

# Mining Specifications of Malicious Behavior

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# Why Understand Malicious Behavior?

- Forensics
  - Understand what a malware does
- Malware Detector
  - Move to behavior-based detectors
  - These need detectors need a high-level specification of malware

# Wide Spectrum of Detectors

- Static detectors:



- Dynamic/hybrid detectors, host IDS:



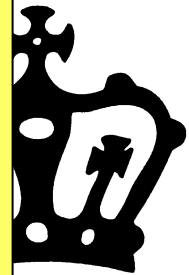
Semantic Web, Behavior-based software detection, Shadow Honeypots [Anagnostakis et al. 2005] [Kanda et al. 2006]

# Misuse Detection

Distinct techniques fundamentally similar...

Sample specification:

- Creates an email with itself attached, and
- Collects email addresses, and
- Sends emails



They **all** require high-quality **specifications of malicious behavior.**

# Key Definitions

**Variants** : New strains of viruses that **borrow code**, to varying degrees, directly from other known viruses.

*Source: Symantec Security Response Glossary*

**Virus family**: a set of variants with a **common code base**.

# Signature-Based Detection

```
lea    eax, [ebp+Data]
push   offset aServices_exe
push   eax
call   _strcat
pop    ecx
lea    eax, [ebp+Data]
pop    ecx
push   edi
push   eax
lea    eax, [ebp+ExistingFileNme]
push   eax
call   ds:CopyFileA
```

```
8D 85 D8 FE FF FF
68 78 8E 40 00
50
E8 69 06 00 00
59
8D 85 D8 FE FF FF
59
57
50
8D 85 D4 FD FF FF
50
FF 15 C0 60 40 00
```

Signature

- Signatures (aka scan-strings) are the most common malware detection mechanism.

# Current Signature Management

McAfee: release daily updates

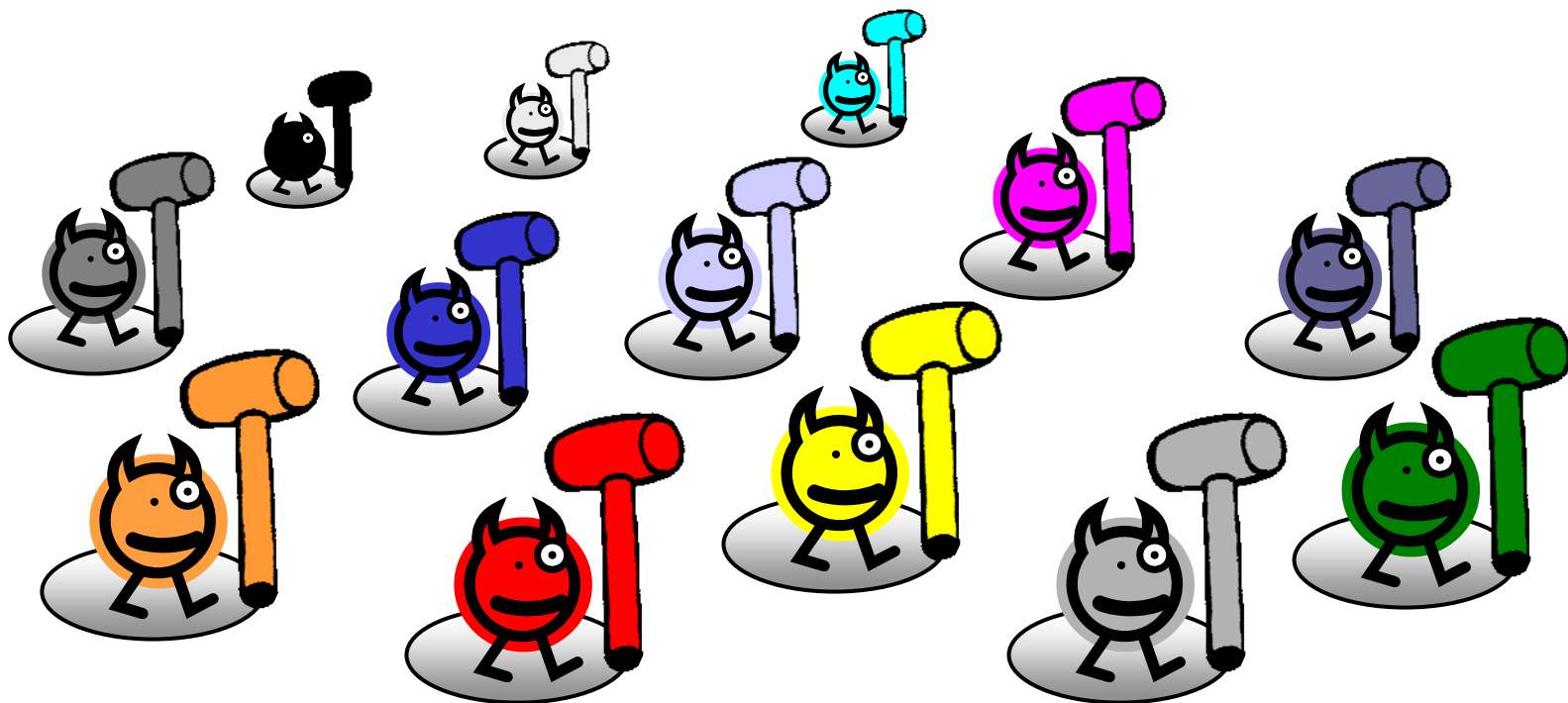
- Trying to move to hourly “beta” updates

DAT File #	Date	Threats Detected	New Threats Added	Threats Updated
4578	Sep. 09	147,382	22	188
4579	Sep. 12	147,828	27	231
4580	Sep. 13	148,000	11	236
4581	Sep. 14	148,368	42	140
4582	Sep. 15	148,721	16	203
4583	Sep. 16	149,050	18	117

*Source: McAfee DAT Readme*

# Signature Detection Does Not Scale

One signature for one malware instance.





# Goals for Better Detection

- Make the malware writer's job as hard as possible.
- Detect malware families, not individual malware instances.
- Move away from syntactic signatures.

# Threat Model

- Malware writers craft their programs so to avoid detection.

Two common **evasion techniques**:

- Program Obfuscation  
(Preserves malicious behavior)
- Program Evolution  
(Enhances malicious behavior)

# Obfuscations for Evasion

Nop insertion

Register renaming

Junk insertion

Instruction reordering

Encryption

Compression

Reversing of branch conditions

Equivalent instruction substitution

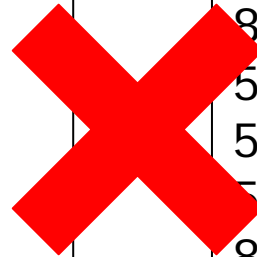
Basic block reordering

# Evasion Through Junk Insertion

```
lea    eax, [ebp+Data]
nop
push   offset aServices_exe
nop
nop
push   eax
call   _strcat
nop
nop
nop
pop    ecx
lea    eax, [ebp+Data]
pop    ecx
push   edi
push   eax
nop
lea    eax, [ebp+ExistingFileName]
push   eax
call   ds:CopyFileA
```

```
8D 85 D8 FE FF FF
68 78 8E 40 00
50
E8 69 06 00 00
59
8D 85 D8 FE FF FF
59
57
50
8D 85 D4 FD FF FF
50
FF 15 C0 60 40 00
```

