

RISC-V EMULATOR:-

RISCEMU is a command-line-based emulator. It can be used to run the RISC-V program.

This application supports all the RV32I and RV64I instructions. It also supports “M” standard extension for integer multiplication and division.

This interpreter supports the following instructions:

Computational Instruction:-

Arithmetic: add, addw, addi, addiw, sub, subw, mul, mulw, mulh, mulhu, mulhsu, lui

Logical: or, ori, and, andi, xor, xori

Compare: slt, slti, sltu, sltui

Shift: sll, sllw, slli, slliw, srl, srlw, srli, srliw, sra, saw, srai, sraiw

DataTransfer Instruction:- sb, sh, sw, sd, lb, lh, lw, ld, lbu, lbh, lbw

Branch Instructions:- beq, bne, bge, blt, bgeu, bltu, jal, jalr

RISC-V Architecture:-

General-purpose registers x1 through x31 are available for use without any restrictions or special functions assigned by the processor hardware. Register x0 is hardwired to return zero when it is read and will discard any value written to it. There are no processor flag registers in the RISC-V ISA. Some operations that modify flags in other processor architectures instead store their results in a RISC-V register. For example, the signed (slt) and unsigned (sltu) RISC-V comparison instructions subtract two

operands and set a destination register to 0 or 1 depending on the sign of the result. A subsequent conditional branch instruction uses the value in that register to determine which code path to take.

Most of the base ISA computational instructions use a three-operand format, in which the first operand is the destination register, the second operand is a source register, and the third operand is either a second source register or an immediate value. This is an example of three-operand instruction:

```
add x1, x2, x3
```

To avoid introducing instructions that are not strictly necessary, many instructions take on extra duties that are performed by dedicated instructions in other processor architectures. For example, RISC-V has no instruction that simply moves one register to another. Instead, a RISC-V addition instruction adds a source register and an immediate value of zero and stores the result in a destination register, producing the same result. The instruction to transfer register x2 to x1 is therefore add x1, x2, x0 assigning the value $(x2 + 0)$ to x1.

RISC-V Base Instruction Set:-

RISC-V base instructions fall into the categories of computational instructions, control flow instructions, and memory access instructions.

Computational Instructions:-

- **Arithmetic Instruction** (add, addi, sub):-Perform addition and subtraction. The immediate value in the addi instruction is a 12-bit signed value. The sub instruction subtracts the second source operand from the first. There

is no subi instruction because addi can add a negative immediate value.

- **Shift Instructions(sll, slli, srl, srli, sra, srai):-** Perform logical left and right shifts (sll and srl), and arithmetic right shifts (sra). Logical shifts insert zero bits into vacated locations. Arithmetic right shifts replicate the sign bit into vacated locations. The number of bit positions to shift is taken from the lowest 5 bits of the second source register or the 5-bit immediate value.
- **Logical Instruction(and, andi, or, ori, xor, xori):-** Perform the indicated bitwise operation on the two source operands. Immediate operands are 12 bits.
- **Compare Instructions(slt, slti, sltu, sltui):-** The set if less than instructions set the destination register to 1 if the first source operand is less than the second source operand: This comparison is in terms of two's complement (slt) or unsigned (sltu) operands. Immediate operand values are 12 bits.
- **lui:-** Load upper immediate. This instruction loads bits 12-31 of the destination register with a 20-bit immediate value. Setting a register to an arbitrary 32-bit immediate value requires two instructions: First, lui sets bits 12-31 to the upper 20 bits of the value. Then addi adds in the lower 12 bits to form the complete 32-bit result. lui has two operands: the destination register and the immediate value

Control Flow Instructions:-

The conditional branching instructions perform comparisons between two registers and, based on the result, may transfer control within the range of a signed 12-bit address offset from the current PC. Two unconditional jump instructions are available,

one of which (jalr) provides access to the entire 32-bit address range.

- **beq, bne, blt, bltu, bge, bgeu**:- Branch if equal (beq), not equal (bne), less than (blt), less than unsigned (bltu), greater or equal (bge), or greater or equal, unsigned (bgeu). These instructions perform the designated comparison between two registers and, if the condition is satisfied, transfer control to the address offset provided in the 12-bit signed immediate value.
- **jal**:- Jump and link. Transfer control to the PC-relative address provided in the 20-bit signed immediate value and store the address of the next instruction (the return address) in the destination register.
- **jalr**:- Jump and link, register. Compute the target address as the sum of the source register and a signed 12-bit immediate value, then jump to that address and store the address of the next instruction in the destination register. When preceded by the auipc instruction, the jalr instruction can perform a PC-relative jump anywhere in the 32-bit address space.

Memory Access Instructions:-

The memory access instructions transfer data between a register and a memory location. The first operand is the register to be loaded or stored. The second is a register containing a memory address. A signed 12-bit immediate value is added to the address in the register to produce the final address for the load or store.

The load instructions perform sign extension for signed values or zero extension for unsigned values. The sign or zero extension operation fills in all 32 bits in the destination register when a

smaller data value (a byte or halfword) is loaded. Unsigned loads are specified by a trailing *u* in the mnemonic.

- **lb, lbu, lh, lhu, lw**:- Load an 8-bit (byte) (lb), a 16-bit (halfword) (lh) or 32-bit (word) (lw) into the destination register. For byte and halfword loads, the instruction will either sign-extend (lb and lh) or zero-extend (lbu and lhu) to fill the 32-bit destination register. For example, the instruction `lw x1, 16(x2)` loads the word at the address $(x2 + 16)$ into register x1.
- **sb, sh, sw**:- Store a byte (sb), halfword (sh), or word (sw) to a memory location matching the size of the data value.

Pseudo Instructions:-

The RISC-V architecture has a truly reduced instruction set, lacking several types of instructions present in the instruction sets of other processor architectures. The functions of many of those more familiar instructions can be performed with RISC-V instructions, though perhaps not in an immediately intuitive manner.

The RISC-V assembler supports several pseudo-instructions, each of which translates to one or more RISC-V instructions providing a type of functionality one might expect in a general-purpose processor instruction set. The following table presents a few of the most useful RISC-V pseudo-instructions:

PSEUDO-INSTRUCTION	RISC-V INSTRUCTION	FUNCTION
Nop	addi x0, x0, 0	No operation
mv rd,rs	addi rd, rs, 0	Copy rs to rd
not rd, rs	ori rd, rs, -1	rd = NOT rs
neg rd, rs	sub rd, x0, rs	rd = -rs
j offset	jal x0, offset	Unconditional jump
jal offset	jal x1, offset	Near function call (20-bit offset)
call offset	auipc x1, offset[31:12] + offset[11] jalr x1, offset[11:0](x1)	Far function call (32-bit offset)
ret	jalr x0, 0(x1)	Return from function
beqz rs, offset	beq rs, x0, offset	Branch if equal to zero
bgez rs, offset	bge rx, x0, offset	Branch if greater than or equal to zero
bltz rs, offset	blt rs, x0, offset	Branch if less than zero
bgt rs, rt, offset	blt rt, rs, offset	Branch if greater than
ble rs, rt, offset	bge rt, rs, offset	Branch if less than or equal
li rd, immed	addi rd, x0, immed	Load 12-bit immediate

RISC-V Extension:-

The instruction set described in this section is named RV32I, which stands for the RISC-V 32-bit integer instruction set. Although the RV32I ISA provides a complete and useful

instruction set for many purposes, it lacks several functions and features available in other processors such as x86 and ARM.

