## TUTORIAL SHEET 5

1. You are given a positive integer $N$. You want to reach $N$ by starting from 1 , and performing one of the following two operations at each step: (i) increment the current number by 1 , or (ii) double the current number. For example, if $N=10$, you can start with 1 , and then go in the sequence $1,2,4,8,9,10$, or $1,2,4,5,10$. Give an efficient algorithm, which given the number $N$ finds the minimum number of such operations to go from 1 to $N$.
2. Suppose you are given two sets of $n$ points, one set $\left\{p_{1}, \ldots, p_{n}\right\}$ on the $y=0$ line and the other set $\left\{q_{1}, \ldots, q_{n}\right\}$ on the $y=1$ line. Now we draw $n$ line segments given by $\left[p_{i}, q_{i}\right], i=1, \ldots, n$. Describe and analyze a divide-and-conquer algorithm to determine how many pairs of these line segments intersect, in $O(n \log n)$ time.
3. You are given a rooted binary tree $T$ where each vertex $v$ has an integer $\operatorname{val}(v)$ stored in it (you can assume that all the integers involved are distinct). A vertex $v$ is said to be a local minimum if $\operatorname{val}(v) \leq \operatorname{val}(w)$ for all the neighbours $w$ of $v$ (i.e., its value is at most the value at its children and at its parent). Show how to find a local minimum in $O(H)$ time, where $H$ is the height of $T$.
Solve the same problem when the graph is an $n \times n$ grid graph. An $n \times n$ grid graph has vertices labelled $(i, j)$, where $1 \leq i, j \leq n$ and $(i, j)$ is adjacent to $(i-1, j),(i+$ $1, j),(i, j-1),(i, j+1)$ (with appropriate restrictions at the boundary). The time taken by the algorithm should be $O(n)$.
4. (a) Let $n=2^{l}-1$ for some positive integer $l$. Suppose someone claims to hold an unsorted array $A[1 \ldots n]$ of distinct $l$-bit strings; thus, exactly one $l$-bit string does not appear in $A$. Suppose further that the only way we can access $A$ is by calling the function $F B(i, j)$, which returns the $j^{\text {th }}$ bit of the string $A[i]$ in $O(1)$ time. Describe an algorithm to find the missing string in $A$ using only $O(n)$ calls to $F B$.
(b) (Hard) Now suppose $n=2^{l}-k$ for some positive integers $k$ and $l$, and again we are given an array $A[1 \ldots n]$ of distinct $l$-bit strings. Describe an algorithm to find the $k$ strings that are missing from $A$ using only $O(n \log k)$ calls to $F B$.
5. You are given $n$ charged particles located at points with coordinates $\{1,2, \ldots, n\}$ on the real line. The point at coordinate $i$ has charge $q_{i}$. Note that the total Coulomb force on a particle at location $i$ is given by

$$
-\sum_{j<i} \frac{q_{i} q_{j}}{(i-j)^{2}}+\sum_{j>i} \frac{q_{i} q_{j}}{(i-j)^{2}} .
$$

Give an $O(n \log n)$ time algorithm to compute the total force on each of the particles.
6. We are given a sequence of $n$ distinct integers $a_{1}, \ldots, a_{n}$ in an array. An inversion is defined as a pair $(i, j)$ such that $i<j$ but $a_{i}>a_{j}$. Give an $O(n \log n)$ time to count the number of inversions in the array.
A strong inversion is defined as a pair $(i, j)$ such that $i<j$ but $a_{i}>2 a_{j}$. Give an $O(n \log n)$ time to count the number of strong inversions in the array.