Signature



# Evasion Through Reordering

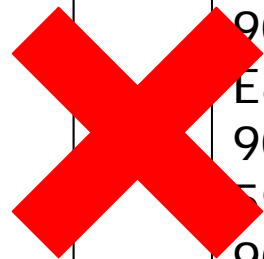
```
    lea    eax, [ebp+Data]
    jmp   label_one

label_two:
    lea    eax, [ebp+Data]
    ...
    push  eax
    call  ds:CopyFileA
    jmp   label_three

label_one:
    ...
    call  _strcat
    ...
    jmp   label_two

label_three: ...
```

```
8D 85 D8 FE FF FF
90*
68 78 8E 40 00
90*
50
90*
E8 69 06 00 00
90*
59
90*
.
.
.
90*
50
90*
FF 15 C0 60 40 00
```

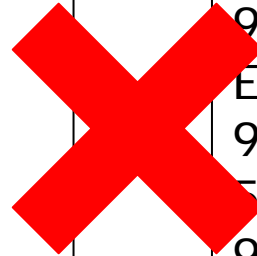


Regex Signature

# Evasion Through Encryption

```
    lea    esi, data_area
    mov    ecx, 37
again:
    xor    byte ptr [esi+ecx], 0x01
    loop  again
    jmp    data_area
    .
    .
    .
data_area:
    db    8C 84 D9 FF ...
    .
    .
    .
    db    FE 14 C1 61 ...
```

```
8D 85 D8 FE FF FF
90*
68 78 8E 40 00
90*
50
90*
E8 69 06 00 00
90*
59
90*
.
.
.
90*
50
90*
FF 15 C0 60 40 00
```



Regex Signature

# Evasion Through Evolution

- Malware writers are good at software engineering:
  - Modular designs
  - High-level languages
  - Sharing of exploits, payloads, and evasion techniques

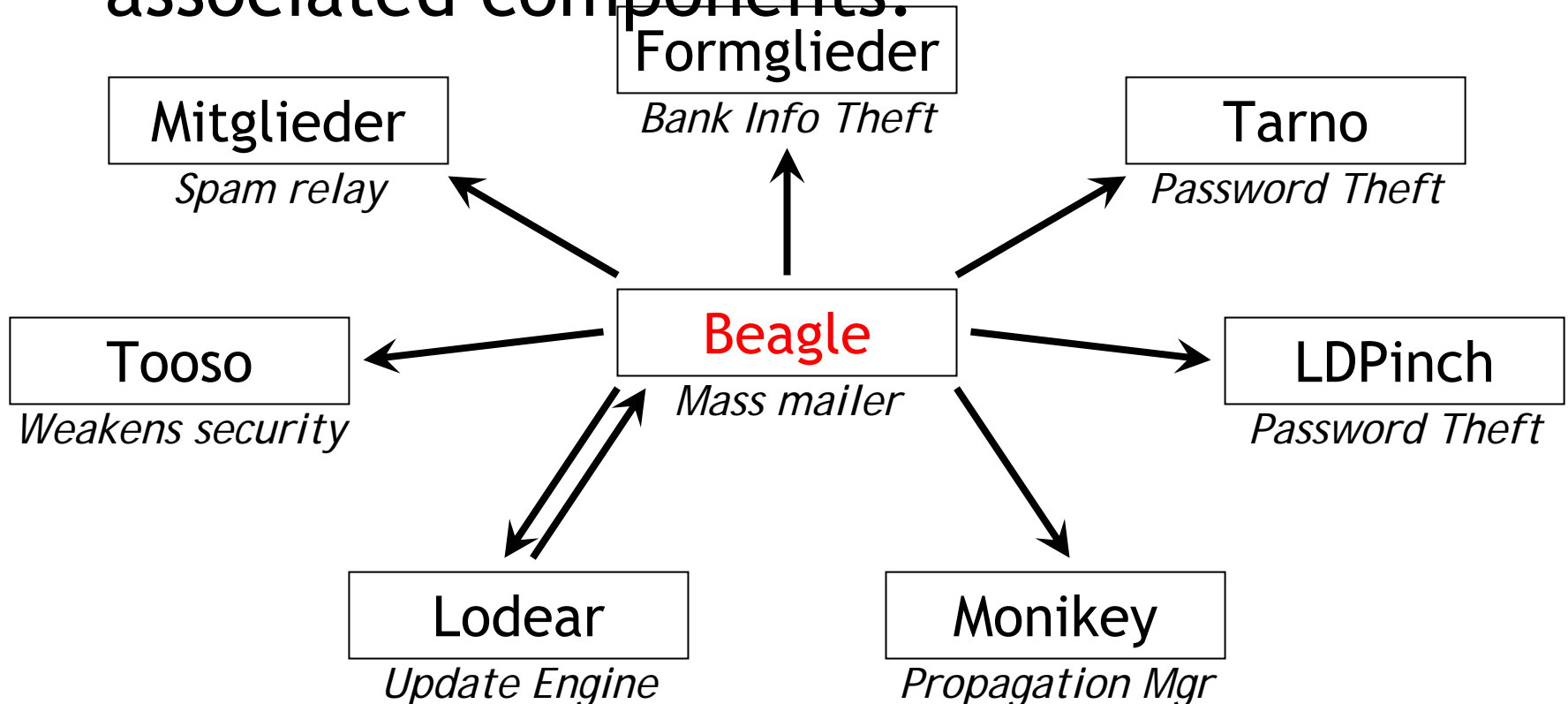
## Example:

Beagle e-mail virus gained additional functionality with each version.

# Beagle Evolution

Source: J. Gordon, infectionvectors.com

- More than 100 variants, not counting associated components.