The RISC-V extensions provide a mechanism for adding capabilities to the base instruction set in an incremental and compatible manner. Implementors of RISC-V processors can selectively include extensions in a processor design to optimize trade-offs between chip size, system capability, and performance.

M Extension:-

The RISC-V M extension adds integer multiplication and division functionality to the base RV32I instruction set. The following instructions are available in this extension:

- **mul**:- Multiply two 32-bit registers and store the lower 32 bits of the result in the destination register.
- **mulh, mulhu, mulhsu**:- Multiply two 32-bit registers and store the upper 32 bits of the result in the destination register. Treat the multiplicands as both signed (mulh), both unsigned (mulhu), or signed rs1 times unsigned rs2 (mulhsu). rs1 is the first source register in the instruction and rs2 is the second.
- **div, divu** :- Perform division of two 32-bit registers, rounding the result toward zero, on signed (div) or unsigned (divu) operands.
- **rem, remu**:- Return the remainder corresponding to the result of a div or divu instruction on the operands.

Division by zero does not raise an exception. To detect division by zero, the code should test the divisor and branch to an appropriate handler if it is zero.

64-bit RISC-V:-

The RISC-V introduction to this point has discussed the 32-bit RV32I architecture and instruction set, with extensions. The RV64I instruction set extends RV32I to a 64-bit architecture. As in RV32I, instructions are 32-bits wide. The RV64I instruction set is almost entirely the same as RV32I except for a few significant differences:

- All of the integer registers are widened to 64 bits.
- Addresses are widened to 64 bits.
- Bit shift counts in instruction opcodes increase in size from 5 to 6 bits.
- Several new instructions are defined to operate on 32-bit values in a manner equivalent to RV32I. These instructions are necessary because most instructions in RV64I operate on 64-bit values and there are many situations in which it is necessary to operate efficiently on 32-bit values. These word-oriented instructions have an opcode mnemonic suffix of w. The w-suffix instructions produce signed 32-bit results. These 32-bit values are sign-extended (even if they are unsigned values) to fill the 64-bit destination register. In other words, bit 31 of each result is copied into bits 32-63.

The following new instructions are defined in RV64I:

- **addw, addiw, subw, sllw, slliw, srlw, srliw, sraw, sraiw**:- These instructions perform equivalently to the RV32I instruction with the same mnemonic, minus the w suffix. They work with 32-bit operands and produce 32-bit results. The result is sign-extended to 64 bits.

- **ld, sd**:- Load and store a 64-bit doubleword. These are the 64-bit versions of the lw and sw instructions in the RV32I instruction set.

The remaining RV32I instructions perform the same functions in RV64I, except addresses and registers are 64 bits in length. The same opcodes, both in assembly source code and binary machine code, are used in both instruction sets.

Design of Emulator:-

This emulator is designed in multiple file formats. Here we have nine .c files that contains the main() function and nine header files that contain definitions of all useful functions.

.c File:- MAINFILE.c, Execution.h, Arithmetic.h, BranchInst.c, Compare.c, DataTransfer.c, Logical.c, Register.c, Shift.c

.h File:- Global_Variable.h, Execution.h, Register.h, Arithmetic.h, Logical.h, BranchInst.h, DataTransfer.h, Shift.h and Compare.h

Global Variable.h:- All the global variables are externally declared here.

Excecution.h:- print() and executeInstruction() functions are externally declared here.

Register.h:- This file contains external declaration of invalidInst(), Rtype(), Itype(), ItypeL(), Btype(), Utype(), Stype and dec() functions.

Arithmetic.h:- This header file contains external declaration of ADD(), ADDI(), ADDW(), ADDIW(), SUB(), SUBW(), MUL(), MULW(), MULH(), MULHU(), MULHSU(), DIV(), DIVW(), REM(), REMU() and LUI() functions.

Logical.h:- This header file contains external declaration of OR(), ORI(), AND(), ANDI(), XOR() and XORI() functions.

BranchInst():- This header file contains external declaration of getPcforLabel(), BEQ(), BNE(), BGE(), BGEU(), BLT(), BLTU(), JAL() and JALR() functions.

DataTransfer.h:- This header file contains external declaration of STOREB(), STOREH(), STOREW(), STORED(), LOADB(), LOADBU(), LOADH(), LOADHU(), LOADW(), LOADWU(), and LOADD() functions.

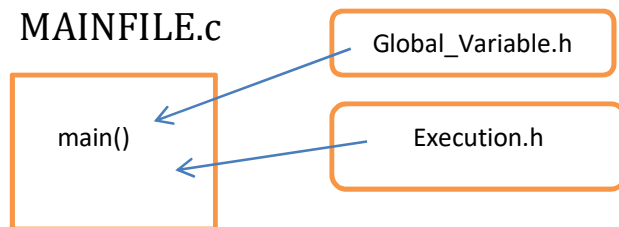
Shift.h:- This header file contains external declaration of SLL(), SLLW(), SLLI(), SLLIW(), SRL(), SRLW(), SRLI(), SRLIW(), SRA(), SRAI(), SRAW() and SRAIW() functions.

Compare.h:- This header file contains external declaration of SLT(), SLTU(), SLTI() and SLTIU() functions.

MAINFILE.c:-

This file contains a main() function that reads your complete file that is given to its argument and stores all instructions and labels in an array of struct data types. It keeps track of which instruction is executing.

It includes the Global_Variable.h and Execution.h file and for every new instruction, it transfers control to the Execution.c file.



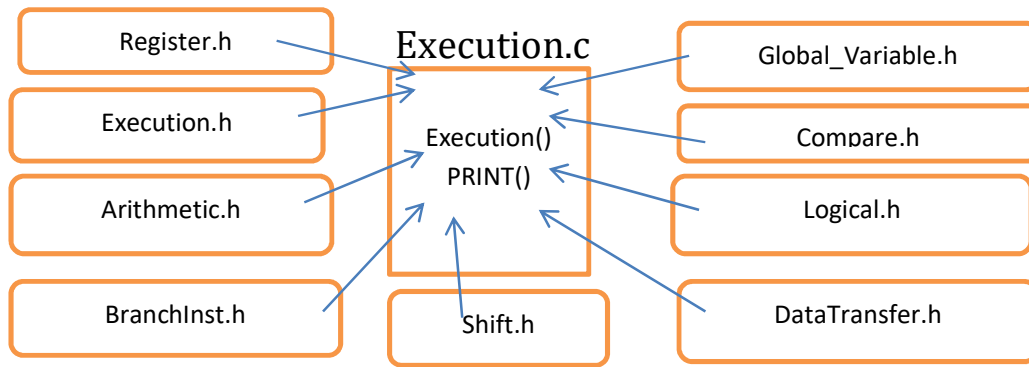
Execution.h:-

It includes Global_Variable.h, Register.h, Arithmetic.h, Logical.h, Shift.h, DataTransfer.h, Compare.h, BranchInst.h and Execution.h header files. This file stores definition of the executeInstruction() and print() functions.

executeInstruction():-This file stores current instruction in a string and identifies which instruction is going to be executed. If it is unable to recognize which instruction it is. Then, it prints the line number where it failed and terminates the program.

