integers involved are distinct). A vertex v is said to be a *local minimum* if $val(v) \leq val(w)$ for all the neighbours w of v (i.e., its value is at most the value at its children and at its parent). Show how to find a local minimum in O(H) time, where H is the height of T.

Solve the same problem when the graph is an $n \times n$ grid graph. An $n \times n$ grid graph has vertices labelled (i, j), where $1 \leq i, j \leq n$ and (i, j) is adjacent to (i - 1, j), (i + 1, j), (i, j - 1), (i, j + 1) (with appropriate restrictions at the boundary). The time taken by the algorithm should be O(n).

- 4. (a) Let $n = 2^{l} 1$ for some positive integer l. Suppose someone claims to hold an unsorted array A[1...n] of distinct l-bit strings; thus, exactly one l-bit string does not appear in A. Suppose further that the only way we can access A is by calling the function FB(i, j), which returns the j^{th} bit of the string A[i] in O(1) time. Describe an algorithm to find the missing string in A using only O(n) calls to FB.
 - (b) (Hard) Now suppose $n = 2^{l} k$ for some positive integers k and l, and again we are given an array A[1...n] of distinct *l*-bit strings. Describe an algorithm to find the k strings that are missing from A using only $O(n \log k)$ calls to FB.
- 5. You are given *n* charged particles located at points with coordinates $\{1, 2, ..., n\}$ on the real line. The point at coordinate *i* has charge q_i . Note that the total Coulomb force on a particle at location *i* is given by

$$-\sum_{j < i} \frac{q_i q_j}{(i-j)^2} + \sum_{j > i} \frac{q_i q_j}{(i-j)^2}.$$

Give an $O(n \log n)$ time algorithm to compute the total force on each of the particles.

6. We are given a sequence of n distinct integers a_1, \ldots, a_n in an array. An inversion is defined as a pair (i, j) such that i < j but $a_i > a_j$. Give an $O(n \log n)$ time to count the number of inversions in the array.

A strong inversion is defined as a pair (i, j) such that i < j but $a_i > 2a_j$. Give an $O(n \log n)$ time to count the number of strong inversions in the array.