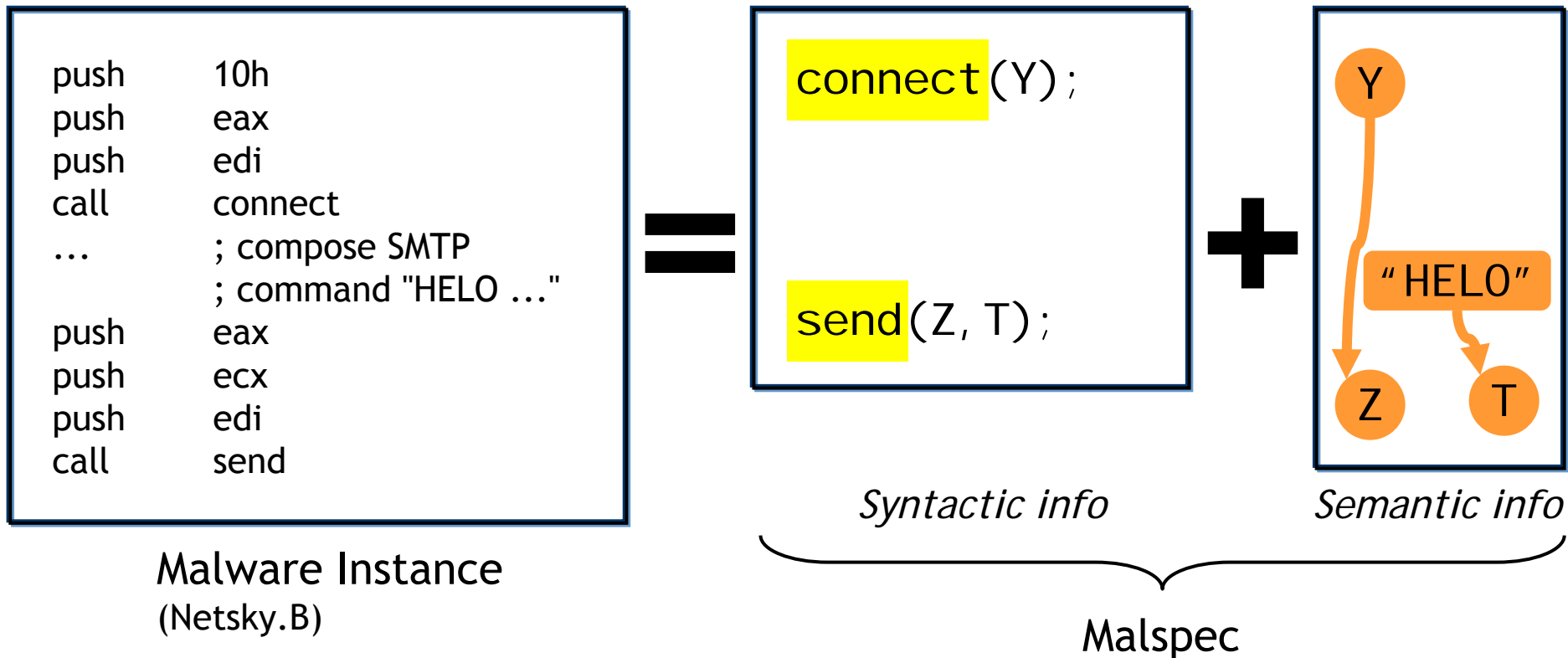
# Describing Malicious Behavior

[Christodorescu et al., Oakland 2005]

- Informal description:  
“Mass-mailing virus”
- A more precision description:  
“A program that:  
sends messages containing copies of  
itself,  
using the SMTP protocol,  
in a large number over a short period  
of time.”

# Malspec

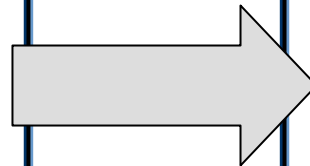
- A specification of behavior.



# Obfuscation Preserves Behavior

```
push    10h
push    eax
push    edi
call    connect
...     ; compose SMTP
        ; command "HELO ..."

push    eax
push    ecx
push    edi
call    send
```



```
push    10h
nop
push    eax
xor     eax, ebx
xor     eax, ebx
push    edi
call    connect
...     ; compose SMTP
        ; command "HELO ..."

push    eax
push   eax
pop     eax
push    ecx
push    edi
call    send
```

- Junk insertion + code reordering.

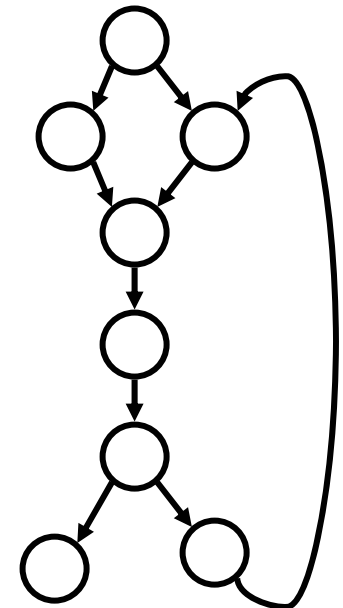
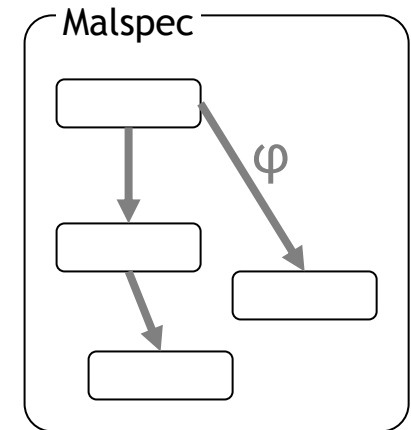
# Detection Using Malspecs

Static detection:

Given an executable binary,  
check whether it satisfies  
the malspec.

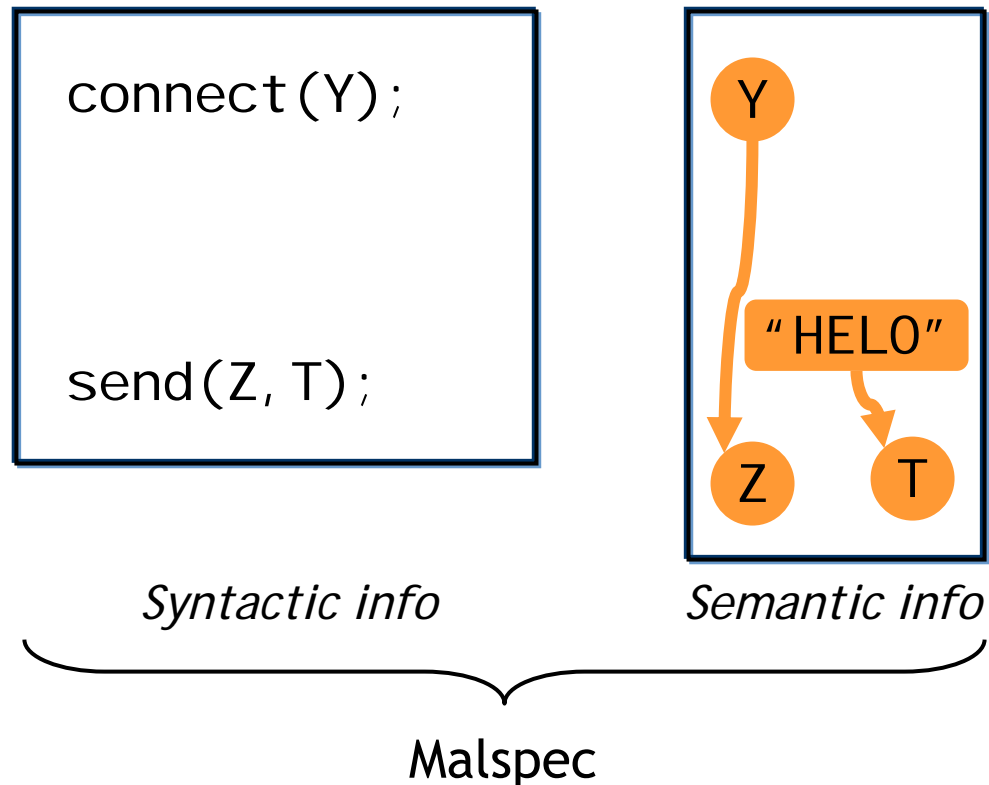
Just like model checking, but...

- Malicious code allows no assumptions to be made
- Real-time constraints



# A Behavior-Based Detector

- Match the syntactic constructs, then check the semantic information.