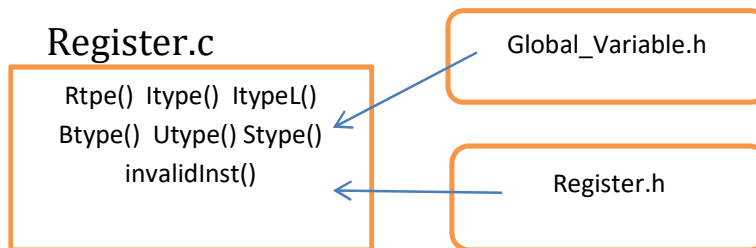
If it recognizes instruction it transfers control to the specific function that is present in the specific .c file.

print():-This function is called when the print label is found. It prints values stored in specific registers.



Register.c:-

This file stores the definition of `Rtype()`, `Itype()`, `ItypeL()`, `Stype()`, `Btype()`, `Utype()` and `invalidInstruction()` functions.



Rtype(): This function is called when the instruction format is `op rd, rs1, rs2`. This function identifies which register is used as the destination register and which as the source registers. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

Itype():-This function is called when the instruction format is `op rd, rs1, imm`. This function recognizes which register is used as the destination register and which is the source register. It also finds the decimal value of immediate. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

ItypeL():-This function is called when the instruction format is `op rd, offset(rs1)`. This function identifies which register is used as the destination register and which is the source register. It finds

the decimal value of the offset. It takes a string in which instruction is stored and the index which points to the destination register as arguments. Its return type is void.

Stype():-This function is called when the instruction format is op rs1, offset(rs2). This function recognizes which registers are used as source registers. It finds the decimal value of the offset. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

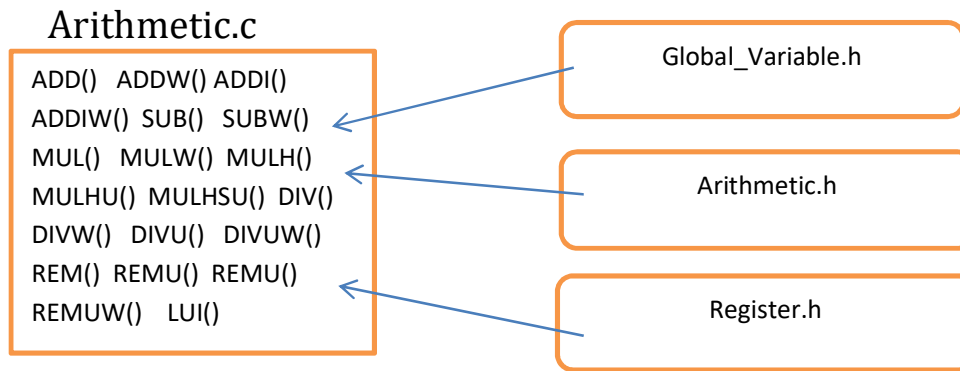
Btype():-This function is called when the instruction format is op rs1, rs2, imm/label. This function recognizes which registers are used as source registers. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is int. It returns the index where it is currently pointing in the string.

Utype():-This function is called when the instruction format is op rd imm/label. This function recognizes which register is used as the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is int. It returns the index where it is currently pointing in the string.

invalidInstruction():- This function is called when any undefined behavior is encountered.

Arithmetic.c:-

This file stores definition of ADD(), ADDI(), ADDW(), ADDIW(), SUB(), SUBW(), MUL(), MULW(), MULH(), MULHU(), MULHSU(), DIV(), DIVW(), DIVU(), DIVUW(), REM(), REMU(), REMW(), and REMUW() functions.



ADDQ:-This function is called when instruction is `add rd, rs1, rs2`. It performs addition on values present in source registers and stores the result in the destination register. It takes a string in which instruction is stored and the index that points to the destination register as arguments. Its return type is void.

ADDIQ:-This function is called when instruction is `addi rd, rs1, imm`. It performs addition on values present in the source register and immediate and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

ADDWQ:-This function is called when instruction is `addw rd, rs1, rs2`. It performs addition on lower 32-bit values present in source registers and stores its sign-extended representation in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

ADDIWQ:-This function is called when instruction is `addiw rd, rs1, imm`. It performs addition on lower 32-bit values present in the source register and immediate stores its sign-extended representation in the destination register. It takes a string in which instruction is stored and index which points destination register as arguments. Its return type is void.

SUBQ:-This function is called when instruction is sub rd, rs1, rs2. It subtracts values present in rs2 from values present in rs1 and stores the result in the destination register. It takes a string in which instruction is stored and the index that points to the destination register as arguments. Its return type is void.

SUBWQ:-This function is called when instruction is sub rd, rs1, rs2. It subtracts lower 32-bit values present in rs2 from lower 32-bit values present in rs1 and stores the result in its sign-extended representation in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

MULQ:-This function is called when instruction is mul rd, rs1, rs2. It performs multiplication on values present in source registers and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

MULWQ:-This function is called when instruction is mulw rd, rs1, rs2. It performs multiplication on lower 32-bit values present in source registers and stores its sign-extended representation in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

MULHQ:-This function is called when instruction is mulh, rs1, rs2. It performs multiplication on values present in source registers and stores its upper 64-bit in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

MULHUQ:-This function is called when instruction is mulhu, rs1, rs2. It performs multiplication on unsigned values present in

source registers and stores its upper 64-bit in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

MULHSU():-This function is called when instruction is mulhsu, rs1, rs2. It performs multiplication on signed values present in the first source register and unsigned values present in the second source register and stores its upper 64-bit in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

DIVQ():-This function is called when instruction is div rd, rs1, rs2. It divides the value present in rs1 by the value present in rs2 and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

DIVW():-This function is called when instruction is divw rd, rs1, rs2. It divides the lower 32-bit value present in rs1 by the lower 32-bit value present in rs2 and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

DIVU():-This function is called when instruction is divu rd, rs1, rs2. It divides the unsigned value present in rs1 by the unsigned value present in rs2 and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

