

COL866: Quantum Computation and Information

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Quantum Computation: Phase estimation

Quantum Computation

Phase estimation

Phase estimation

Suppose a unitary operator U has an eigenvector $|u\rangle$ with eigenvalue $e^{2\pi i\varphi}$. The goal is to estimate φ .

- We will use the assumption that there are black-boxes that:
 - prepare the state $|u\rangle$, and
 - perform the controlled- U^{2^j} operation.
- We will describe a phase estimation procedure that uses two registers:
 - A t -qubit register initially in state $|0\dots 0\rangle$ (the value of t to be decided later), and
 - a register that begins in the state $|u\rangle$.

Quantum Computation

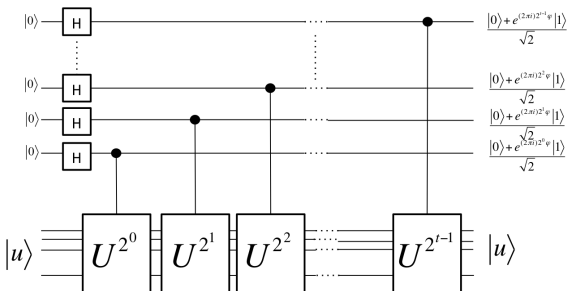
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- Claim 1: The final state of the first register in the circuit below is given by:

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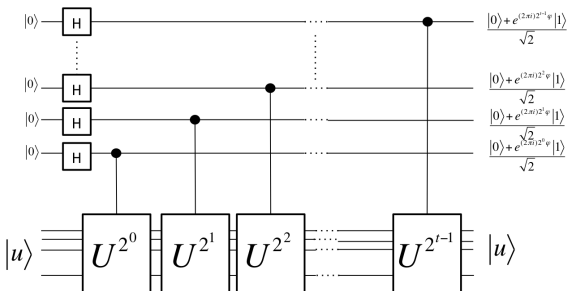
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Claim 2

It is sufficient to run the phase estimation technique with $t = n + \log\left(2 + \frac{1}{2\varepsilon}\right)$ in order to obtain φ accurate to n bits with probability at least $(1 - \varepsilon)$.

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It is sufficient to run the phase estimation technique with $t = n + \lceil \log(2 + \frac{1}{2\varepsilon}) \rceil$ in order to obtain φ accurate to n bits with probability at least $(1 - \varepsilon)$.

Proof sketch

- Let $0 \leq b \leq 2^t - 1$ be an integer such that $\frac{b}{2^t} = [0 \cdot b_1 \dots b_t]$ is the best t bit approximation to φ that is less than φ . Let $\delta = \varphi - \frac{b}{2^t}$ (which implies $0 \leq \delta \leq 2^{-t}$).
- Claim 2.1: Applying the inverse Fourier transform on the first register in state $\frac{1}{2^{t/2}} \sum_{k=0}^{2^t-1} e^{(2\pi i)\varphi k} |k\rangle$ ends in the following state:

$$\frac{1}{2^t} \sum_{k,l=0}^{2^t-1} e^{-\frac{(2\pi i)kl}{2^t}} e^{(2\pi i)\varphi k} |l\rangle.$$

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- Claim 2.2: Let α_l be the amplitude of $|(b+l) \bmod 2^t\rangle$. Then
$$\alpha_l = \frac{1}{2^t} \left(\frac{1 - e^{(2\pi i)(2^t\varphi - (b+l))}}{1 - e^{(2\pi i)(\varphi - (b+l)/2^t)}} \right) = \frac{1}{2^t} \left(\frac{1 - e^{(2\pi i)(2^t\delta - l)}}{1 - e^{(2\pi i)(\delta - l/2^t)}} \right).$$

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$$\Pr[|m - b| > e] \leq \frac{1}{2(e-1)}.$$

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- The claim follows by setting $t = n + p$ and $\varepsilon = \frac{1}{2(2^p - 1)}$. □

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- The phase estimation protocol works when the second register is set to the eigenstate $|u\rangle$. In general, this may not be feasible.
- Observation: Any general state $|\psi\rangle$ may be written in terms of the eigenstates of U as $\sum_u c_u |u\rangle$.
- Exercise: The phase estimation procedure takes state $(|0\rangle)(\sum_u c_u |u\rangle)$ to $\sum_u c_u |\tilde{\varphi}_u\rangle |u\rangle$. If $t = n + \lceil \log(2 + \frac{1}{2\varepsilon}) \rceil$, then the probability of measuring φ_u accurate to n bits at the end of the phase estimation procedure is at least $|c_u|^2(1 - \varepsilon)$.

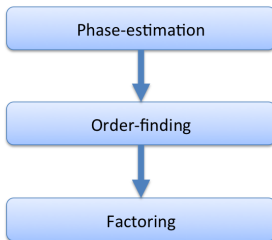
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- Phase estimation enables us to design quantum algorithms for the **order-finding** and **factoring** problems.



End