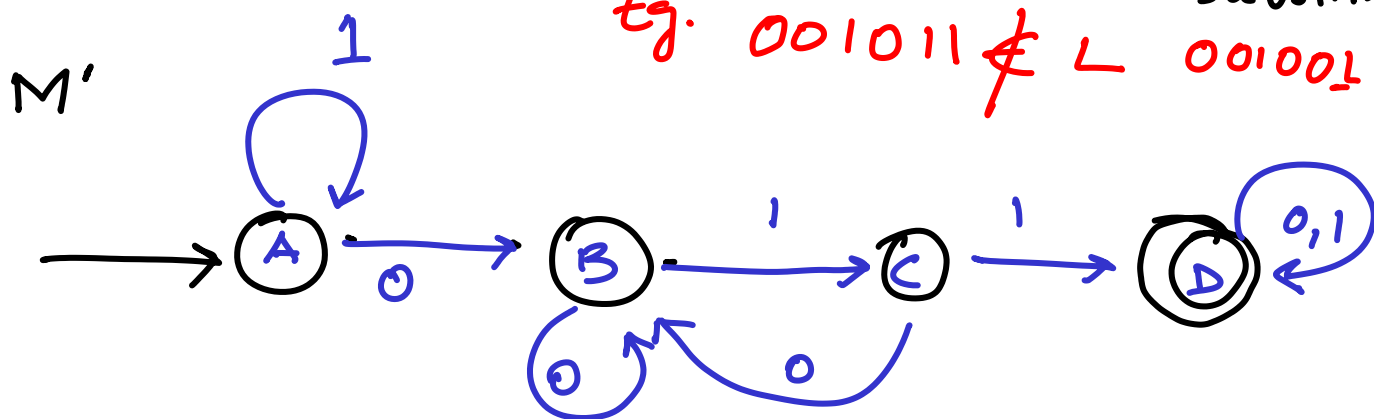


Design a DFA M for the language $L = \{ w \in \{0,1\}^* \mid w \text{ does not contain } 011 \text{ as a substring} \}$

Eg. $001011 \notin L$ $001001 \in L$



With a little effort, you can see that M' accepts all strings containing the pattern 011

(Try proving it rigorously)

Observation: M can be obtained from M' by flipping the accept/non-accept states since

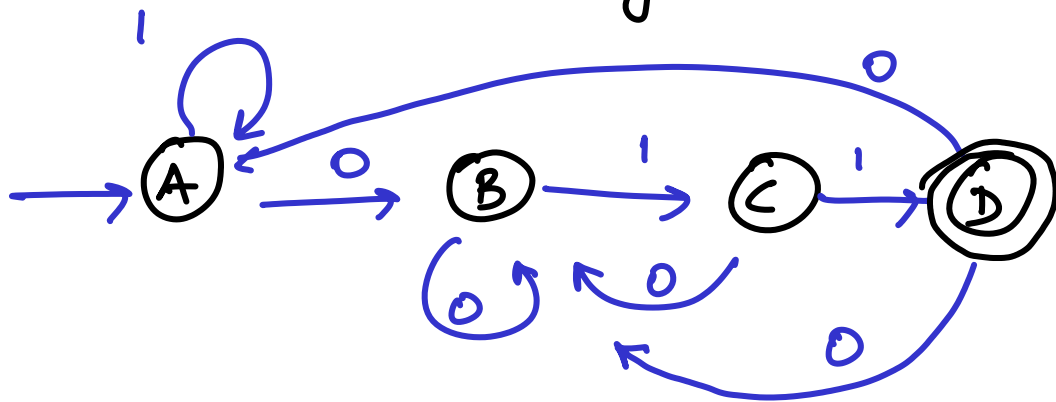
$$L(M) = \Sigma^* - L(M')$$

\Rightarrow if there is a DFA for L then
 there is a DFA for $\bar{L} = \Sigma^* - L$
 (Complement of L)

Problem 2: Design a machine that
 accepts all strings ending with 011

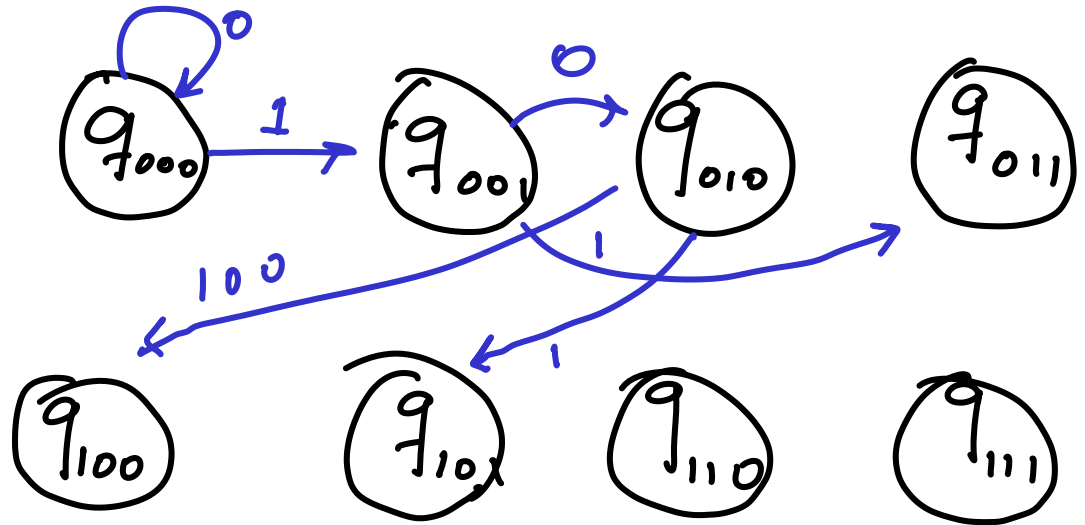
Example $01101101101 \notin L$ $\{0, 00, 10,$
 $001100011 \in L$ $11\} \notin L$

You can modify the previous machine



Prove that it works correctly

Another clean soln



$$\delta(q_{x_1 x_2 x_3}, 0) = q_{x_2 x_3 0}$$

$$\delta(q_{x_1 x_2 x_3}, 1) = q_{x_2 x_3 1}$$

$x_1 x_2 x_3$ are the most recent 3 inputs

Add some more states to initialize