DIVUWQ:-This function is called when instruction is divuw rd, rs1, rs2. It divides the lower unsigned 32-bit value present in rs1 by the lower unsigned 32-bit value present in rs2 and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

REMQ:-This function is called when instruction is rem rd, rs1, rs2. It divides the value present in rs1 by the value present in rs2 and stores the remainder in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

REMWQ:-This function is called when instruction is remw rd, rs1, rs2. It divides the lower 32-bit value present in rs1 by the lower 32-bit value present in rs2 and stores the remainder in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

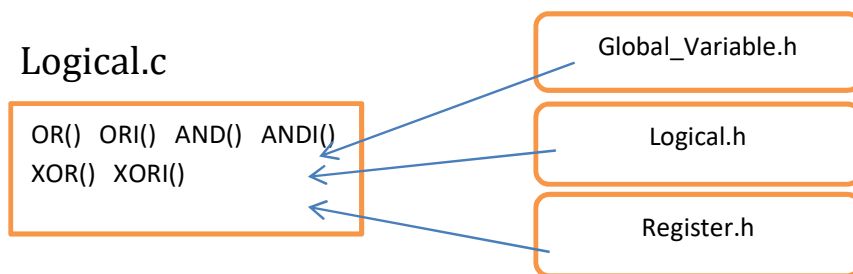
REMUQ:-This function is called when instruction is remu rd, rs1, rs2. It divides the unsigned value present in rs1 by the unsigned value present in rs2 and stores the remainder in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

REMUWQ:-This function is called when instruction is remuw rd, rs1, rs2. It divides the lower unsigned 32-bit value present in rs1 by the lower unsigned 32-bit value present in rs2 and stores the remainder in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

LUI():-This function is called when instruction is lui rd, imm. It stores 12-31 bits of destination register with 20-bit immediate. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

Logical.c:-

This file stores definition of OR(), ORI(), AND(), ANDI(), XOR() and XORI() functions.



OR():-This function is called when instruction is or rd, rs1, rs2. It performs or operation on values present in source registers and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

ORI():-This function is called when instruction is ori rd, rs1, imm. It performs or operation on values present in the source register and immediate and stores result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

AND():-This function is called when instruction is and rd, rs1, rs2. It performs and operation on values present in source registers and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

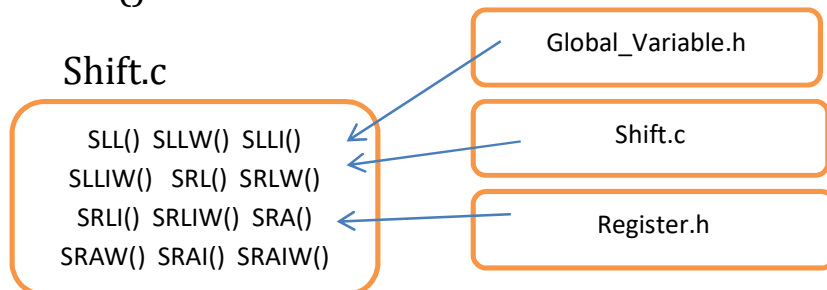
ANDI():-This function is called when instruction is `andi rd, rs1, imm`. It performs and operation on values present in the source register and immediate and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

XOR():-This function is called when instruction is `xor rd, rs1, rs2`. It performs xor operation on values present in source registers and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

XORI():-This function is called when the instruction is `xori rd, rs1, imm`. It performs xor operation on values present in the source register and immediate and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

Shift.c:-

This file stores definition of `SLL()`, `SLLW()`, `SLLI()`, `SLLIW()`, `SRA()`, `SRAW()`, `SRAI()`, `SRAIW()`, `SRL()`, `SRLW()`, `SRLI()` and `SRLIW()` functions.



SLL():-This function is called when instruction is `sll rd, rs1, rs2`. It shifts left the value present in `rs1` by the value present in `rs2` and stores the result in the destination register. It takes a string in

which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SLLW():-This function is called when instruction is sllw rd, rs1, rs2. It shifts left lower 32-bit value present in rs1 by lower 32-bit value present in rs2 and stores the result in the destination register. It takes a string in which instruction is stored and index which points destination register as arguments. Its return type is void.

SLLI():-This function is called when instruction is slli rd, rs1, imm. It shifts left the value present in rs1 by the value of immediate and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SLLIW():-This function is called when instruction is slliw rd, rs1, rs2. It shifts left the lower 32-bit value present in rs1 by the value immediate and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SRA():-This function is called when instruction is sra rd, rs1, rs2. It shifts right the value present in rs1 by the value present in rs2 arithmetically and stores the result in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SRAW():-This function is called when instruction is saw rd, rs1, rs2. It shifts right the lower 32-bit value present in rs1 by the lower 32-bit value present in rs2 arithmetically and stores it in the destination register. It takes a string in which instruction is

stored and an index that points to the destination register as arguments. Its return type is void.

SRAIQ:-This function is called when instruction is srai rd, rs1, imm. It shifts right the value present in rs1 by the value of immediate arithmetically and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SRAIWQ:-This function is called when instruction is sraiw rd, rs1, rs2. It shifts right arithmetically the lower 32-bit value present in rs1 by the value of immediate and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SRLQ:-This function is called when instruction is srl rd, rs1, rs2. It shifts right the value present in rs1 by the value present in rs2 logically and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SRLWQ:-This function is called when instruction is srlw rd, rs1, rs2. It shifts right the lower 32-bit value present in rs1 by the lower 32-bit value present in rs2 logically and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

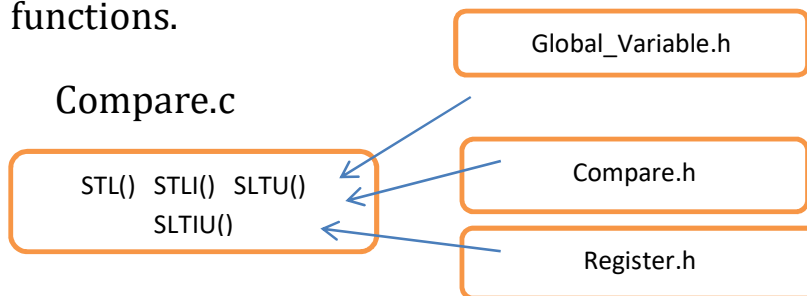
SRLIQ:-This function is called when instruction is srli rd, rs1, imm. It shifts the right value present in rs1 by the value of immediate logically and stores it in the destination register. It takes a string in which instruction is stored and an index that

points to the destination register as arguments. Its return type is void.

SRLIW():-This function is called when instruction is srlw rd, rs1, rs2. It shifts right the lower 32-bit value present in rs1 by the value of immediate logically and stores it in the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

Compare.c:-

This file stores definition of SLT(), SLTI(), SLTU() and SLTIU() functions.



