

## Problem Set 1

1. Prove that  $\log n! = \theta(n \log n)$ . Use this to find out if  $\lceil \log n \rceil!$  and  $\lceil \log \log n \rceil!$  are polynomially bounded. A function  $f(n)$  is said to be polynomially bounded if there exists an integer  $k$  such that  $f(n) = O(n^k)$ .

2. Iterated logarithmic function is defined as  $\log^*(n) = \min\{i \geq 0 : \log^i n \leq 1\}$ . For example,  $\log^* 16 = 3$  where  $\log$  is taken w.r.t base 2. Which is asymptotically larger:  $\log(\log^* n)$  or  $\log^*(\log n)$ ?

3. Write an efficient algorithm that checks whether a given singly linked list contains a loop. A loop is a sequence of nodes  $v_1, v_2, \dots, v_k$  such that  $v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k \rightarrow v_1$ .

4. Consider a stack with an additional operation,  $MULTIPOP(S, k)$  which removes the  $k$  top objects of stack  $S$ , popping the entire stack if the stack contains fewer than  $k$  objects. The cost of the operation  $MULTIPOP(S, k)$  is  $k$ , while that of  $PUSH(S, x)$  and  $POP(S)$  is 1. Now consider a sequence of  $n$  stack operations on an initially empty stack, where each operation is either  $PUSH, POP$  or  $MULTIPOP$ . Prove that the total cost of these  $n$  operations is  $\theta(n)$ .

5. Design a data structure *SpecialStack* that supports all the stack operations like  $push(), pop(), isEmpty(), top()$  and an additional operation  $getMin()$  which should return minimum element from the *SpecialStack*. All these operations of *SpecialStack* must be  $O(1)$ . To design *SpecialStack*, you should only use standard Stack data structure and no other data structure like arrays, list, etc.

6. Describe a  $O(n)$ -time algorithm that, given a set  $S$  of  $n$  integers in sorted order and another integer  $x$ , determines whether or not there exist two elements in  $S$  whose sum is exactly  $x$ .

7. Given an unsorted array  $A$  of size  $n$  that may contain duplicates and a number  $k < n$ , design an  $O(n)$  algorithm that returns *true* if array contains duplicates within a distance of  $k$ , i.e. there exists  $i, j \in [0, n - 1]$  such that  $A[i] = A[j]$  and  $|i - j| < k$ .