

Lecture 14: September 16

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This lecture's notes illustrate some uses of various L^AT_EX macros. Take a look at this and imitate.

14.1 Some theorems and stuff

Proposition 14.1 *Let $\alpha(n)$ be a function such that $\alpha(n) \rightarrow \infty$ as $n \rightarrow \infty$, then for a random graph $G(n, p)$ it holds*

$$pN - \alpha(n)n \leq e(G) \leq pN + \alpha(n)n$$

almost asymptotically surely (a.a.s).

Proof: Let

$$\begin{aligned} X &= \sum_{e \in \binom{V}{2}} X_e \\ X &= Np \\ \text{Var}(X) &= Npq \end{aligned}$$

where v is the number of vertices, and p and q holds their standard meanings ($q = 1 - p$)
Applying Chebyshev's inequality,

$$\begin{aligned} P(|X - Np| \geq \alpha(n)n) &\leq \frac{\text{Var}(X)}{(\alpha(n))^2 n^2} \\ P(|X - Np| \geq \alpha(n)n) &\leq \frac{Npq}{(\alpha(n))^2 n^2} \\ P(|X - Np| \geq \alpha(n)n) &\leq \frac{pq}{(\alpha(n))^2} \end{aligned}$$

where pq is a constant. Thus, $\frac{pq}{\alpha(n)^2} \rightarrow 0$ as $n \rightarrow \infty$. ■

14.2 Thresholds for connectivity

Definition 14.2 *Let Q be a graph property and let $r, s, t : N \rightarrow \mathbb{R}$ be functions. We say that*

1. t is a lower threshold function for the graph property Q in the random graph model $G(n,p)$ if for any $p \leq t \forall n$ $G(n,p)$ doesn't have property Q a.a.s
2. s is an upper threshold function for the graph property Q in the random graph model $G(n,p)$ if for any $p \geq s \forall n$ $G(n,p)$ doesn't have property Q a.a.s

Note that s and t may coincide. Without loss of generality we can assume that $t \leq s$. Our aim is to bridge the gap between t and s to get better results.

We need a technical result which handles the situation of non-independent random variables. Let X_1, \dots, X_k be random variables. Define

$$\Delta_1 = \sum_{\substack{i,j \in \{1, \dots, k\} \\ i \neq j}} E(X_i, X_j)$$

We say $X_i \sim X_j$, $i \neq j$ if X_i and X_j are not independent.

$$\Delta_2 = \sum_{\substack{i,j \in \{1, \dots, k\} \\ i \sim j}} E(X_i, X_j)$$

Theorem 14.3 Let X_1, \dots, X_k be random variables (not necessarily independent) and let $X = \sum_{i=1}^k X_i$. We have

1.
$$\text{Var}(X) = \sum_i \text{Var}(X_i) + \sum_{\substack{i,j \in \{1, \dots, k\} \\ i \neq j}} \text{Cov}(X_i, X_j)$$

2. If X_i 's are Bernoulli random variables then

- (a)
$$\text{Var}(X) = E(X) - E(X)^2 + \Delta_1$$

- (b)
$$\text{Var}(X) \leq E(X) + \Delta_2$$

- (c)
$$P(X = 0) \leq \frac{1}{E(X)} - 1 + \frac{\Delta_1}{(E(X))^2}$$

- (d)
$$P(X = 0) \leq \frac{1}{E(X)} + \frac{\Delta_2}{(E(X))^2}$$

Note: (a) and (b) can be concluded using straight-forward calculations. (c) and (d) can be derived from (a) and (b) respectively, using Chebyshev's inequality.

Theorem 14.4 Let

$$s(n) = \frac{\ln(n) - \alpha(n)}{n}$$

$$t(n) = \frac{\ln(n) + \alpha(n)}{n}$$

where $\alpha(n) \rightarrow \infty$ as $n \rightarrow \infty$ and $\alpha(n) \leq \ln(n)$.

s is a lower threshold function and t is an upper threshold function for the property that $G(n,p)$ has no isolated vertices.

Proof:

1. Upper threshold

Consider $p \geq t$

We proved that the probability that $G(n,p)$ has isolated vertices is at most $n(1-p)^{(n-1)}$. This approaches zero as $n \rightarrow \infty$ when p is constant. But here p is a function of n . Now,

$$\begin{aligned} n(1-p)^{(n-1)} &\leq n(1-t)^{(n-1)} \leq \frac{n}{1-t}(1-t)^n \leq \frac{n}{1-t}e^{(-t)} \\ n(1-p)^{(n-1)} &\leq \frac{n}{1-t}e^{(-\ln(n)-\alpha(n))} \leq \frac{1}{1-t}e^{-\alpha(n)} \end{aligned}$$

As $n \rightarrow \infty$, $t \rightarrow 0$.

$$\frac{1}{1-t}e^{-\alpha(n)} \rightarrow 0$$

Hence proved.

2. Lower threshold

Let $p \leq s$. Let X be the random variable which counts the number of isolated vertices. We need to prove that

$$P(X = 0) \leq \frac{1}{E(X)} - 1 + \frac{\Delta_1}{(E(X))^2}$$

and further prove that R.H.S $\rightarrow 0$ as $n \rightarrow \infty$.

$$E(X) = n(1-p)^{(n-1)} \geq n(1-s)^{(n-1)} = \frac{n}{1-s}(1-s)^n \geq \frac{n}{1-s}e^{-n(s+s^2)}$$

(Using $1-x \geq e^{(-x-x^2)}$ for $x \in [0, 1/3]$)

$$E(X) = \frac{1}{1-s}e^{\alpha(n)+\beta(n)}, \text{ where } \beta(n) = \frac{\ln^2(n)-2\ln(n)+\alpha(n)}{n}$$

Now, $\beta n \rightarrow 0$ as $n \rightarrow \infty$

$\implies E(X) \rightarrow \infty$ as $n \rightarrow \infty$

$\implies \frac{1}{E(X)} \rightarrow 0$ as $n \rightarrow \infty$

Now consider, $\frac{\Delta_1}{E(X)^2}$

$$\begin{aligned} \Delta_1 &= \sum_{v,w \in V, v \neq w} P(X_v = 1 | \text{not } X_w = 1) \\ \implies \Delta_1 &= \sum_{v,w \in V, v \neq w} (1-p)^{2(n-2)+1} \end{aligned}$$

Using the fact that $(1-p)^{2(n-2)+1}$ edges are likely to exist if no edge exists with a vertex either from v or w .

$$\begin{aligned} \Delta_1 &= n(n-1)(1-p)^{2n-3} \\ \frac{\Delta_1}{E(X)^2} &= \frac{n(n-1)(1-p)^{2n-3}}{n^2(1-p)^{2(n-1)}} = \left(1 - \frac{1}{n}\right) * \frac{1}{1-p} \end{aligned}$$

This tends to 1, as $n \rightarrow \infty$.

Therefore, we have $\frac{1}{E(X)} - 1 + \frac{\Delta_1}{E(X)^2} \rightarrow 0$ as $n \rightarrow \infty$

Hence Proved. ■