SLTQ:-This function is called when instruction is slt rd, rs1, rs2. It compares the content of both source registers. If the content of rs1 is less than the content of rs2 it sets the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SLTIQ:-This function is called when instruction is slti rd, rs1, imm. It compares the content of rs1 and immediate. If the content of rs1 is less than immediate it sets the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

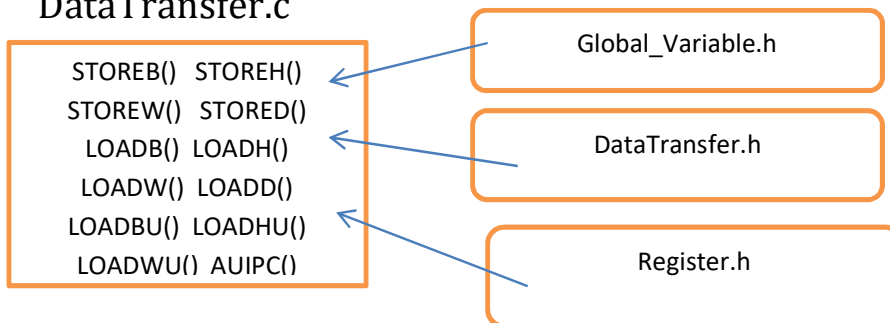
SLTU():-This function is called when instruction is sltu rd, rs1, rs2. It compares unsigned values stored in both source registers. If the unsigned value of rs1 is less than the unsigned value of rs2 it sets the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

SLTUI():-This function is called when instruction is sltui rd, rs1, rs2. It compares unsigned values of rs1 and immediate. If unsigned content of rs1 is less than immediate it sets the destination register. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

DataTransfer.c:-

This file stores definition of STOREB(), STOREH(), STOREW(), STORED(), LOADB(), LOADH(), LOADW(), LOADD(), LOADBU(), LOADHU() and LOADWU() functions.

DataTransfer.c



STOREB():-This function is called when instruction is sb rs1, offset(rs2). It stores lower 8-bit of rs1 to the memory location [rs2]+offset. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

STOREH():-This function is called when instruction is sh rs1, offset(rs2). It stores lower 16-bit of rs1 to the memory location [rs2]+offset. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

STOREW():-This function is called when instruction is sw rs1, offset(rs2). It stores the lower 32-bit of rs1 to the memory location [rs2]+offset. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

STORED():-This function is called when instruction is sd rs1, offset(rs2). It stores the content of rs1 to the memory location [rs2]+offset. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

LOADB():-This function is called when instruction is lb rd, offset(rs1). It loads a byte from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

LOADH():-This function is called when instruction is lh rd, offset(rs1). It loads a halfword from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

LOADW():-This function is called when instruction is lw rd, offset(rs1). It loads a word from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an

index that points to the destination register as arguments. Its return type is void.

LOADD():-This function is called when instruction is lb rd, offset(rs1). It loads a doubleword from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

LOADBU():-This function is called when instruction is lbu rd, offset(rs1). It loads the unsigned value of byte from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

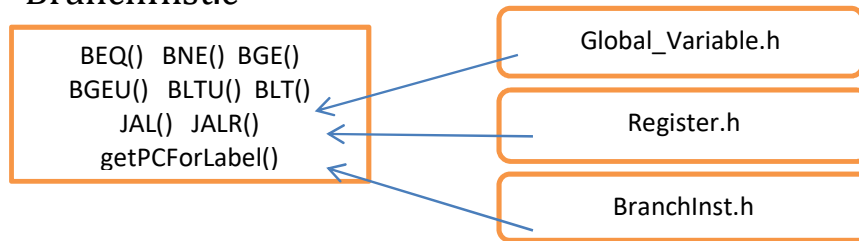
LOADHU():-This function is called when instruction is lhu rd, offset(rs1). It loads the unsigned value of halfword from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

LOADWU():-This function is called when instruction is lwu rd, offset(rs1). It loads the unsigned value of the word from memory location [rs1]+offset to register rd. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

BranchInst.c:-

This file stores definition of BEQ(), BNE(), BGE(), BGEU(), BLT(), BLTU(), JAL() and JALR() functions.

BranchInst.c



BEQ():-This function is called when instruction is `beq rs1, rs2, imm/label`. It compares the content of `rs1` and `rs2` if both are equal it sets `pc` to respective label or `pc+immediate`. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

BNE():-This function is called when instruction is `bne rs1, rs2, imm/label`. It compares the content of `rs1` and `rs2` if both are not equal it sets `pc` to respective labels or `pc+immediate`. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

BGE():-This function is called when instruction is `bge rs1, rs2, imm/label`. It compares the content of `rs1` and `rs2` if the content of `rs1` is greater than or equal to the content of `rs2` it sets `pc` to respective label or `pc+immediate`. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

BGEU():-This function is called when instruction is `bgeu rs1, rs2, imm/label`. It compares the content of `rs1` and `rs2` if the unsigned value of `rs1` is greater than or equal to the unsigned value of `rs2` it sets `pc` to the respective label or `pc+immediate`. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

BLT():-This function is called when instruction is `blt rs1, rs2, imm/label`. It compares the content of `rs1` and `rs2` if the content of

rs1 is less than the content of rs2 it sets pc to the respective label or pc+immediate. It takes a string in which instruction is stored and an index that points to the first source register as arguments. Its return type is void.

JAL():-This function is called when instruction is jal rd, imm/label. It stores the location of the next instruction in rd sets pc to respective label or pc+immediate. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

JALR():-This function is called when instruction is jalr rd, offset(rs1). It stores the location of the next instruction in rd sets pc to ([rs1]+offset)/4. It takes a string in which instruction is stored and an index that points to the destination register as arguments. Its return type is void.

Execution of Arithmetic Instruction:-

Different formats of Arithmetic instruction

op rd, rs1, rs2
op rd, rs1, imm
op rd, imm

When instruction is in format op rd, rs1, rs2.

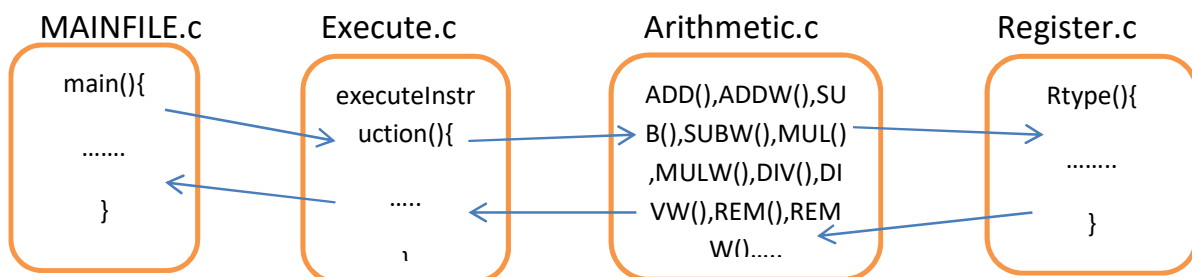
add rd, rs1, rs2
addw rd, rs1, rs2
sub rd, rs1, rs2
subw rd, rs1, rs2
mul rd, rs1, rs2
mulw rd, rs1, rs2
div rd, rs1, rs2

```

divu rd, rs1, rs2
divw rd, rs1, rs2
divuw rd, rs1, rs2
rem rd, rs1, rs2
remw rd, rs1, rs2
remu rd, rs1, rs2
remuw rd, rs1, rs2

