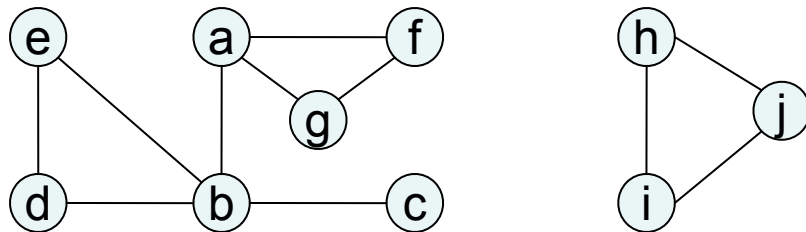


Lecture 5: Undirected connectivity

An undirected graph is *connected* if there is a path between any pair of nodes.



This graph has 2 *connected components*.

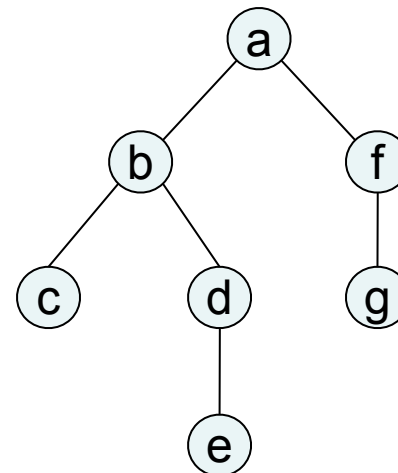
`explore(G,v)` returns the connected component containing `v`.

To explore the rest of the graph, restart `explore()` elsewhere.

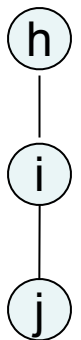
DFS decomposes an undirected graph into its connected components!

```
procedure dfs(G)
for all v in V:
    visited[v] = false
for all v in V:
    if not visited[v]:
        explore(G,v)
```

`explore(G,a)`



`explore(G,h)`



Running time analysis

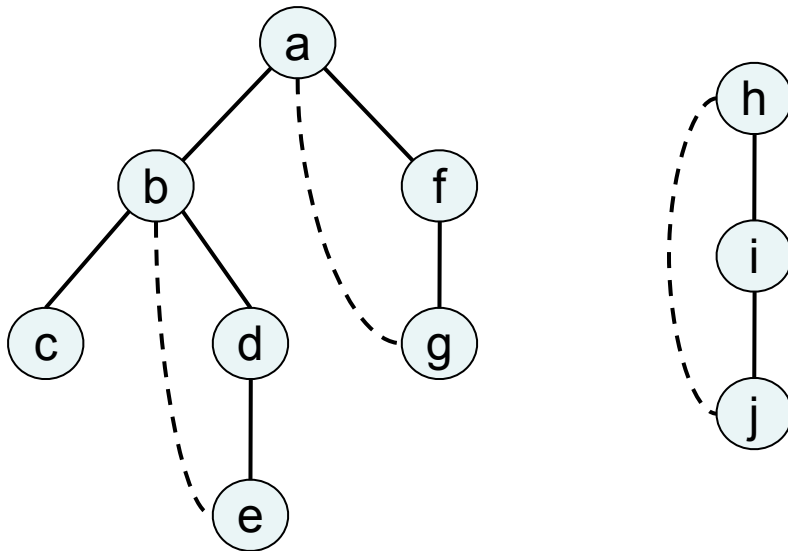
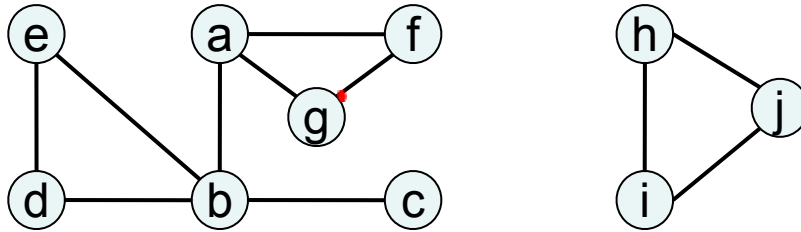
```
procedure explore(G,v)
visited[v] = true
for each edge (v,u) in E:
    if not visited[u]:
        explore(G,u)
```

```
procedure dfs(G)
for all v in V:
    visited[v] = false
for all v in V:
    if not visited[v]:
        explore(G,v)
```

How long does $\text{dfs}(G)$ take?

$\text{explore}(G,v)$ is called exactly once for each node v .

DFS search forest



Terminology:

DFS search forest consisting of two *DFS search trees*

- *tree edge*: traversed by DFS
- - - *back edge*: not traversed (led to a node already visited)