# Check the Semantic Info

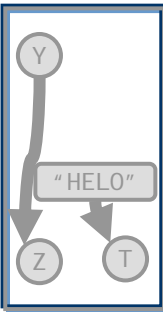
Program (Netsky.O):

```
push    10h
push    eax
push    [ebp+s]
call    connect
...
push    ebx
lea    eax, [ebp+s]
push    eax
call    send_email
```

*send\_email()*

```
... ; compose SMTP
      ; command "HELO ..."
lea    eax, [ebp+arg1]
push   eax
lea    eax, [ebp+buffer]
push   eax
call   SMTP_send_and_rcv
```

```
connect(Y);
send(Z, T);
```



Syntactic info

Semantic info

Malspec

*SMTP\_send\_and\_rcv()*

```
push    eax
push    [ebp+arg1]
mov     eax, [ebp+arg2]
push    [eax]
call    send
```

# Check with the Oracle

- Assume we have an oracle that can validate value predicates.

Does  
**eax before == ebx after**  
for the code sequence:

```
push eax  
call foo  
mov ebx, [ebp+4]
```

?



Yes.

# Check the Semantic Info

Program (Netsky.O):

```
push    10h
push    eax
push    [ebp+s]
A: call    connect
```

```
...
push    ebx
lea    eax, [ebp+s]
push    eax
call    send_email
```

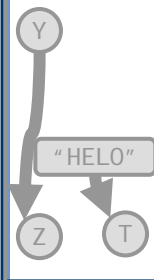
*send\_email()*

```
... ; compose SMTP
; command "HELO ..."
lea    eax, [ebp+arg1]
push    eax
lea    eax, [ebp+buffer]
push    eax
call    SMTP_send_and_rcv
```

connect(Y);

send(Z, T);

Syntactic info



Malspec

*SMTP\_send\_and\_rcv()*

```
push    eax
push    [ebp+arg1]
mov     eax, [ebp+arg2]
push    [eax]
B: call    send
```



# Query the Oracle

Program (Memory)

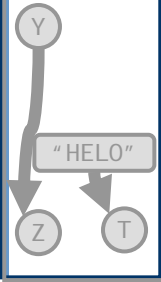
connect(x)

A:

```
push
push
push
call
...
push
lea
push
call
```

Does  
`memory[ebp@A+4] ==`  
`memory[ebp@B+4]` hold  
for the code sequence  
between A and B?

Yes.



Semantic info

pec

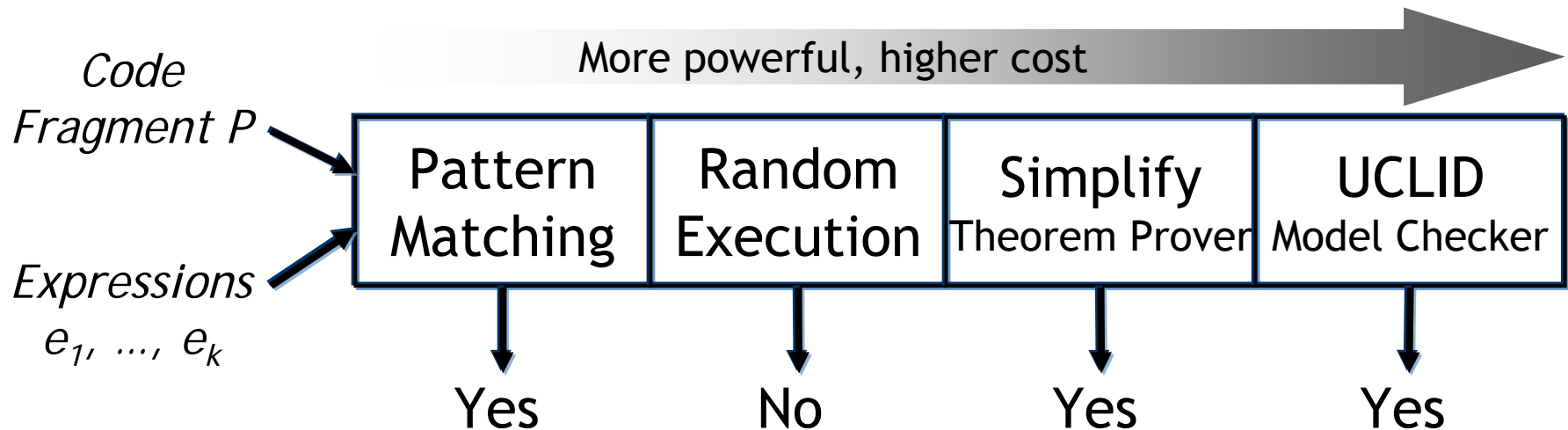
and\_rcv()

```
arg1]
ebp+arg2]
```

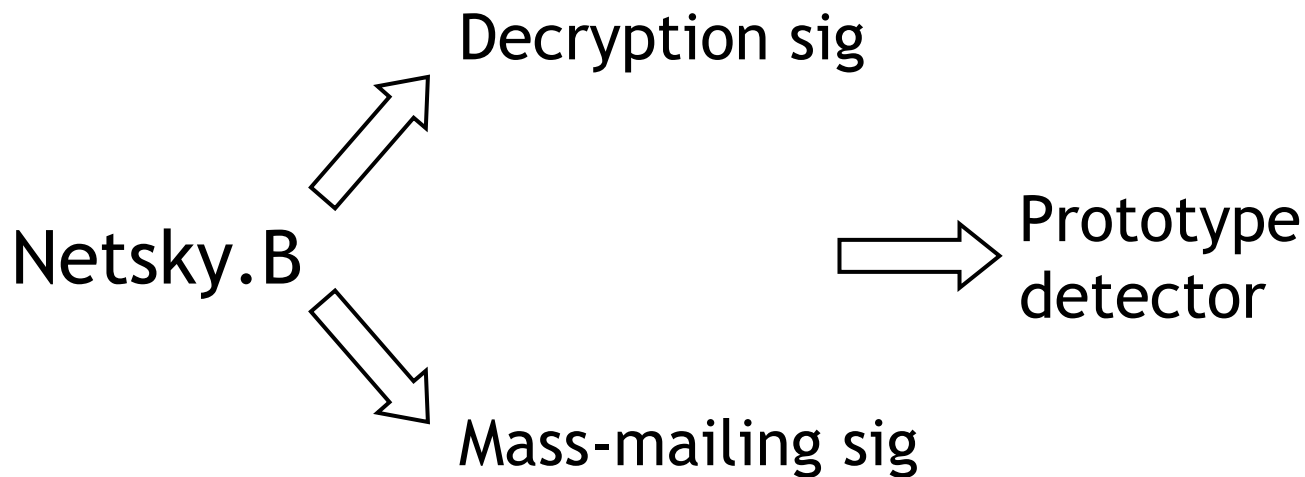
# A Recipe for an Oracle



- Instance of **program verification problem**:  
*Does program  $P$  respect property  $\phi$  ?*



# Evaluation of Malspecs



Netsky.C	✓
Netsky.D	✓
Netsky.O	✓
Netsky.P	✓
Netsky.T	✓
Netsky.W	✓

McAfee uses individual signatures for each worm.

**Malspecs provide forward detection.**

# Additional Information

- Papers

- M. Christodorescu and S. Jha, Testing Malware Detectors, *International Symposium on Testing and Analysis (ISSTA)*, 2004
- M. Christodorescu, S. Seshia, S. Jha, D. Song, and R. Bryant, *Semantics-Aware Malware Detection*, *IEEE Symposium on Security and Privacy (Oakland)*, 2005.