```

In the process of execution of a program when pc points one of these instruction, control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Arithmetic.c file where control is transferred to Rtype() function(resides in Register.c file). Rtype() function identifies which register is used as destination register and which are source registers and after that, it returns control to Arithmetic.c. Here, the appropriate operation is performed on the content of the source registers and the result is stored in the destination register. Now, control is returned to executeInstruction() function and from here control is returned to main() function where pc is increased to point to next instruction



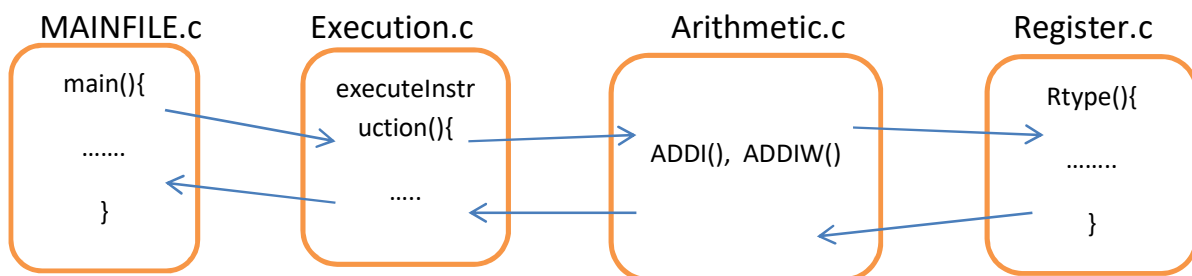
Control Flow of Execution of Arithmetic Instruction(op rd, rs1, rs2)

When instruction is in format op rd, rs1, imm.

addi rd, rs1, imm

addiw rd, rs1, imm

In the process of execution of a program when pc points to one of these instruction controls is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Arithmetic.c file where control is transferred to Itype() function(resides in Register.c file). Itype() function recognizes which register is used as a destination register and which is a source register. It finds the value of immediate and also ensures that immediate is a 12-bit value and after that, it returns control to Arithmetic.c. Here, the appropriate operation is performed on the content of the source register and immediate and result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.

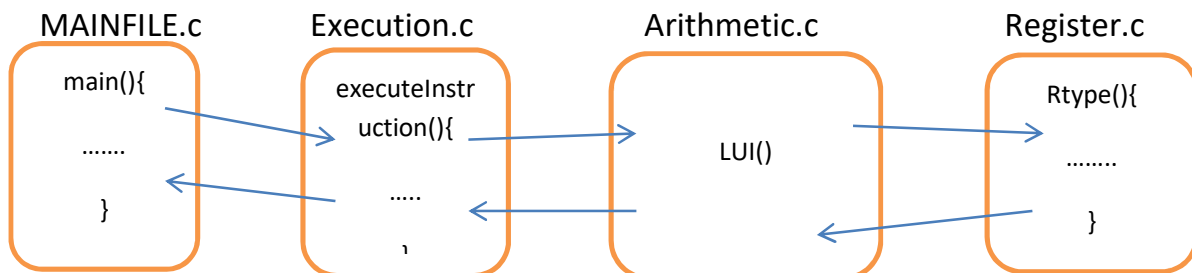


Control Flow of Execution of Arithmetic Instruction(op rd, rs1, imm)

When instruction is in format op rd, imm.

lui rd, imm

In the process of execution of a program when pc points this instruction, control is transferred from main() function(resides in MAINFILE.c) to executeInstruction()function(resides in Execution.c file). Where it recognizes which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Arithmetic.c file where control is transferred to Utype() function(resides in Register.c file). Utype() function identifies which register is used as the destination register and finds the value of immediate and also ensures that immediate is 20-bit and after that, it returns control to Arithmetic.c. Here, the appropriate operation is performed. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



Control Flow of Execution of Arithmetic Instruction(op rd, imm)

Execution of Logical Instruction:-

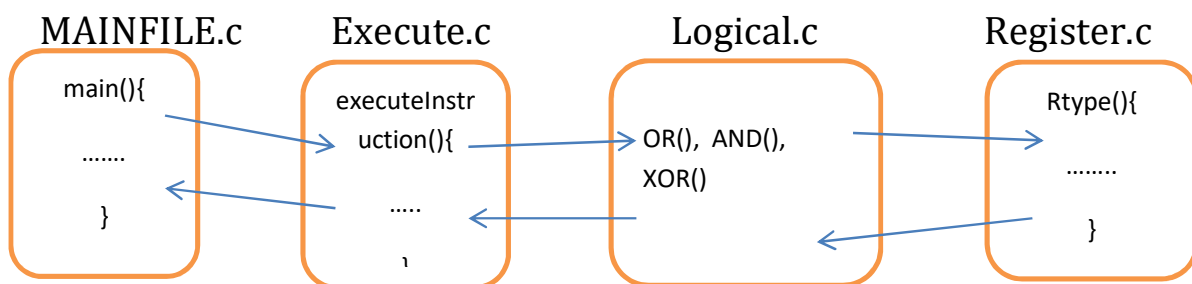
Different formats of Logical instruction

op rd, rs1, rs2
op rd, rs1, imm

When instruction is in format op rd, rs1, rs2.

or rd, rs1, rs2
and rd, rs1, rs2
xor rd, rs1, rs2

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction()function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Logical.c file where control is transferred to Rtype() function(resides in Register.c file). Rtype() function recognizes which register is used as destination register and which are source registers and after that, it returns control to Arithmetic.c. Here, the appropriate operation is performed on the content of the source register, and the result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.

