# COL702: Advanced Data Structures and Algorithms 

## (Semester-I-2022-23)

## Grading

| Grading component | $\#$ | \% |
| :--- | :--- | :--- |
| Homework | 5 | $20 \%$ (best 4 out of 5) |
| Quiz | 5 | $20 \%$ (best 4 out of 5) |
| Minor | 1 | $20 \%$ |
| Final | 1 | $35 \%$ |
| Comprehension quiz | See below | $5 \%$ |

- Comprehension quizzes are online quizzes conducted through Moodle/Gradescope that students are supposed to take for a better understanding of the material.


## Logistics

| Grading <br> component | Comments |
| :--- | :--- |
| Homework | •Can be done in groups of size at most 2. <br> Submissions to be made over Gradescope. Please avoid missing <br> deadlines. <br> We will prefer if you submit typed solutions. Latex source files will be <br> provided. <br> May include a programming component. <br> Quiz <br> Comprehension <br> quiz <br> -Conducted using online Gradescope feature or in class. <br> - We may create multiple versions of the quiz.We quiz per lecture component. There may be around 20 such quiz. <br> been completed. However, we will have a flexible deadline. |

## Important points

- Please register on Piazza (no code required).
- You will be registered on Gradescope by syncing with Moodle.
- Course webpage:
- https://www.cse.iitd.ac.in/~rjaiswal/Teaching/2022/COL702/
- Textbook
- -Algorithms by Dasgupta, Papadimitriou, and Vazirani.
- Algorithm Design by Jon Kleinberg and Eva Tardos.
- Audit:
- You will need to score at least $40 \%$ for Audit-pass. You should plan to take all the grading components even if you audit.


## Analyzing algorithms

A royal mathematical challenge (1202):

Suppose that rabbits take exactly one month to become fertile, after which they produce one child per month, forever. Starting with one rabbit, how many are there after $n$ months?


Leonardo da Pisa, aka Fibonacci

## The proliferation of rabbits

Rabbits take one month to become fertile, after which they produce one child per month, forever.

|  | Fertile | Not fertile |
| :--- | :---: | :---: |
| Initially |  |  |
| One month |  |  |
| Two months |  |  |
| Three months |  |  |
| Four months |  |  |
| Five months |  |  |

## The Fibonacci sequence

$$
F_{1}=1, \quad F_{2}=1, \quad F_{n}=F_{n-1}+F_{n-2}
$$

These numbers grow very fast: $F_{30}>10^{6}$ !
In fact, $F_{n} \approx 2^{0.694 n} \approx 1.6^{n}$, exponential growth.

## The Fibonacci sequence

$$
F_{1}=1, F_{2}=1, F_{n}=F_{n-1}+F_{n-2}
$$

Can you see why $F_{n}<2^{n}$ ?

## Computing Fibonacci numbers

```
function Fib1(n)
if n = 1 return 1
if n = 2 return 1
return Fib1(n-1) + Fib1(n-2)
```

A recursive algorithm

Two questions we always ask about algorithms:
Does it work correctly?
How long does it take?

## Running time analysis

```
function Fib1(n)
if n = 1 return 1
if n = 2 return 1
return Fib1(n-1) + Fib1(n-2)
```

Exponential time... how bad is this?
Eg. Computing $F_{200}$ needs about $2^{140}$ operations. How long does this take on a fast computer?

## IBM Summit



Can perform up to 200 quadrillion ( $=200 \times 10^{15}$ ) operations per second.

## Is exponential time all that bad?

The Summit needs $2^{82}$ seconds for $F_{200}$.

Time in seconds
$2^{10}$
$2^{20}$
$2^{30}$
$2^{40}$
$2^{45}$
$2^{51}$
$2^{57}$
$2^{60}$

Interpretation
17 minutes
12 days
32 years
cave paintings
homo erectus discovers fire extinction of dinosaurs creation of Earth origin of universe

## Post mortem

What takes so long?
Let's unravel the recursion...

```
function Fib1(n)
if n = 1 return 1
if n = 2 return 1
return Fib1(n-1) + Fib1(n-2)
```



The same subproblems get solved over and over again!

## A better algorithm

Subproblems: $F_{1}, F_{2}, \ldots, F_{n}$. Solve them in order and save their values!

```
function Fib2(n)
Create an array fib[1..n]
fib[1] = 1
fib[2] = 1
for i = 3 to n:
    fib[i] = fib[i-1] + fib[i-2]
return fib[n]
```

[1] Does it return the correct answer?
[2] How fast is it?

## Polynomial vs. exponential

Polynomial running times:

Exponential running times:

To an excellent first approximation:
polynomial is reasonable exponential is not reasonable

This is one of the most fundamental dichotomies in the analysis of algorithms.