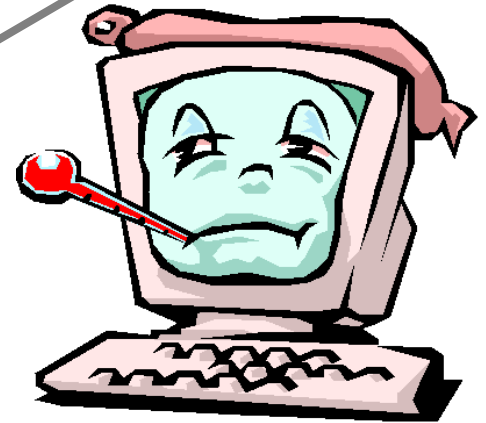
- Website

- <http://www.cs.wisc.edu/wisa/>

# Problem 2: Spec. Imprecision

Too general = false positives

→ Angry users



→ Infected machines

Too specific = false negatives

# Our Automatic Solution

**MINIMAL**: a technique for mining malicious specifications

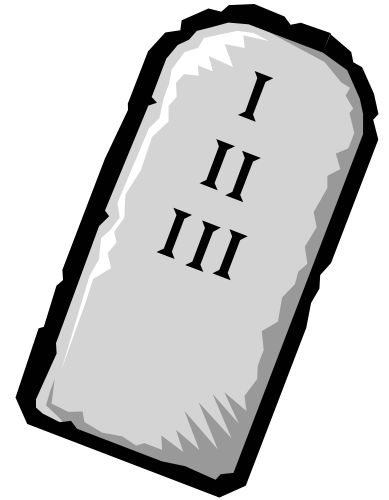
- (Mostly) automatic
- Flexible specification language
- Fast
- Performs well (compared to a human expert)

# Specification Language

# What's In a Specification?

Requirements for obfuscation resilience:

1. Describe **only relevant** operations
2. **Capture dependencies** where present
3. **Preserve independence** of operations





# Specifying Malicious Operations

- We chose **system calls**
  - Compatible with specifications for behavior-based detectors
  - Define interface between trusted OS and untrusted programs
- Mining algorithm is not restricted to the system-call interface.

# Specifying Malicious Constraints

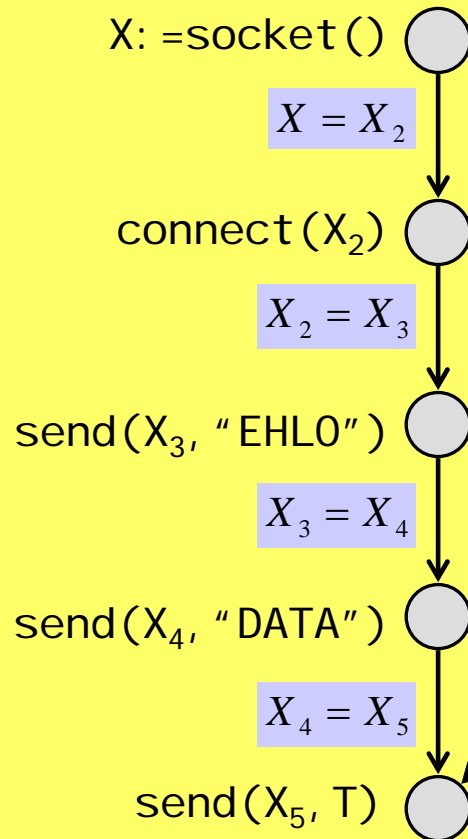
- Program operations are insufficient to distinguish malicious from benign.
- We need to capture relations between operations:

`F=open("file"); read(F, buf); send(S, buf)`

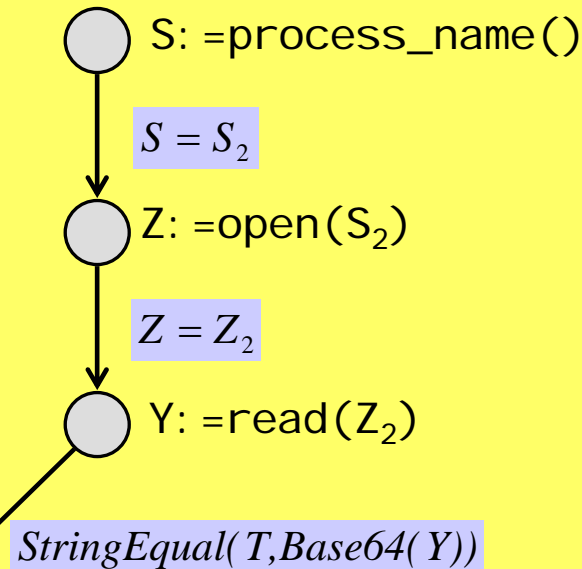
**Constraints** = logical formulas over system-call arguments