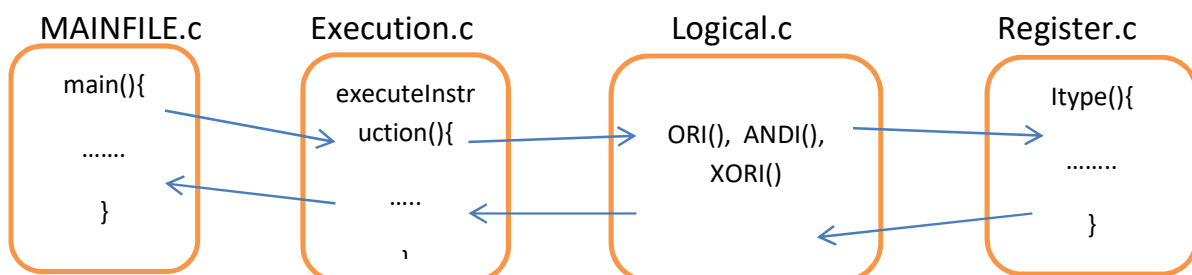


Control Flow of Execution of Logical Instruction(op rd, rs1, rs2)

When instruction is in format op rd, rs1, imm.

or rd, rs1, imm
andi rd, rs1, imm
xori rd, rs1, imm

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Logical.c file where control is transferred to Itype() function(resides in Register.c file). Itype() function recognizes which register is used as a destination register and which is a source register. It finds the value of immediate and also ensures that immediate is a 12-bit value and after that, it returns control to Logical.c. Here, an appropriate operation is performed on the content of the source register and the immediate and result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



Control Flow of Execution of Logical Instruction(op rd, rs1, imm)

Execution of Compare Instruction:-

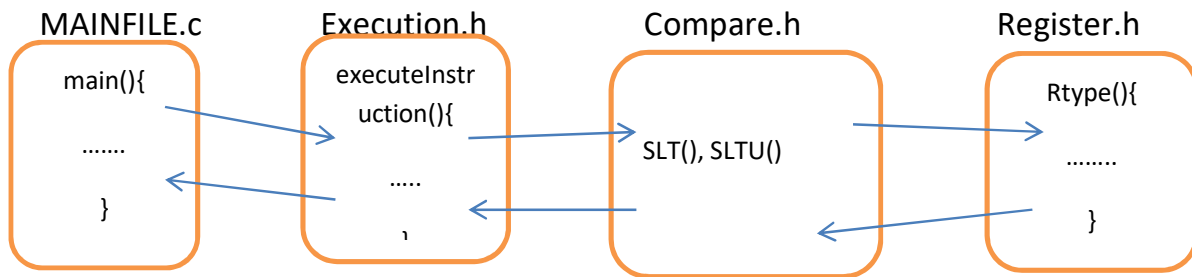
Different formats of Arithmetic instruction

op rd, rs1, rs2
op rd, rs1, imm

When instruction is in format op rd, rs1, rs2.

slt rd, rs1, rs2
sltu rd, rs1, rs2

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Compare.c file where control is transferred to Rtype() function(resides in Register.c file). Rtype() function recognizes which register is used as destination register and which are source registers and after that, it returns control to Compare.c. Here, an appropriate operation is performed on the content of the source register, and the result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



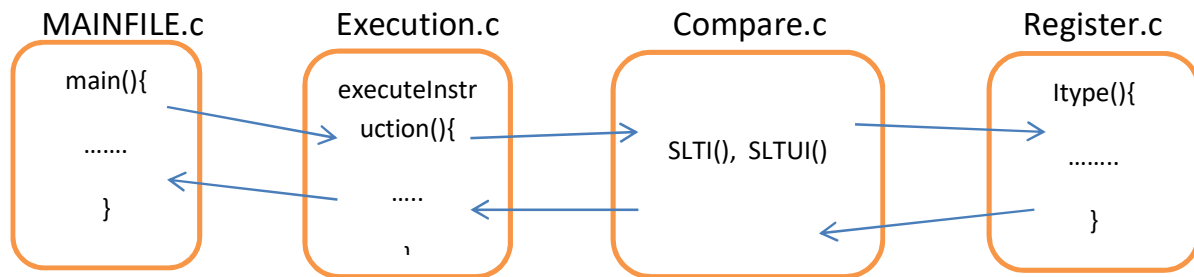
Control Flow of Execution of Compare Instruction(op rd, rs1, rs2)

When instruction is in format op rd, rs1, imm.

```

slti rd, rs1, imm
sltui rd, rs1, imm
  
```

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction()function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in Compare.c file where control is transferred to Itype() function(resides in Register.c file). Itype() function recognizes which register is used as the destination register and which is as source register finds the value of immediate and also ensures that immediate is a 12-bit value and after that, it returns control to Logical.c. Here, an appropriate operation is performed on the content of the source register and the immediate and result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



Control Flow of Execution of Compare Instruction(op rd, rs1, imm)

Execution of Shift Instruction:-

Different formats of Shift instruction:-

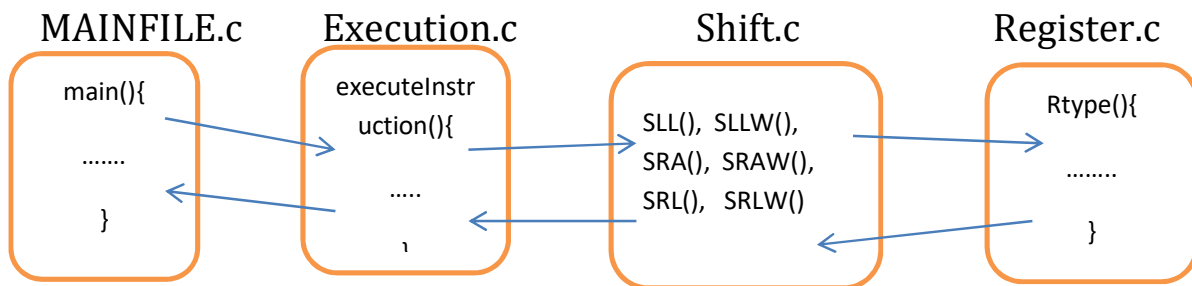
op rd, rs1, rs2
op rd, rs1, imm

When instruction is in format op rd, rs1, rs2.

sll rd, rs1, rs2
sllw rd, rs1, rs2
sra rd, rs1, rs2
sraw rd, rs1, rs2
srl rd, rs1, rs2
srlw rd, rs1, rs2

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in the Shift.c file where control is transferred to Rtype() function(resides in Register.c file). Rtype() function recognizes which register is used as the destination register and

which are source registers and after that, it returns control to Shift.c. Here, an appropriate operation is performed on the content of the source register and the result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



Control Flow of Execution of Shift Instruction(op rd, rs1, rs2)

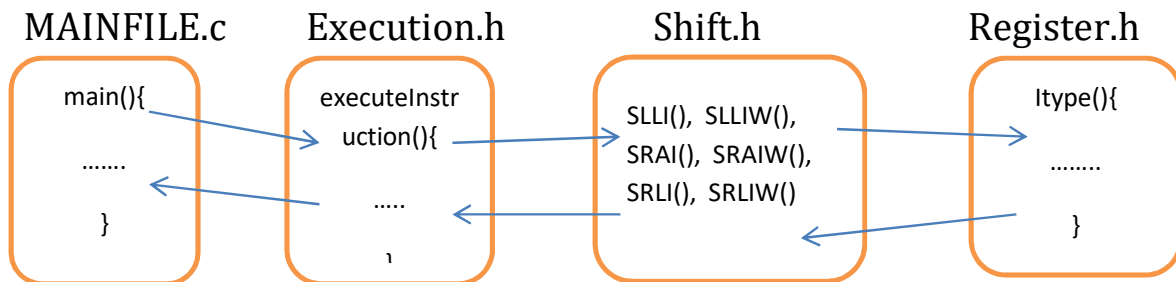
When instruction is in format op rd, rs1, imm.

```

slli rd, rs1, rs2
slliw rd, rs1, rs2
srai rd, rs1, rs2
sraiw rd, rs1, rs2
srli rd, rs1, rs2
srliw rd, rs1, rs2
  
