

Lecture 5: September 24

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5.1 Threshold for connectedness in Random Graphs

Theorem 5.1 Let $\alpha = \alpha(n)$ be a function with $\alpha(n) \rightarrow \infty$ as $n \rightarrow \infty$. Then $S(n) := \frac{\ln(n) - \alpha(n)}{n}$ is a lower threshold function and $t(n) := \frac{\ln(n) + \alpha(n)}{n}$ is an upper threshold function for connectivity of $G(n, p)$.

Proof:

1. Lower Threshold Function: If $p \leq S$ then by Theorem 4.3 $G(n, p)$ has isolated vertices approximately almost surely. Therefore, $G(n, p)$ graph are not connected.
2. Upper Threshold Function: We assume $\alpha(n) \leq \ln(n)$ for all n . For each $k \leq n$, let X_k be the random variable which counts the number of connected components of size k .

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Let $Y = \sum_{k \leq n} X_k$ and let $X := \sum_{k \leq \lfloor \frac{n}{2} \rfloor} X_k$. Note here Y counts all the connected components in $G(n, p)$ whereas

X counts all the connected components of size at most $\lfloor \frac{n}{2} \rfloor$.

We thus have $Y \geq 1$ implies $X \geq 1$ because there is some smaller connected component inside the bigger component.

Therefore (Using Markov inequality), $P(Y \geq 1) \leq P(X \geq 1) \leq \frac{E(X)}{1} = E(X)$

Claim 5.2 $E(X) \rightarrow 0$ as $n \rightarrow \infty$

Proof: Consider $S \subseteq V$,

$P(S \text{ forms a maximal connected component in } G(n, p)) \leq P(\text{no vertex in } S \text{ is connected to } S^c)$

$P(\text{no vertex in } S \text{ is connected to } S^c) = (1 - p)^{|S|(n-|S|)} \leq e^{-p|S|(n-|S|)} \leq e^{-t|S|(n-|S|)}$

Note the last inequality holds if $p \geq t$ (and we can this Equation 1).

Now we are going to analyze the expectation of each $X_k \forall k \leq \lfloor \frac{n}{2} \rfloor$

$$E(X_k) \leq e^{-tk(n-k)} \binom{n}{k}$$

Now we need to give a good estimation of the Binomial Coefficient. Using Stirling's formula :

$$E(X_k) \leq \left(\frac{ne}{k}\right)^k e^{-tk(n-k)}$$

Next inserting the formula for t we obtain :

$$E(X_k) \leq e^{-\alpha(n)} \left(\frac{e^{1-(1-\frac{1}{n})\alpha(n)} + \frac{k}{n}(\ln(n) + \alpha(n))}{k} \right)^k = e^{-\alpha(n)} (B_k)^k$$

We can use the above Equation 2 and now turn our attention to B_k .

Case 1: When $k \leq \lfloor n^{\frac{3}{4}} \rfloor$ then :

$$\frac{k}{n}(\ln(n) + \alpha(n)) \leq \frac{1}{n^{\frac{1}{4}}}(\ln(n) + \alpha(n)) \leq \frac{2 \ln(n)}{n^{\frac{1}{4}}} \rightarrow 0$$

The last approximation can be realized using the Taylor Series expansion. Henceforth, the numerator in (2) is at most a constant say $c > 0$. If $c \leq k$, then with $\theta = \frac{c}{c+1}$ we have $B_k \leq \frac{c}{k} \leq \theta < 1$ for any $k \geq c+1$.

So, $\sum_{k=c+1}^{\lfloor n^{\frac{3}{4}} \rfloor} (B_k)^k \leq \sum_{k=c+1}^{\lfloor n^{\frac{3}{4}} \rfloor} \theta^k \leq c_1$ for some constant $c_1 > 0$

Case 2: When $\lfloor n^{\frac{3}{4}} \rfloor \leq k \leq \lfloor \frac{n}{2} \rfloor$ then:

$$B_k \leq e^{1-\frac{1}{4} \ln(n) - (\frac{1}{2} - \frac{1}{k})\alpha(n)}$$

Since $k \geq 2$, $\frac{1}{2} - \frac{1}{k} \geq 0$ so the exponent tends to infinity. So $B_k \leq \theta$ for some $\theta < 1$

Therefore, $\sum_{k=\lfloor n^{\frac{3}{4}} \rfloor}^{\frac{n}{2}} (B_k)^k < c_2$ for some constant $c_2 > 0$.

Thus $E(X) = \sum_{k=1}^{\frac{n}{2}} E(X_k) \leq e^{-\alpha(n)}(c_1 + c_2 + c_3) \rightarrow 0$ as $n \rightarrow \infty$.

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