# A Sample Specification (Malspec)

## Send Email

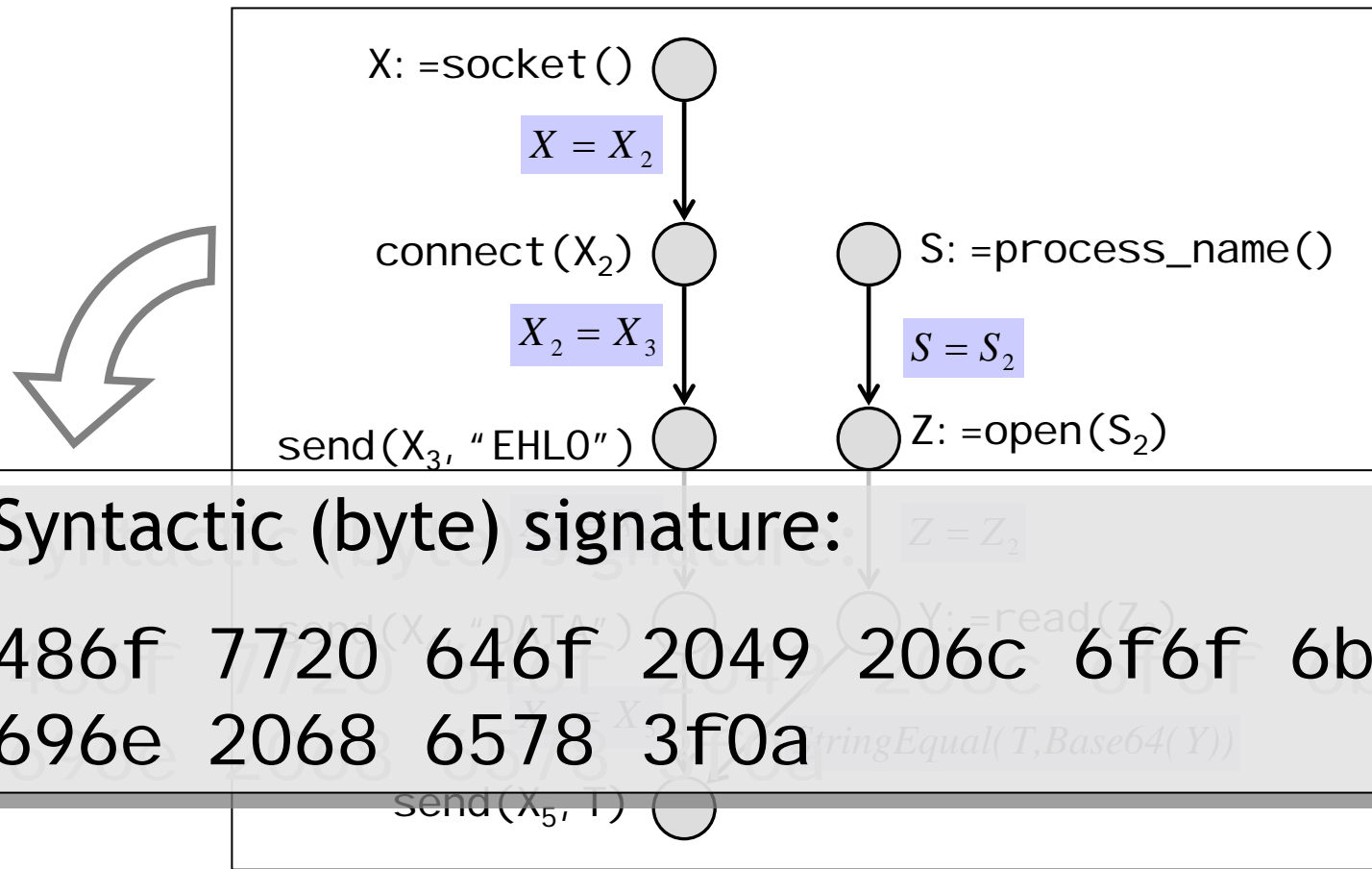


## Read/copy self



# A Sample Specification (Malspec)

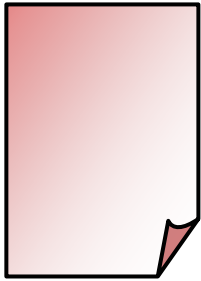
- Rich specification can be “dumbed down”



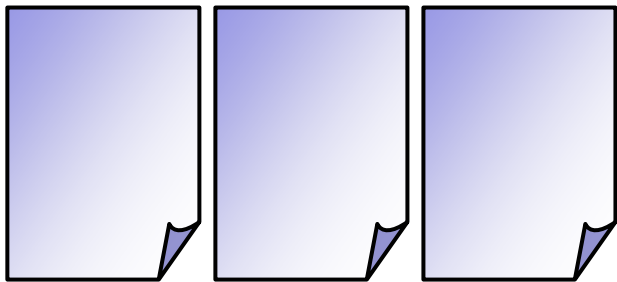
# Mining Algorithm

# The Specification Mining Problem

Known malware



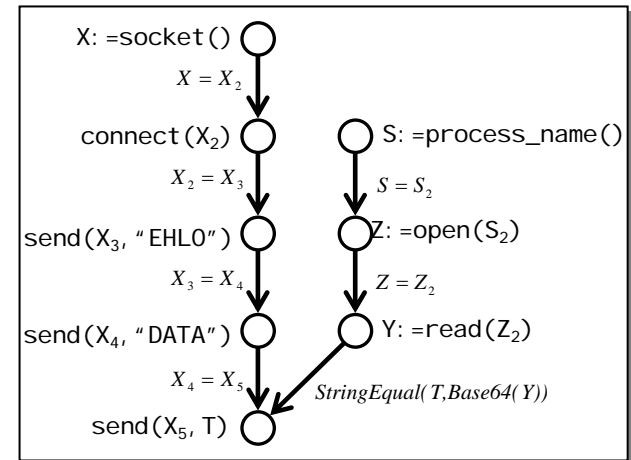
Known benign programs



MINIMAL

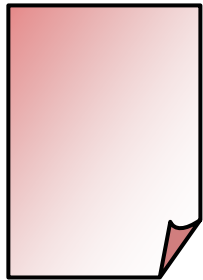


Specification of malicious behavior

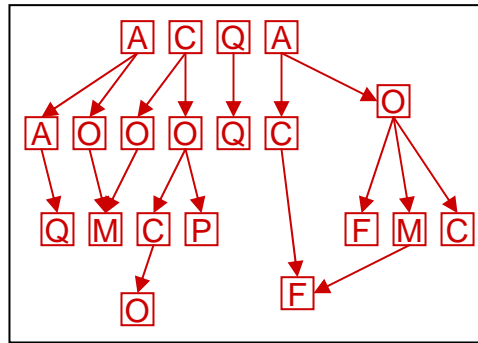


# The Basic Mining Operation

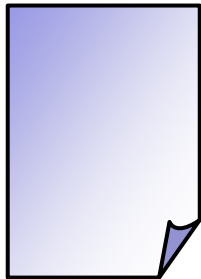
Known malware



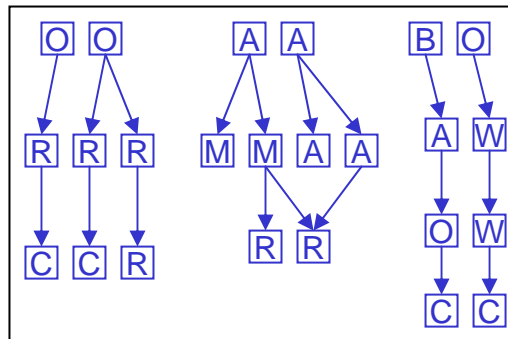
Malware dependence graph



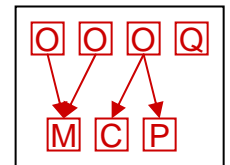
Known benign program



Benign dependence graph



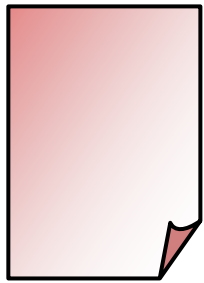
Minimal malspec



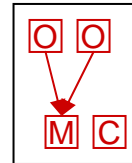
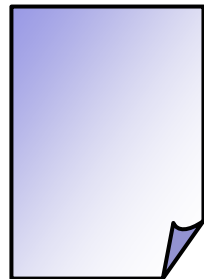
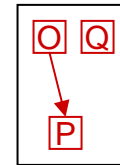
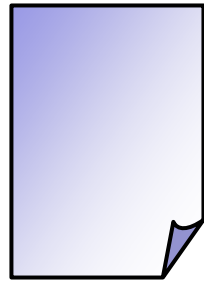
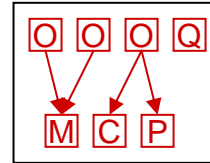
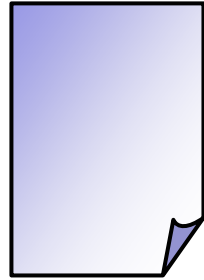
**Step 1** Compute dependence graphs

**Step 2** Compute graph difference

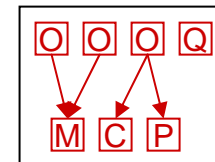
# Multi-Program Mining



vs.



