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**COL351: Analysis and Design of Algorithms****Instructor:** Ragesh Jaiswal

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1. Is the factor-2 approximation for the greedy algorithm for the Minimum Makespan problem shown in class tight?
2. Recall the Minimum Makespan problem discussed in class. We saw a greedy approximation algorithm and showed an approximation guarantee of 2. Consider the following slightly modified algorithm:

**GreedySortMakespan**

- Consider jobs in decreasing order of duration
- While all jobs are not assigned
  - Assign the next job (as per the order defined) to a machine with least load

Let  $OPT$  denote the value of an optimal solution and  $G$  be the value of the above greedy solution. Then we will argue that  $G \leq (4/3) \cdot OPT$ . WLOG, let us assume that  $d(1) \geq d(2) \geq \dots \geq d(n)$ .

Let us call an problem instance  $(d(1), \dots, d(n))$  *nice* if as per the above greedy algorithm, the job with the maximum finishing time is  $n$ . We will first argue that it is sufficient to analyze the approximation guarantee for nice instances.

Claim 1.1: If our greedy algorithm gives a factor  $c$  approximation for nice instances, then it gives factor  $c$  approximation for all instances.

We can now focus on only nice instances. Again, for any nice instance, we use  $G$  to denote the value of the greedy solution and  $OPT$  to denote the value of an optimal solution. Let us break the analysis into two cases: (i)  $d(n) \leq OPT/3$ , and (ii)  $d(n) > OPT/3$ . For case (i), we can make the following claim:

Claim 1.2: If  $d(n) \leq OPT/3$ , then  $G \leq (4/3) \cdot OPT$ .

Now, consider the case when  $d(n) > OPT/3$ . We first show that  $n \leq 2m$ .

Claim 1.3: If  $d(n) > OPT/3$ , then  $n \leq 2m$ .

Once, we realize the above the next thing we show is that in this case, the greedy algorithm gives an optimal solution.

Claim 1.4 If  $d(n) > OPT/3$ , then **GreedySortMakespan** returns an optimal solution.

Combining all the above claims, we get that the above greedy algorithm gives an approximation guarantee of  $(4/3)$ .

3. Given a sorted array  $A$  containing  $n$  distinct integers, design an algorithm to determine if there is an index  $i$  such that  $A[i] = i$ .