

Can we use M-N theorem to show that certain languages are not regular?

$L = \{0^i 1^i \mid i \geq 0\}$. How many eqv classes does it require in R_L

$x R_L y \iff \forall z \in \Sigma^* \quad x \cdot z \in L \text{ exactly when } y \cdot z \in L$

$0^{i_1} \not\sim_{R_L} 0^{i_2} \quad i_1 \neq i_2$

Prob 5 (c) $L^R = \{x^R \mid x \in L\}$

x^R is the reverse of x and L is regular

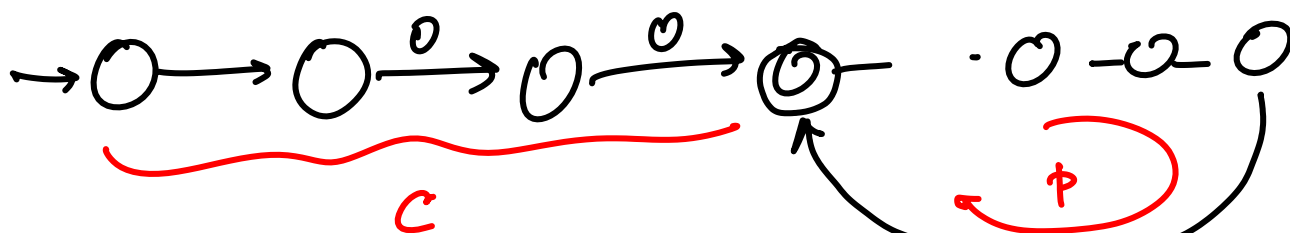
$(a_1 a_2 a_3 \dots a_k)^R = a_k a_{k-1} \dots a_1$ Is L^R regular

Prob 1

A set of integers is linear if it is of the form $\{c + p \cdot i \mid i \geq 0\}$ for some fixed c, p .

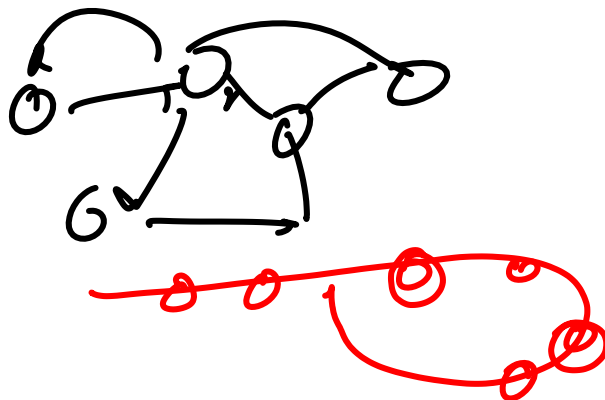
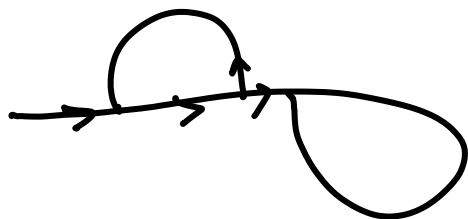
A semilinear set is a finite union of linear sets.

Let $R \subset \mathbb{O}^n$ be a regular language
show that R is semilinear

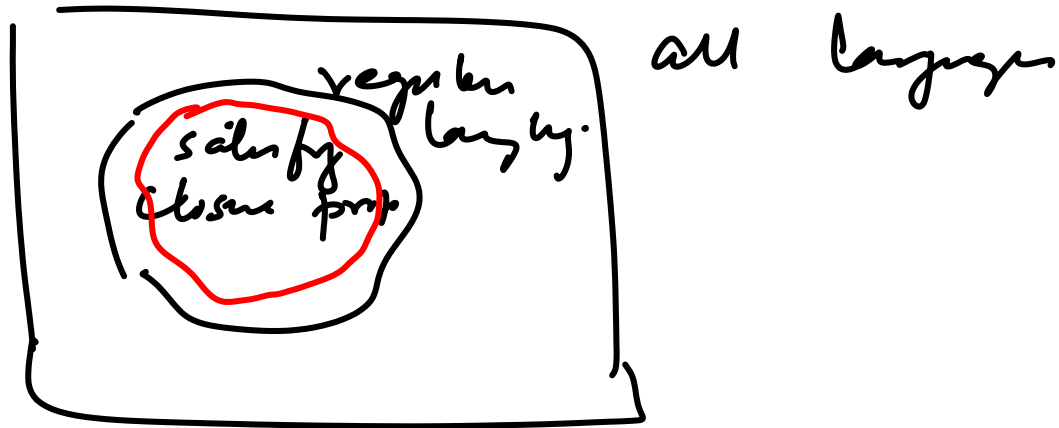


A DFA for linear set

DFA for R



7. What is the relationship between class of regular languages and the least class of languages closed under union, intersection, complement and that contains all finite sets.



This class, say \mathcal{C} must contain all finite languages and complement of finite languages

$\{ L_1, L_2, \dots, \bar{L}_1, \bar{L}_2, \dots \}$

Is \mathcal{C} closed under union & intersection