Maximal union of malspecs:

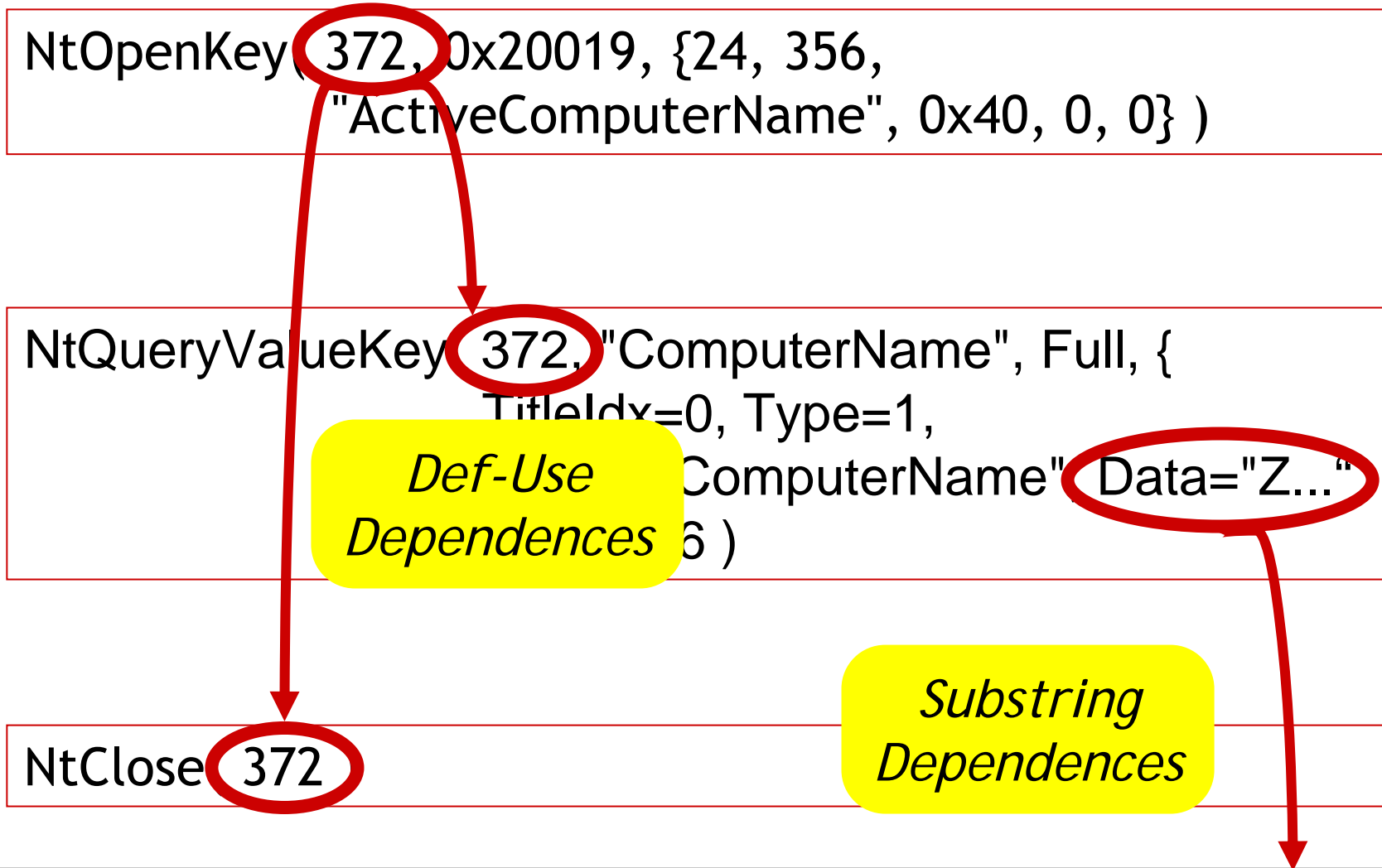




# System-Call Dependence Graph

- We use a **dynamic analysis** to construct the dependence graph
  - Static analysis too imprecise on binary code
- Steps:
  1. Collect system-call trace
  2. **Infer dependencies** between system calls
  3. Construct (an underapproximation of the) dependence graph

# Discovering Dependences



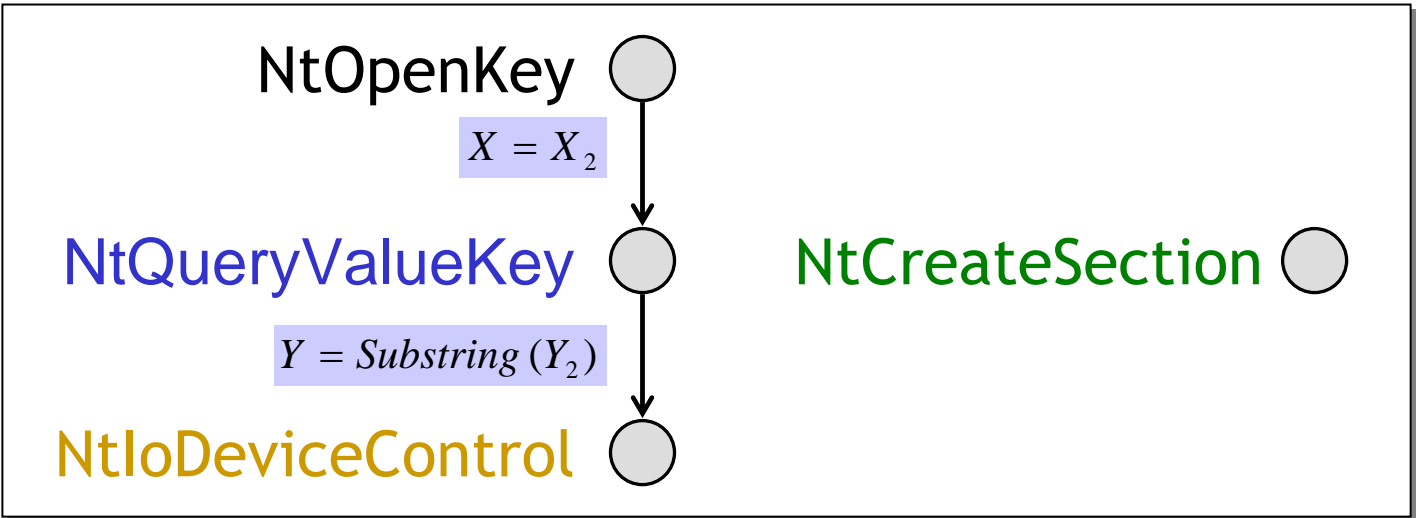
# Discovering Local Constraints

- Access to well-defined resources:
  - Windows registry
  - Access to self
  - System files/directories

```
NtCreateFile ( ..., { ..., "I-Worm.Mydoom.l.exe" ... }, ... )
```

# Dependence Graph Example

```
NtOpenKey( 372, ... )  
NtQueryValueKey( 372, ..., { ..., Data="Z...", ... } )  
NtCreateSection( ... )  
NtIoDeviceControl( ..., OutBuffer="..Z.." ... )
```



# Graph Differencing

Problem:

Find the smallest subgraph of malicious operations that does not appear in any benign graph.