```

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in the Shift.c file where control is transferred to

Itype() function(resides in Register.c file). Itype() function recognizes which register is used as the destination register and which is as the source register finds the value of immediate and also ensures that immediate is a 12-bit value and after that, it returns control to Shift.c. Here, an appropriate operation is performed on the content of the source register and the immediate and result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



Control Flow of Execution of Shift Instruction(op rd, rs1, imm)

Execution of DataTransfer Instruction:-

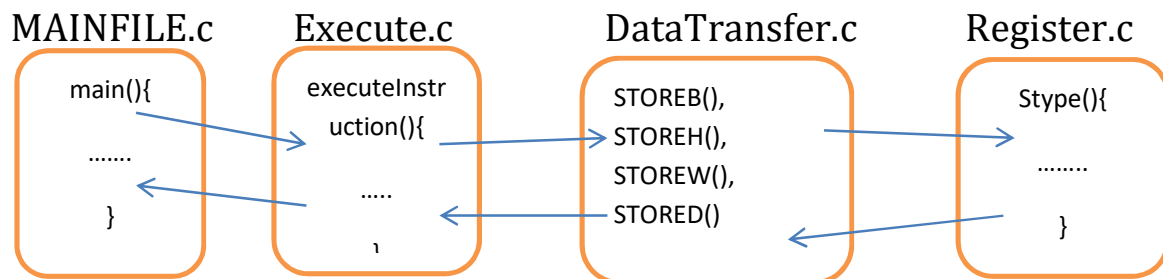
Different formats of DataTransfer instruction:-

op rs1, imm(rs2)
op rd, imm(rs1)

When instruction is in format op rs1, imm(rs2).

sb rs1, imm(rs2)
sh rs1, imm(rs2)
sw rs1, imm(rs2)
sd rs1, imm(rs2)

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to corresponding function that resides in the DataTransfer.c file where control is transferred to the Stype() function(resides in Register.c file). Stype() function recognizes which registers are used as source registers. It also finds value of immediate and also ensures that it can't be more than 12-bit value and after that it returns control to DataTransfer.c. Here, appropriate operation is performed. Now, control is returned to the executeInstruction() function and from here control is returned to main() function where pc is increased to point to the next instruction.



Control Flow of Execution of DataTransfer Instruction(op rs1,imm(rs2))

When instruction is in format `op rd, imm(rs1)`.

`lb rd, imm(rs1)`

`lh rd, imm(rs1)`

`lw rd, imm(rs1)`

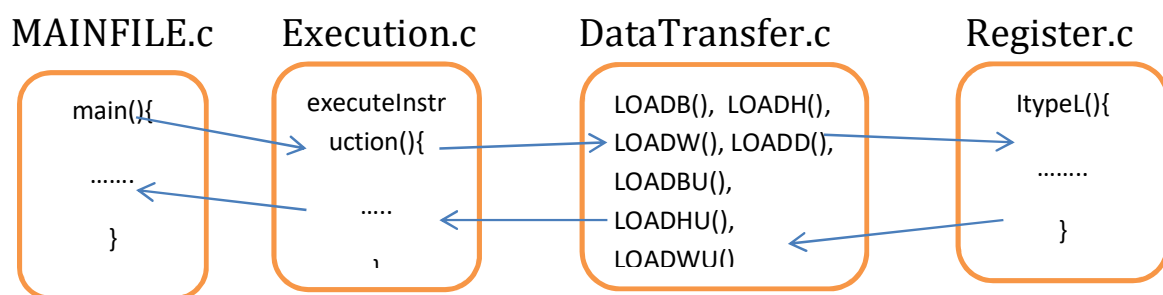
`ld rd, imm(rs1)`

```

lbu rd, imm(rs1)
lhu rd, imm(rs1)
lwu rd, imm(rs1)

```

In the process of execution of a program when pc points to one of these instruction control is transferred from main() function(resides in MAINFILE.c) to the executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from the executeInstruction() function to the corresponding function that resides in DataTransfer.c file where control is transferred to ltypeL() function(resides in Register.c file). ltypeL() function recognizes which register is used as a destination register and which is a source register. It finds the value of immediate and also ensures that immediate is a 12-bit value and after that, it returns control to DataTransfer.c. Here, the appropriate operation is performed on the content of the source register and the immediate and result is stored in the destination register. Now, control is returned to the executeInstruction() function and from here control is returned to the main() function where pc is increased to point to the next instruction.



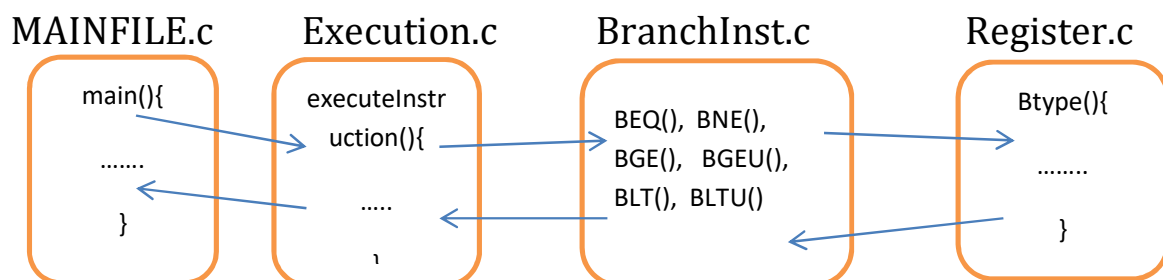
Control Flow of Execution of DataTransfer Instruction(op
rd, imm(rs1))

Execution of Branch Instruction:-

When instruction is in format op rs1, rs2, imm/label.

beq rs1, rs2, imm/label
bne rs1, rs2, imm/label
bge rs1, rs2, imm/label
bgeu rs1, rs2, imm/label
blt rs1, rs2, imm/label
bltu rs1, rs2, imm/label

In the process of execution of a program when pc points one of these instruction control is transferred from main() function(resides in MAINFILE.c) to executeInstruction() function(resides in Execution.c file). Where it identifies which operation is going to be performed and then transfers control from executeInstruction() function to the corresponding function that resides in BranchInst.c file where control is transferred to Btype() function(resides in Register.c file). Btype() function recognizes which registers are used as source registers. It also finds the value of immediate and ensures that it can't be more than 12-bit value and after that, it returns control to BranchInst.c. Here, the appropriate operation is performed. Now, control is returned to the executeInstruction() function and from here control is returned to main() function.



Control Flow of Execution of Branch Instruction(op rs1, rs2, imm/label)