

$$R_{ij}^{(k)} = R_{ik}^{(k-1)} \begin{bmatrix} R_{kk}^{(k-1)*} \\ R_{kj}^{(k-1)} \end{bmatrix} + R_{ij}^{(k-1)}$$

$1^* 0 1 (0 + 11)^*$

| | R_{11} | R_{12} | R_{13} | R_{21} | R_{22} | R_{23} | R_{31} | R_{32} | R_{33} |
|------------------|---|---|-------------|-------------|------------|----------|-------------|----------|---------------------|
| $k=0$ | $1 + \epsilon$ | 0 | \emptyset | \emptyset | ϵ | 1 | \emptyset | 1 | $0 + \epsilon$ |
| q_1 is allowed | $R_{11}^{(0)}$ $(R_{11}^{(0)})^*$ $R_{11}^{(0)}$ $(1 + \epsilon) \cdot 1^*$ $\cdot (1 + \epsilon)$ $= 1^*$ | R_{11}^0 $(R_{11}^0)^*$ $R_{12}^0 + R_{11}^0$ $(1 + \epsilon) \cdot$ $(1 + \epsilon)^* \cdot 0$ $+ 0$ $= 1^* 0$ | \emptyset | \emptyset | ϵ | 1 | \emptyset | 1 | $0 + \epsilon$ |
| 2 | 1^* | $1^* 0$ | $1^* 0 1$ | \emptyset | ϵ | 1 | \emptyset | 1 | $0 + \epsilon + 11$ |

$$(1+\varepsilon)^x = \varepsilon + (1+\varepsilon) + (1+\varepsilon)(1+\varepsilon) + \dots + (1+\varepsilon)^i$$

$$= \varepsilon + 1 + \varepsilon + 1 + \varepsilon + \varepsilon + \underbrace{\varepsilon^2}_{\varepsilon} + \dots$$

$$= 1^x$$

Summary

1. Defined DFA / NFA

Regular Language

DFA \sim NFA

2. R.e.

R.e. \sim DFA

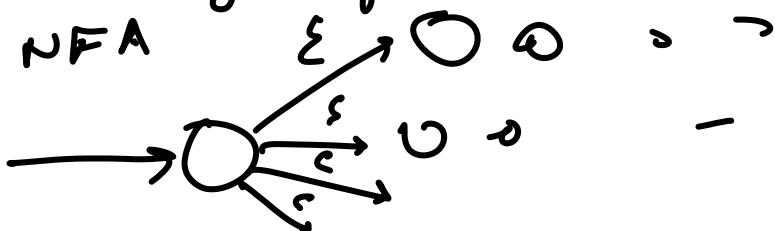
What languages are not regular?

1. There exists languages that are not regular.

$$L = \{ 0^i 1^i \mid i \geq 0 \}$$

$$\{ \epsilon + 01 + 0011 + 000111 + \dots + 0^{100} 1^{100} \}$$

Claim Any finite language is regular



Claim : If L_1 is regular and L_2 is regular so is $L_1 \cup L_2$

$\pi_1 + \pi_2$ is regular

$\pi_1 \downarrow$

π_2

$$(L_1) \cup (L_2 \cup L_3 \cup \dots \cup L_k)$$

Claim If L_1, L_2 are regular so is $L_1 \cdot L_2$

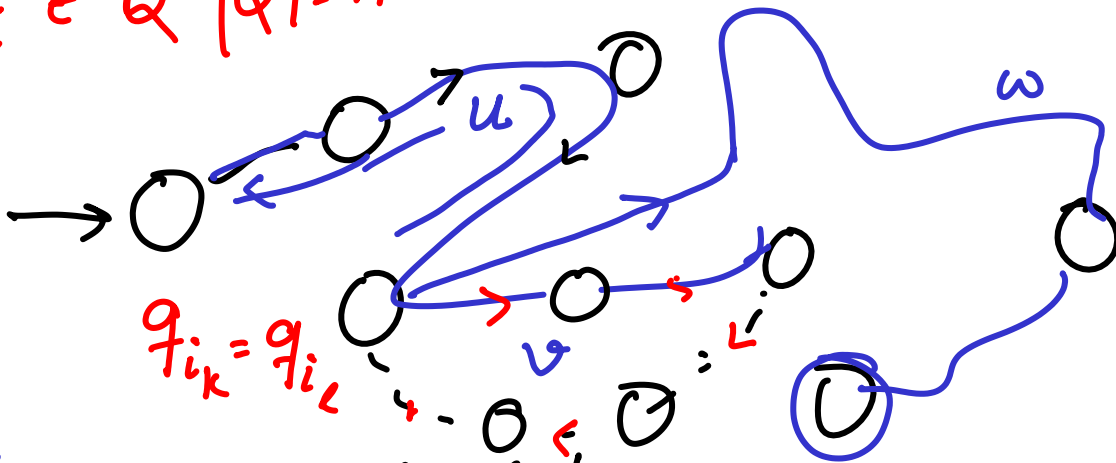
Claim If L is regular, so is L^*

Suppose we have a DFA M for a regular language L with n states. Given a string η (including initial state) $w_1 w_2 w_3 \dots w_n$ length n , we have sequence of state transitions $w_i \in \Sigma$

$$q_0 = q_{i_1} \xrightarrow{w_1} q_{i_2} \xrightarrow{w_2} q_{i_3} \dots q_{i_k} \dots q_{i_l} \dots q_{i_{n+1}}$$

$\delta(q_{i_1}, w_1) = q_{i_2}$

$q_{i_j} \in Q \quad |Q| = n$



$x = w_1 w_2 \dots w_n \in L$

s.t. $u v^i w \in L \quad i \geq 0$ s.t. $x \in L$

Pumping Lemma

Given any string x of length $\geq n$ (# states) we can partition $x = uvw$ s.t. $u v^i w \in L$ for $i \geq 0$. $|v| \geq 1$ $|u \cdot v| \leq n$