Solution:

*Minimal Contrast Subgraph*

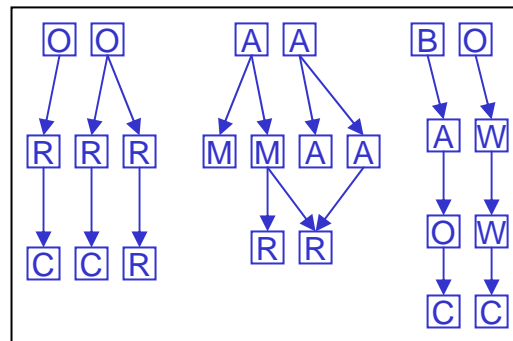
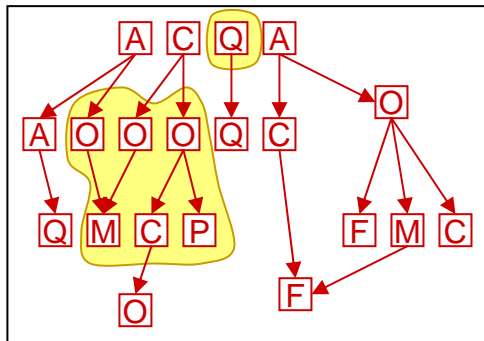
[Ting, Bailey “Mining Minimal Contrast Subgraph Patterns”, SDM 2006]

# Minimal Contrast Subgraphs

- Idea:
  - Minimal contrast subgraphs and maximal common edge sets are duals.
- Finding maximal common edge set:
  - Consider all edge sets (order by size)
  - Eliminate edge-set candidates as early as possible

# Mining Contrast Subgraphs

Malware  
dependence  
graph



Benign  
dependence  
graph

- Size of graphs:  $N = 100K-1.5M$  nodes
- Worst-case complexity:  $O(N!)$

# Heuristics Reduce Problem Size

- Normalize dependence graph
  - Replace system-call sequences with shorter equivalents
- Eliminate disconnected subgraphs
- Eliminate trivial subgraphs

[see paper for details]



# Evaluation

# Evaluating MINIMAL

- Goals:
  - Compare MINIMAL malspecs with those from human expert
  - Use mined malspecs with behavior-based detector

# Experimental Setup

- Trace collection in Windows 2000:
  - Malware samples run with no user input (cf. expected execution model)
  - Benign samples run with normal user input
  - Execution for 1 or 2 minutes
- 16 malware samples:
  - Netsky, MyDoom, Beagle
- 6 benign programs:
  - Firefox, Thunderbird, installers

# MINIMAL vs. Human Expert

MINIMAL  
malspecs



Behavioral features as given by  
Symantec:

A screenshot of a Symantec security response page for the malware W32.Netsky@mm. The page is titled "W32.Netsky@mm" and shows a "Risk Level 2: Low". It includes a "Printer Friendly Page" link and three tabs: "SUMMARY", "TECHNICAL DETAILS", and "REMOVAL". The "SUMMARY" tab is active, displaying the following information:

**Discovered:** February 16, 2004  
**Updated:** February 13, 2007 12:17:32 PM  
**Also Known As:** WORM\_NETSKY.A [Trend]  
**Type:** Worm  
**Systems Affected:** Windows 2000, Windows 95, Windows 98, Windows Me, Windows NT, Windows Server 2003, Windows XP

When W32.Netsky@mm runs, it does the following:

1. Creates a mutex named "AdmMoodownJKIS003." This mutex allows only one instance of the worm to execute in memory.

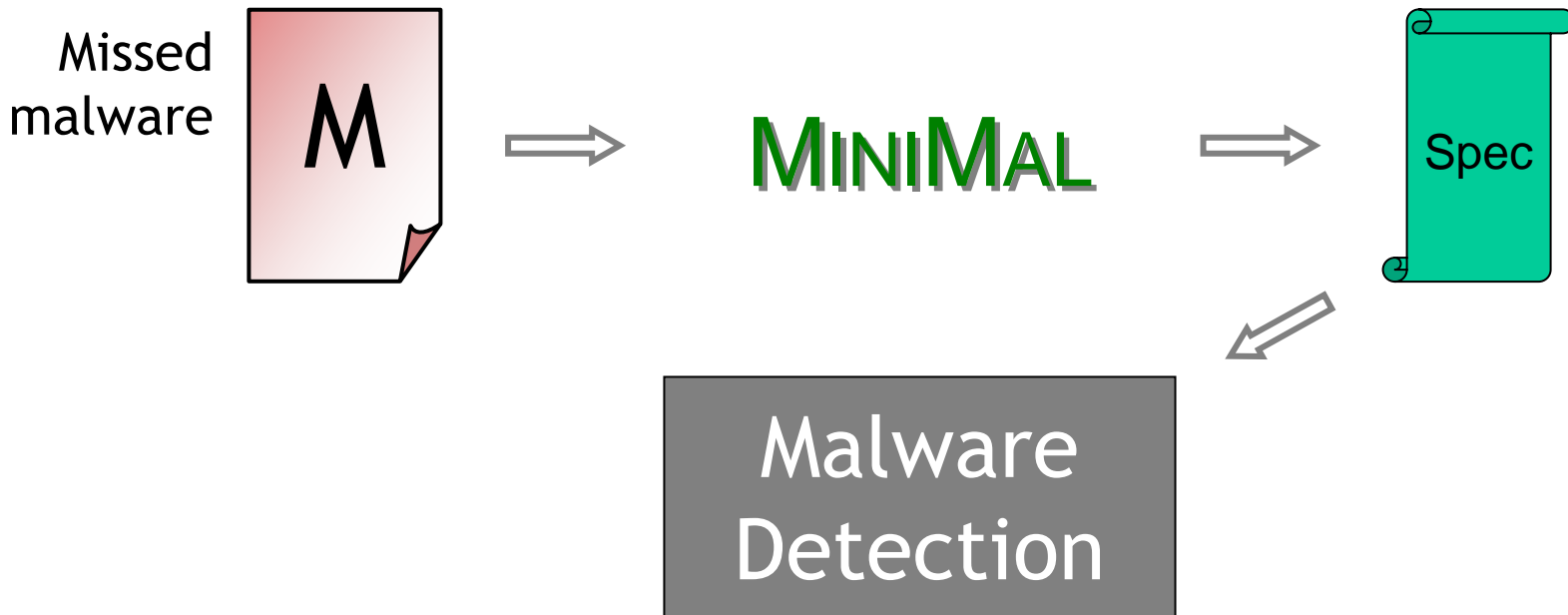
On the right side of the page, there are links for "PRINT THIS PAGE" and "RATE THIS PAGE", a "TOP THINGS TO" section with a list of links, and a "Search This" section with a search box and the text "Search by name Example: W32.B".

# Mined Malspecs for Netsky.A



	MINIMAL malspecs
Create mutex	✓
Self-installation	✓
Modify boot sequence	✓
Terminate antivirus	✓
Email self as ZIP file	
Copy self to network drive	

# MINIMAL Specs in Detection



- Using mined malspecs in semantics-aware malware detection:

Netsky.A malspec → Netsky.D, E, F, ...

# Future Work

## ~~Limitations of~~ MINIMAL

- Sensitive to **test environment**
  - Malicious behavior might not be observed during tracing.
- **Underapproximation** of dependence graph
  - Complex constraints are not discovered.
- Sensitive to **test-set selection**
  - Not all differences are malicious behaviors.

*Questions?*

# Mining Specifications of Malicious Behavior

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